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Uhrzeit:

Einbarung

Tag:

Mittwochs, 17:30 – 19:00

Veranstaltungen (und drei Ausweich-

Okt.: 10./17./24./31. (15Uhr30);

Nov.: 10./17./24./31. (15Uhr30);

Dez.: 1./3.(15Uhr30)/15;

Jan.: 12./14.(15Uhr30)/19./26.(Klausur)

– Raum: Hörsaal HS28; Hauptgebäude

Ausweichtermine: HS28; Hauptgebäude

**Klausur findet am
26.01.2011 um
17Uhr30 in HS28**

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**Die Anmeldung ist
ab sofort im PISWI
frei geschaltet!**



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Chair for Future Communication
Prof. Dr. K. Tutschku
Faculty for Computer Science

050069

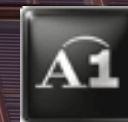
VO Netzwerktechnologie für Multimedia Anwendungen

Lecture 7: Multimedia-Networking

Prof. K. Tutschku (kurt.tutschku@univie.ac.at)

Bachelor Informatik (Medieninformatik)
WS 2010/11

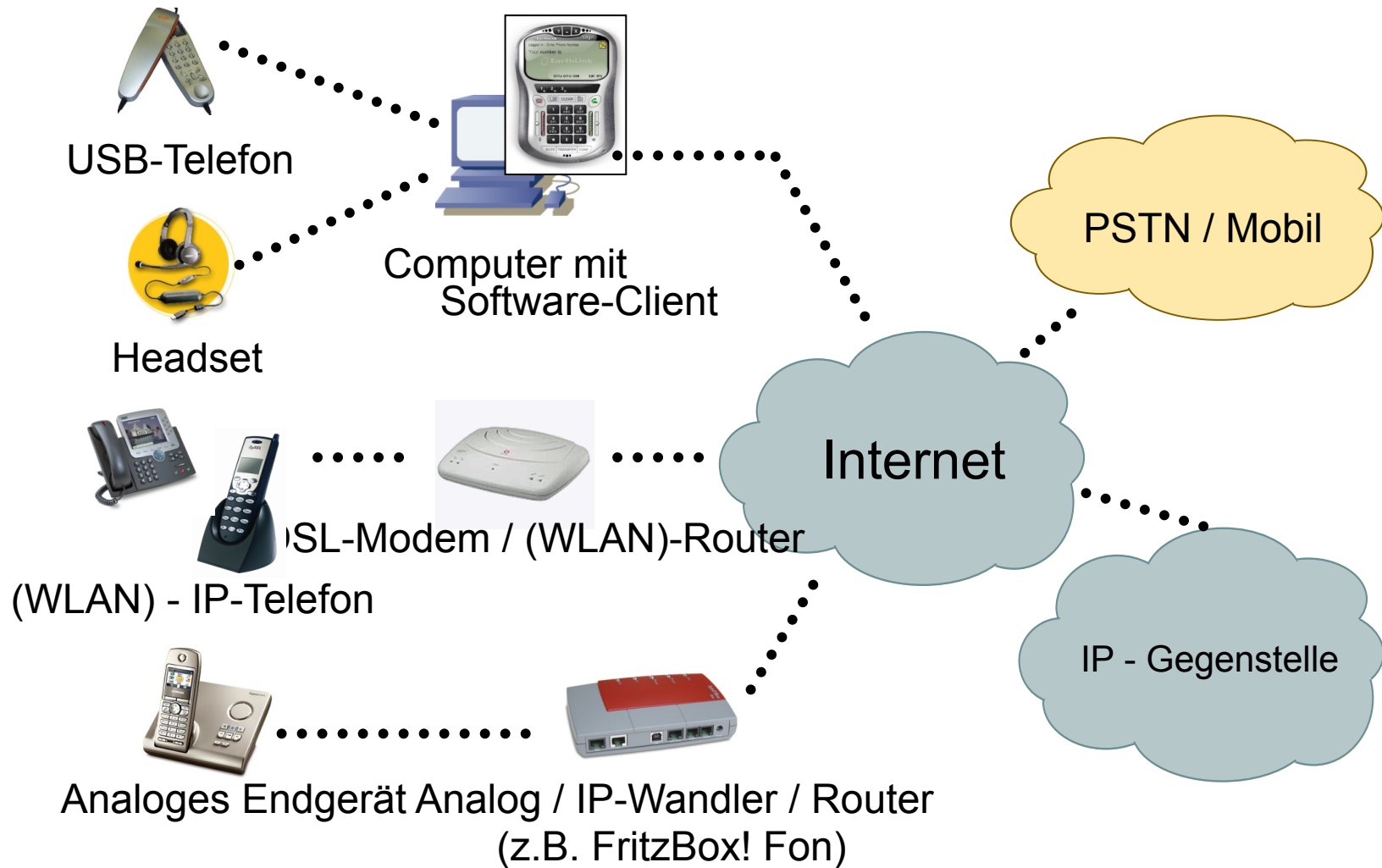
Endowed by





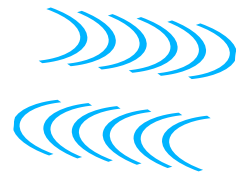
Overview:

- ▶ 2.1 Multimedia Networking Applications
- ▶ 2.2 Streaming stored audio and video
- ▶ 2.3 Real-time Multimedia: Internet Phone study
- ▶ 2.4 Protocols for Real-Time Interactive Applications
 - RTP, RTCP
- ▶ 2.5 IP Telephony, SIP, and H.323
- ▶ 2.6 Distributing Multimedia: content distribution networks





Mobilteil



Basisstation



DECT
Luftschnittstelle



Siemens
M34 USB
Adapter

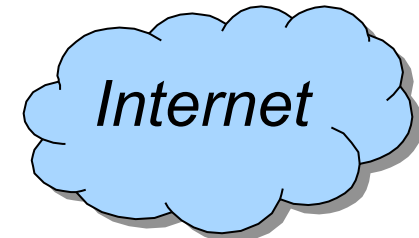
Skype-
Client

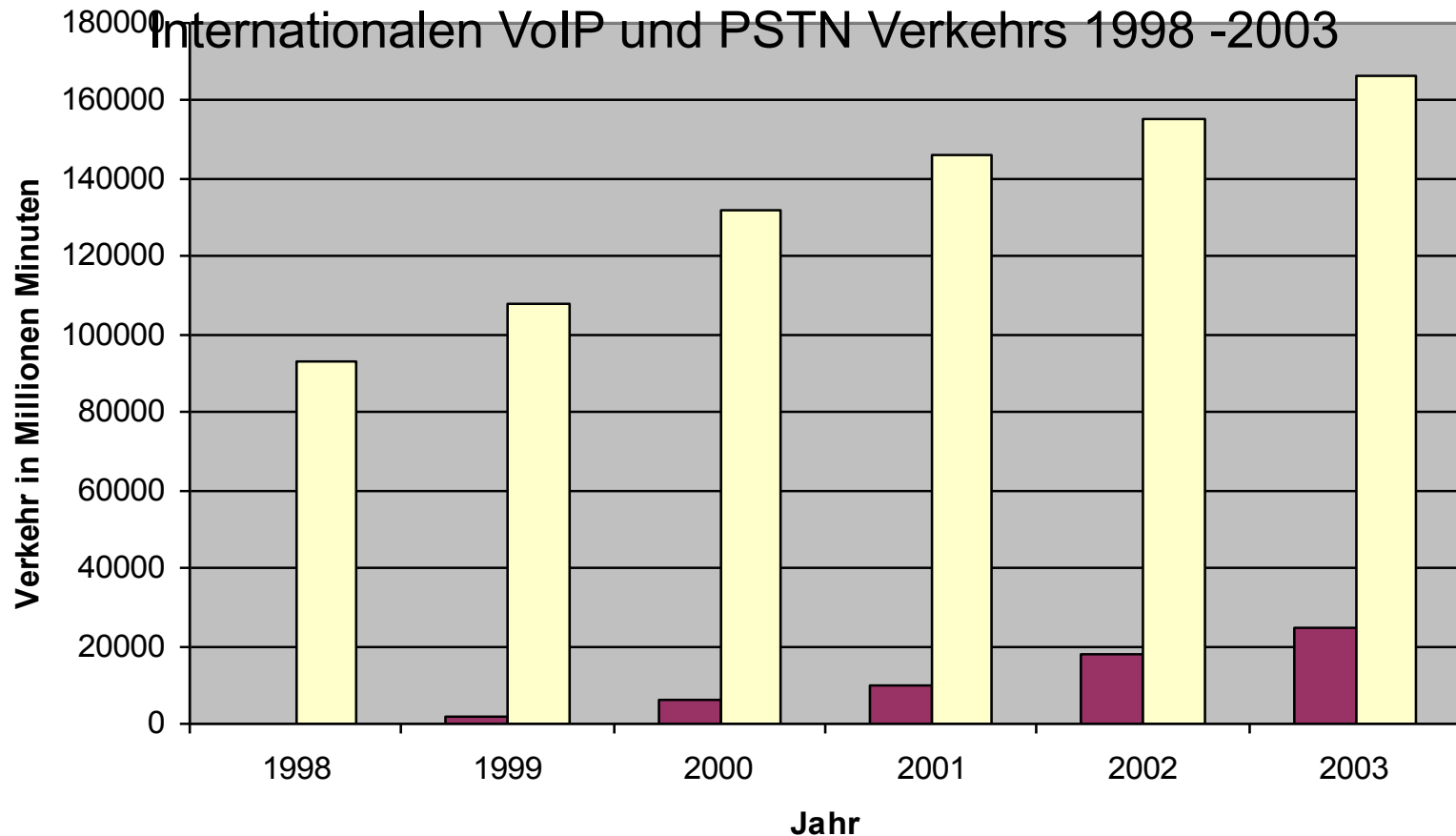


ISDN-
Anschluss



DSL-
Modem





VoIP Verkehr PSTN Verkehr

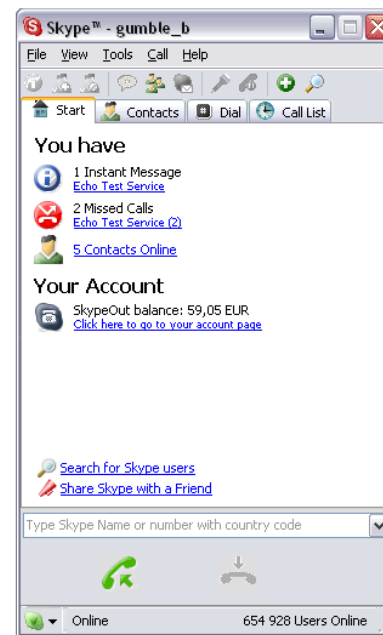
Quelle: TeleGeography 2004 (PriMetrica, Inc.)



- Softwareunterstützung für VoIP



SIP-fähige Clients



Skype



ICQ



Netmeeting



- **China**

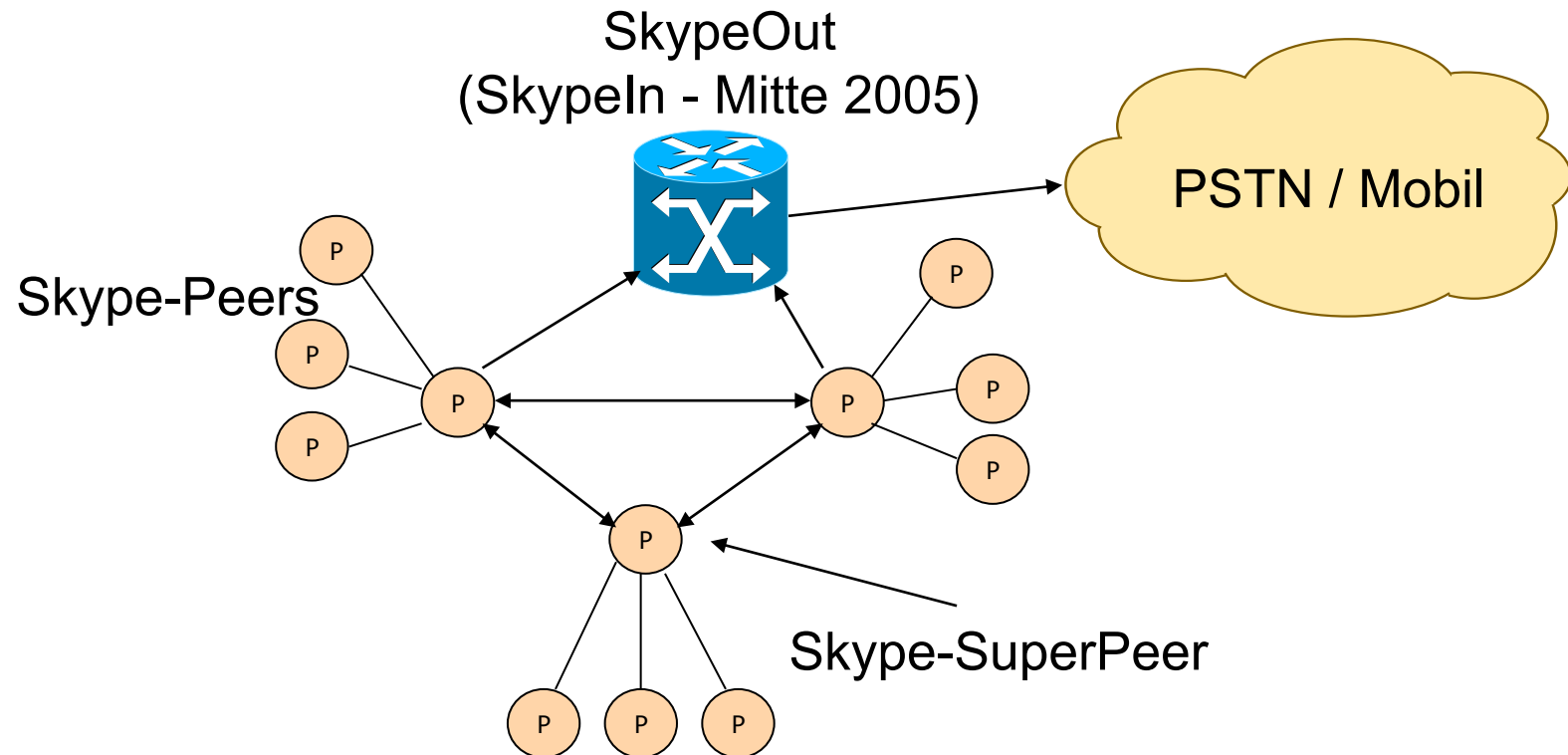
- 50% des Ferngesprächsverkehrs über VoIP
- China Unicom: 29 Mrd VoIP von insgesamt 54 Mrd. Gesprächsminuten
- Auch China Mobile setzt auf IP-Telefonie

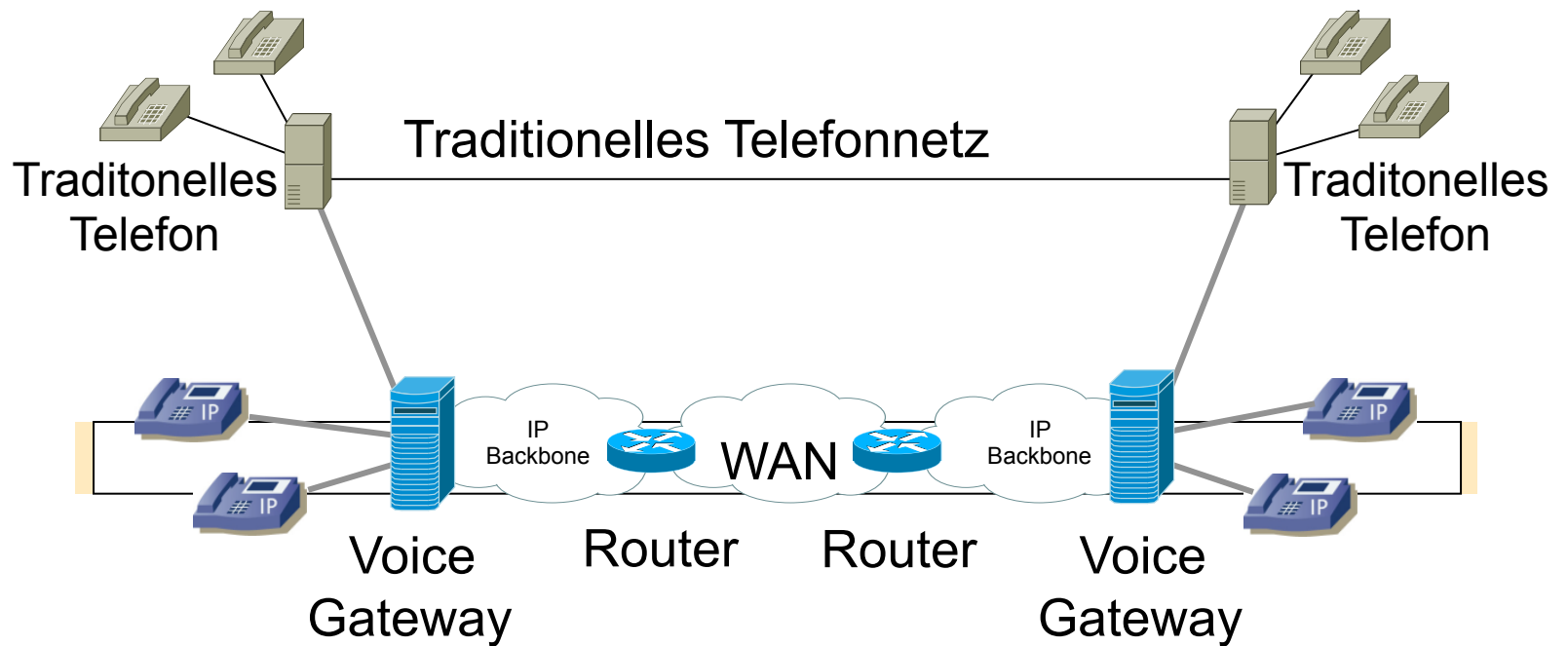
- **Australien**

- Mehr Internettelefonsysteme als klassische Nebenstellenanlagen verkauft
- Anstieg der Umsätze von VoIP-Systemen um 175% gegenüber 13,2% weniger Verkauf traditioneller Anlagen



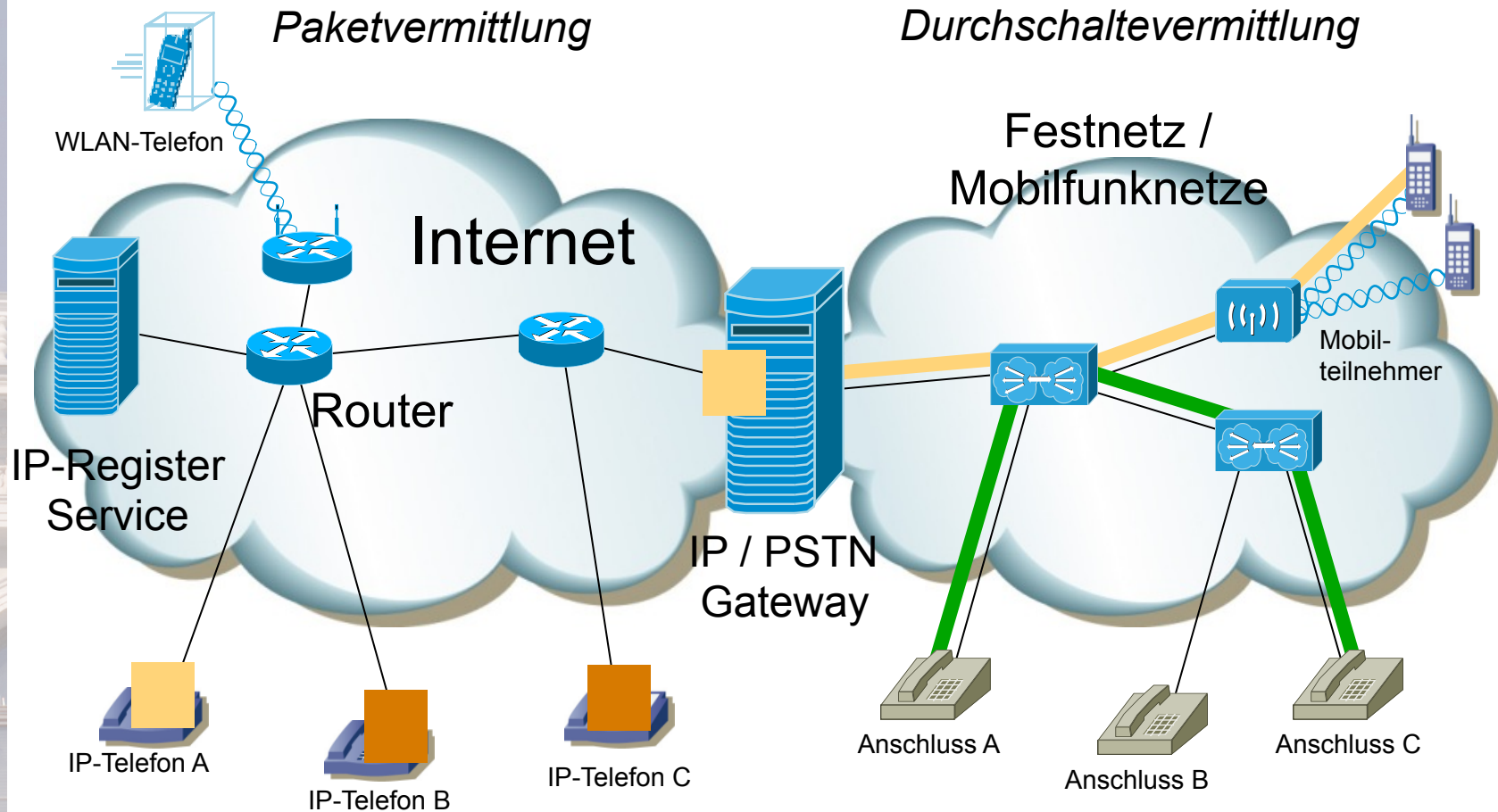
- Skype ist aus der KaZaa-Entwicklung hervorgegangen





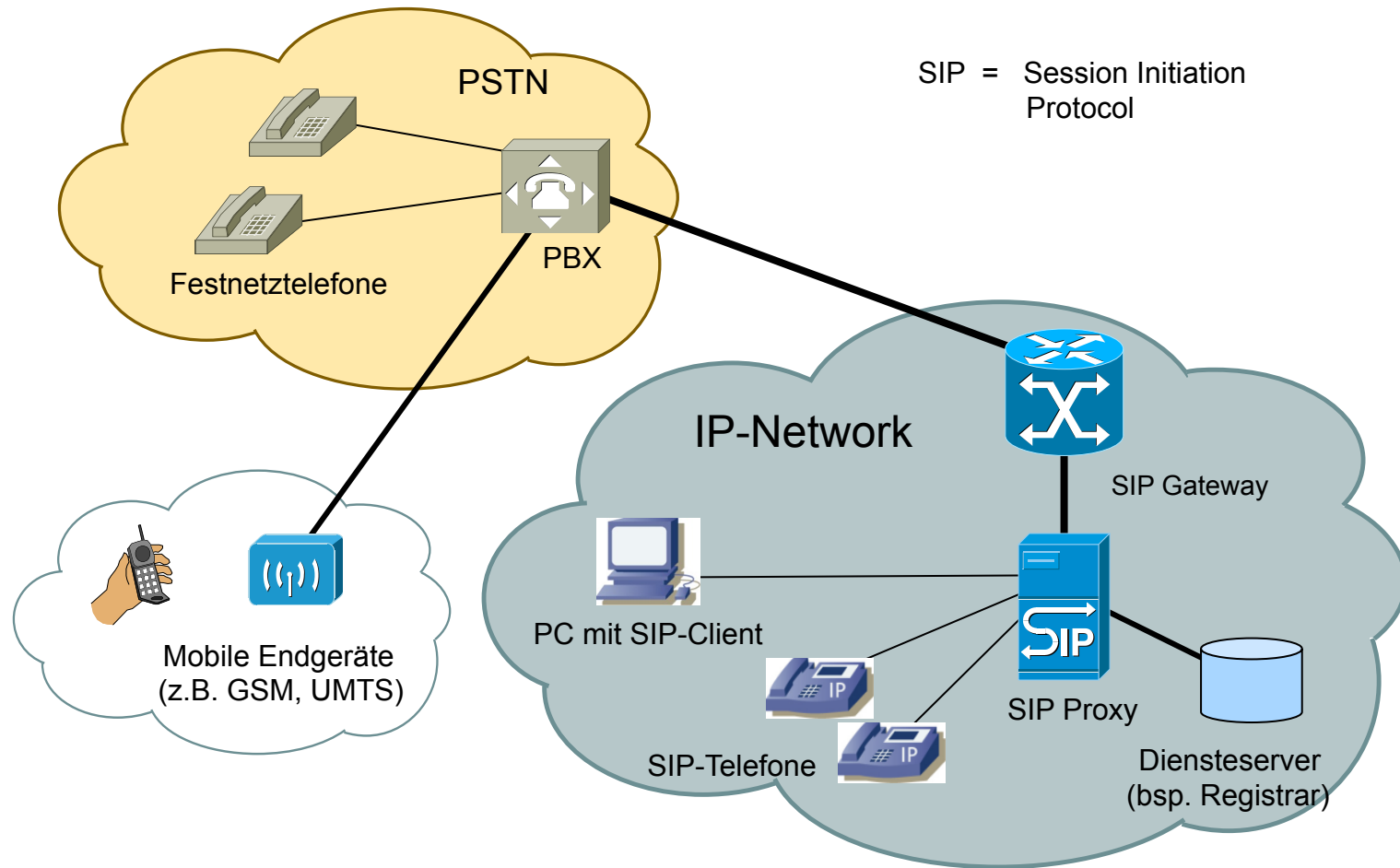


Vergleich klassische vs. VoIP-Telefonie





SIP = Session Initiation Protocol





- **In PSTN-Netzen ist Erreichbarkeit über eindeutige Telefonnummer gewährleistet**
- **Bei VoIP eine Vielzahl an unterschiedlichen Erreichbarkeits-merkmalen:**
 - IP-Adresse
 - SIP-String / VoIP-Nummer
 - Festnetznummer (Wohnortsbereich, RegTP Nummerngasse (032), anbieterspezifisch (z.B. sipgate 01801))
- **Aufgrund der nicht garantierten Ausfallssicherheit und Zuordnung auch keine Erreichbarkeit gewährleistet**
(Notrufnummernproblematik)



- **Session Initiation Protocol (SIP)**
 - Universales Initialisierungsprotokoll für den Verbindungsaufbau
 - Nicht festgelegt auf ein spezielles Medium (Video, Audio)
 - Sehr einfaches Protokoll, dadurch leicht in Hardware implementierbar
 - IETF - Standard
- **H.323**
 - Umfassender Standard für Multimediaübertragung, dadurch kompliziert zu implementieren
 - Spezialisiert auf Sprach- und Multimediadienste
 - Objektorientiert, basiert auf QSIG - Standard



- Bestehende Anbieter für SIP-Telefonie



- Serverlösungen für SIP-Telefonie



Cisco CallManager



- **Multiparty Multimedia Session Control (MMUSIC) working group of the Internet Engineering Task Force (IETF)**
 - RFC 2543 (1999)
 - RFC3261 (updated)
- **SIP long-term vision**
 - All telephone calls and video conference calls take place over the Internet
 - People are identified by names or e-mail addresses, rather than by phone numbers.
 - You can reach the callee, no matter where the callee roams, no matter what IP device the callee is currently using.

- **Works according to the client – server principle**

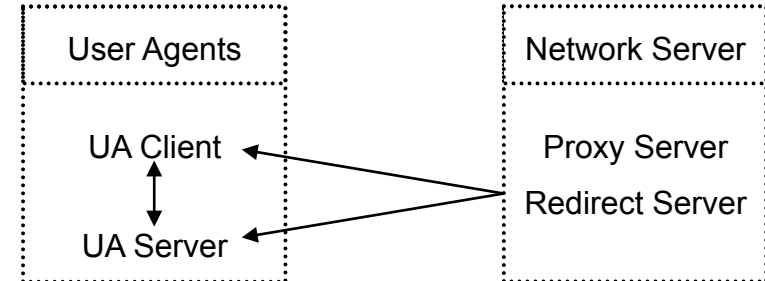




- **Determine current IP address of callee.**
 - Provides mnemonic sip addresses [sip:]<user>@(<host>|<domain>)
 - me@132.187.10.51
 - sip:0123-45-67-89@telefon.com
 - Additional parameters e.g. for transport protocol <URI>;tag1;tag2;...
 - Maps them to current IP addresses
- **Setting up a call**
 - Provides mechanisms for caller to let callee know she wants to establish a call
 - Provides mechanisms so that caller and callee can agree on media type and encoding.
 - Provides mechanisms to end call
- **Call management**
 - Add new media streams during call
 - Change encoding during call
 - Invite others (multi-party conference)
 - Transfer and hold calls



- **User agents**
 - Want to communicate with each other
 - Examples
 - Application on a user's computer
 - Cellular phone
 - PSTN gateway
- **SIP proxies or SIP redirect servers**
 - Help to find other users
- **Registrars**
 - Map mnemonic part of sip addresses to IP numbers
 - Usually collocated with SIP server
- **SIP gateways for interoperability with PSTN**





- **Client side methods**

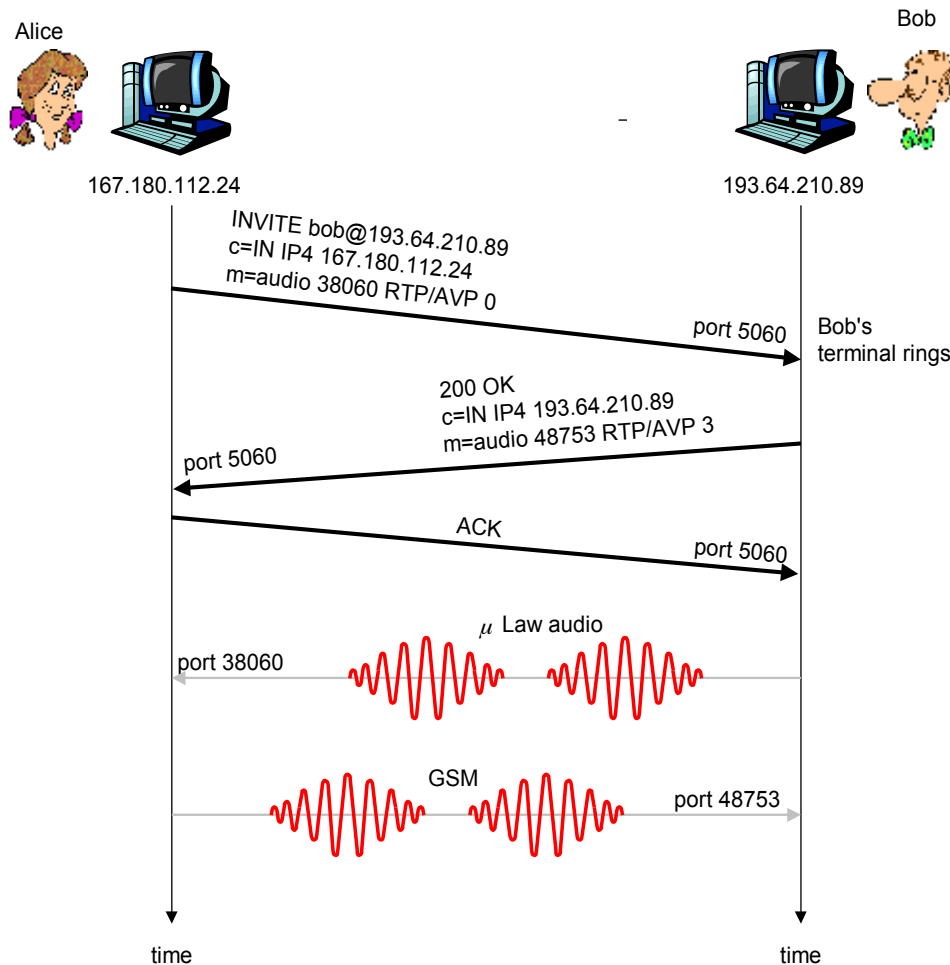
- INVITE: request for session setup
- ACK: acknowledgement of a response
- CANCEL: cancels requests
- BYE: terminates connection
- OPTIONS: checks capabilities of peer
- REGISTER: registers at SIP registrar

- **Server side status codes (6 categories, similar for other protocols)**

- 1xx: information about progress of transaction
 - E.g. 180 „Ringing“, 181 „Call is Being Forwarded“, 182 „Queued“
- 2xx: success of transaction (200 „OK“)
- 3xx: deviation, more effort required
 - E.g. 301 „Moved Permanently“, 302 „Moved Temporarily“
- 4xx: flaw in request
 - E.g. 404 „Not Found“, 420 „Bad Extension“, 486 „Busy Here“
- 5xx: flaw at server side
 - E.g. 500 „Internal Server Error“, 504 „Server Time Out“
- 6xx: general fault
 - E.g. 600 „Busy Everywhere“, 603 „Decline“, 604 „Does Not Exist Anywhere“



Setting up a call to a known IP address



- Alice's SIP invite message indicates her port number & IP address. Indicates encoding that Alice prefers to receive (PCM ulaw)
- Bob's 200 OK message indicates his port number, IP address & preferred encoding (GSM)
- SIP messages can be sent over TCP or UDP; here sent over RTP/UDP.
- Default SIP port number is 5060.



- **When Bob starts SIP client, client sends SIP REGISTER message to Bob's registrar server
(similar function needed by Instant Messaging)**

Register Message:

```
REGISTER sip:domain.com SIP/2.0  
Via: SIP/2.0/UDP 193.64.210.89  
From: sip:bob@domain.com  
To: sip:bob@domain.com  
Expires: 3600
```



Name Translation and User Location

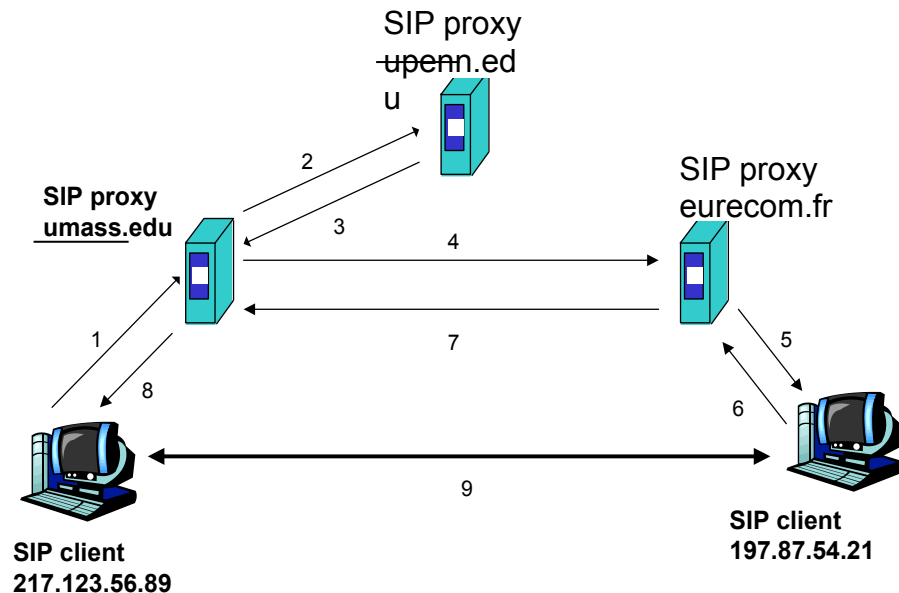


- **Caller wants to call callee, but only has callee's name or e-mail address.**
- **Need to get IP address of callee's current host:**
 - User moves around
 - DHCP protocol
 - User has different IP devices (PC, PDA, car device)
- **Caller asks SIP server (like DNS server)**
- **Result can be based on:**
 - Time of day (work, home)
 - Caller (don't want boss to call you at home)
 - Status of callee (calls sent to voicemail when callee is already talking to someone)
- **Alice sends invite message to her proxy server containing Bob's address "sip:bob@domain.com"**
- **Proxy responsible for routing SIP messages to callee possibly through multiple proxies.**
- **Callee sends response back through the same set of proxies.**
- **Proxy returns SIP response message to Alice containing Bob's IP address**



Caller jim@umass.edu places a call to keith@upenn.edu

- (1) Jim sends INVITE message to umass SIP proxy.
- (2) Proxy forwards request to upenn proxy/registrar server.
- (3) upenn server returns redirect response, indicating that it should try keith@eurecom.fr



- (4) umass proxy sends INVITE to eurecom SIP proxy/registrar.
- (5) eurecom proxy/registrar forwards INVITE to 197.87.54.21, which is running keith's SIP client.
- (6-8) SIP response sent back
- (9) media sent directly between clients.
- Note: also a SIP ack message, which is not shown.

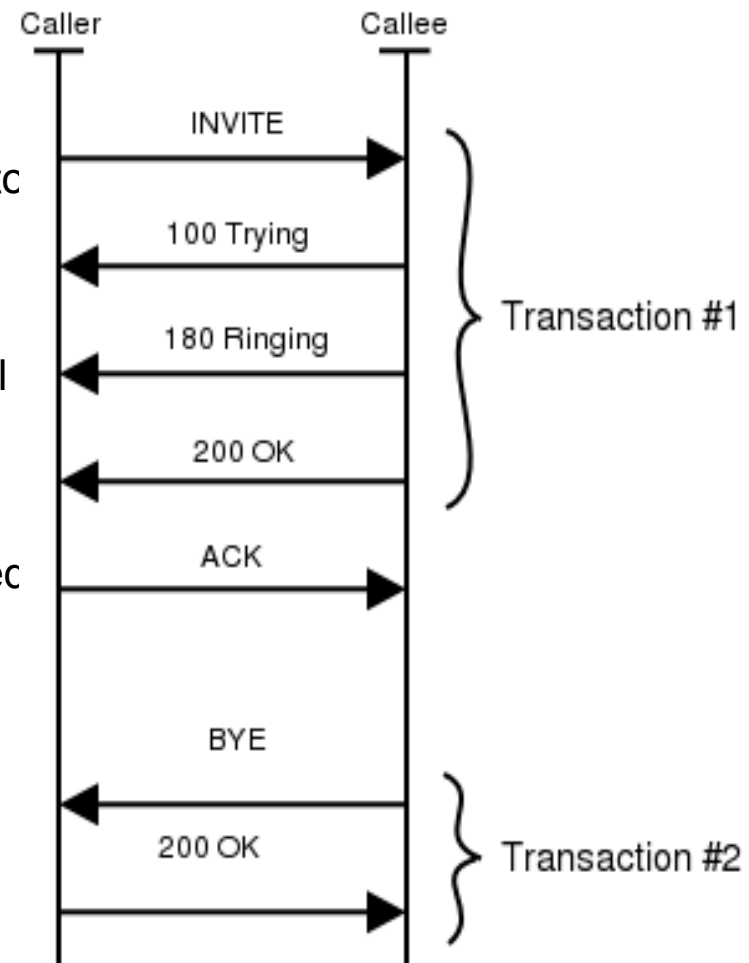


- **Codec negotiation:**
 - Suppose Bob doesn't have PCM ulaw encoder.
 - Bob will instead reply with 606 Not Acceptable Reply and list encoders he can use.
 - Alice can then send a new INVITE message, advertising an appropriate encoder.
- **Rejecting the call**
 - Bob can reject with replies "busy," "gone," "payment required," "forbidden".
- **Media can be sent over RTP or some other protocol.**



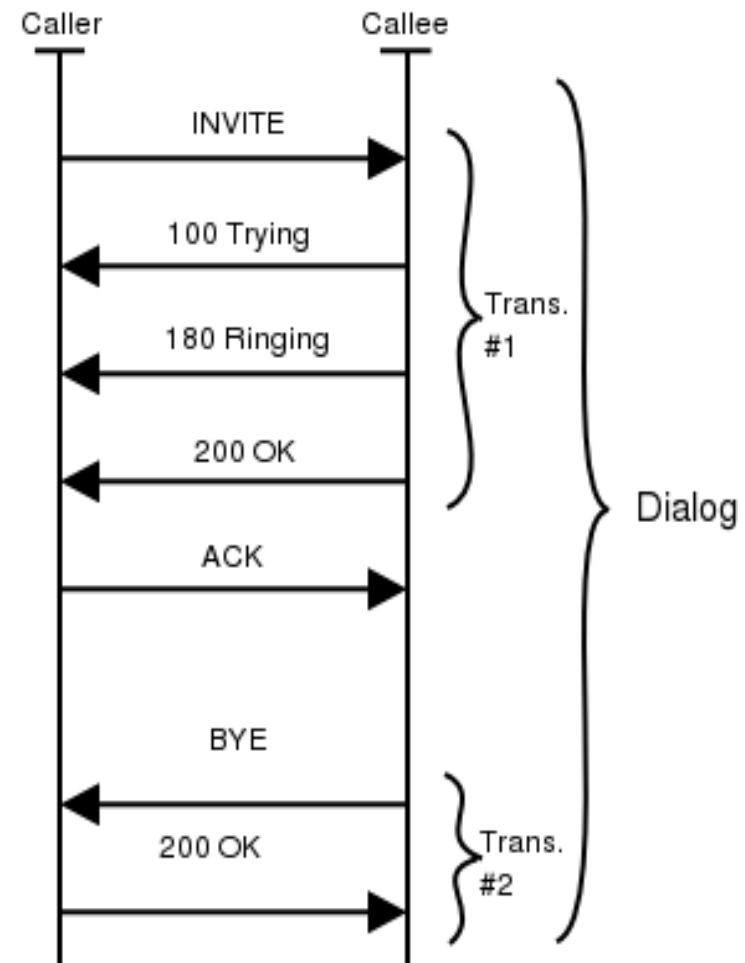
- **Transaction**

- Sequence of SIP messages
- One request and all responses to that request
- Zero or more provisional responses and one or more final responses (e.g. when proxy server forks)
- Final positive ack may be omitted but not final negative ack
- Transaction identifier
 - hash of all important message header fields (obsolete)
 - Directly contained in msg.



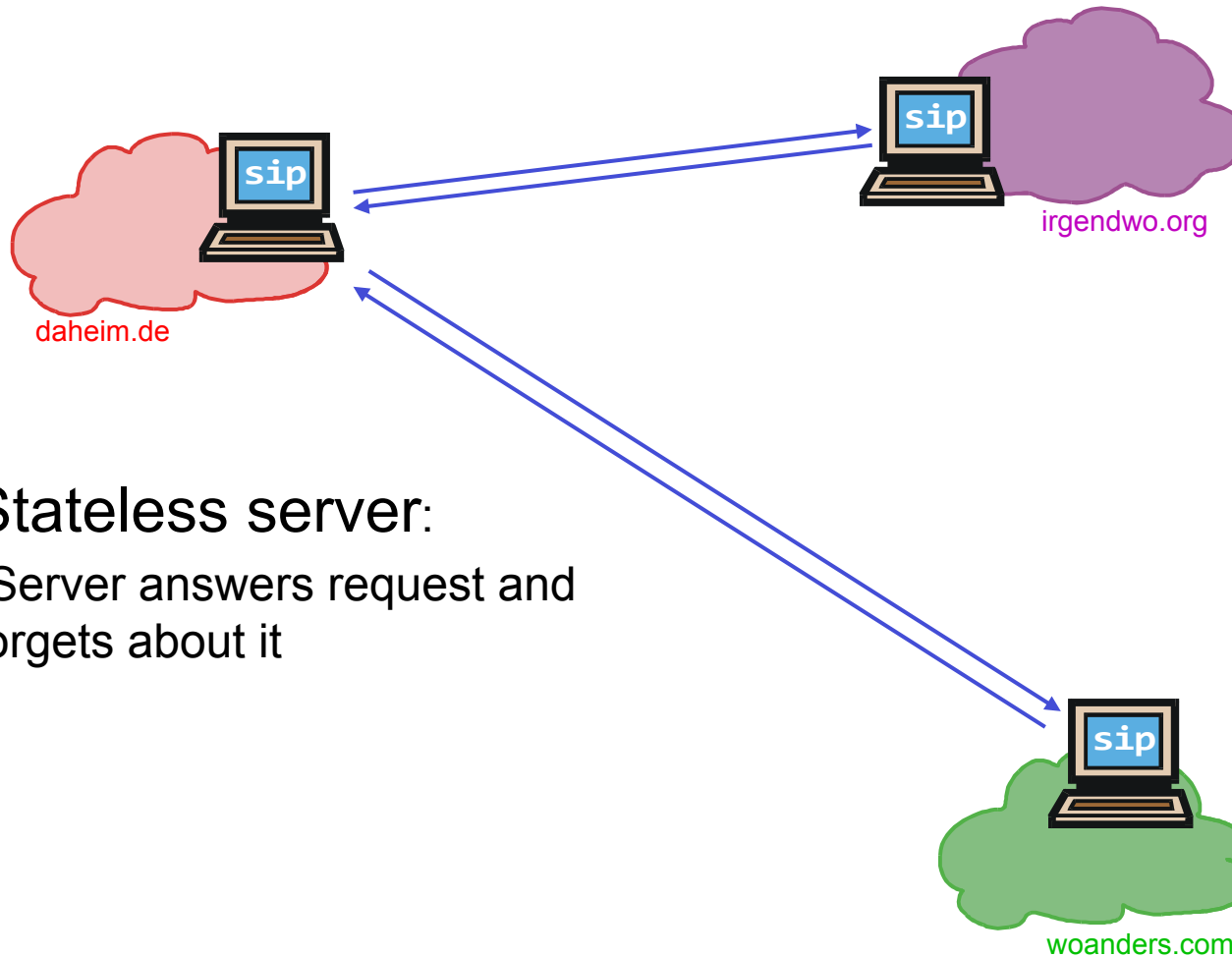


- Represents a peer-to-peer SIP relationship between two user agents
- *A dialog is a sequence of transactions*
- Dialogs are identified using Call-ID, From tag, and To tag
- CSeq header field numbers request / messages within a dialog.
- The number must be monotonically increased for each message sent within a dialog otherwise the peer will handle it as out of order request or retransmission





Network Server (1): Redirect Server

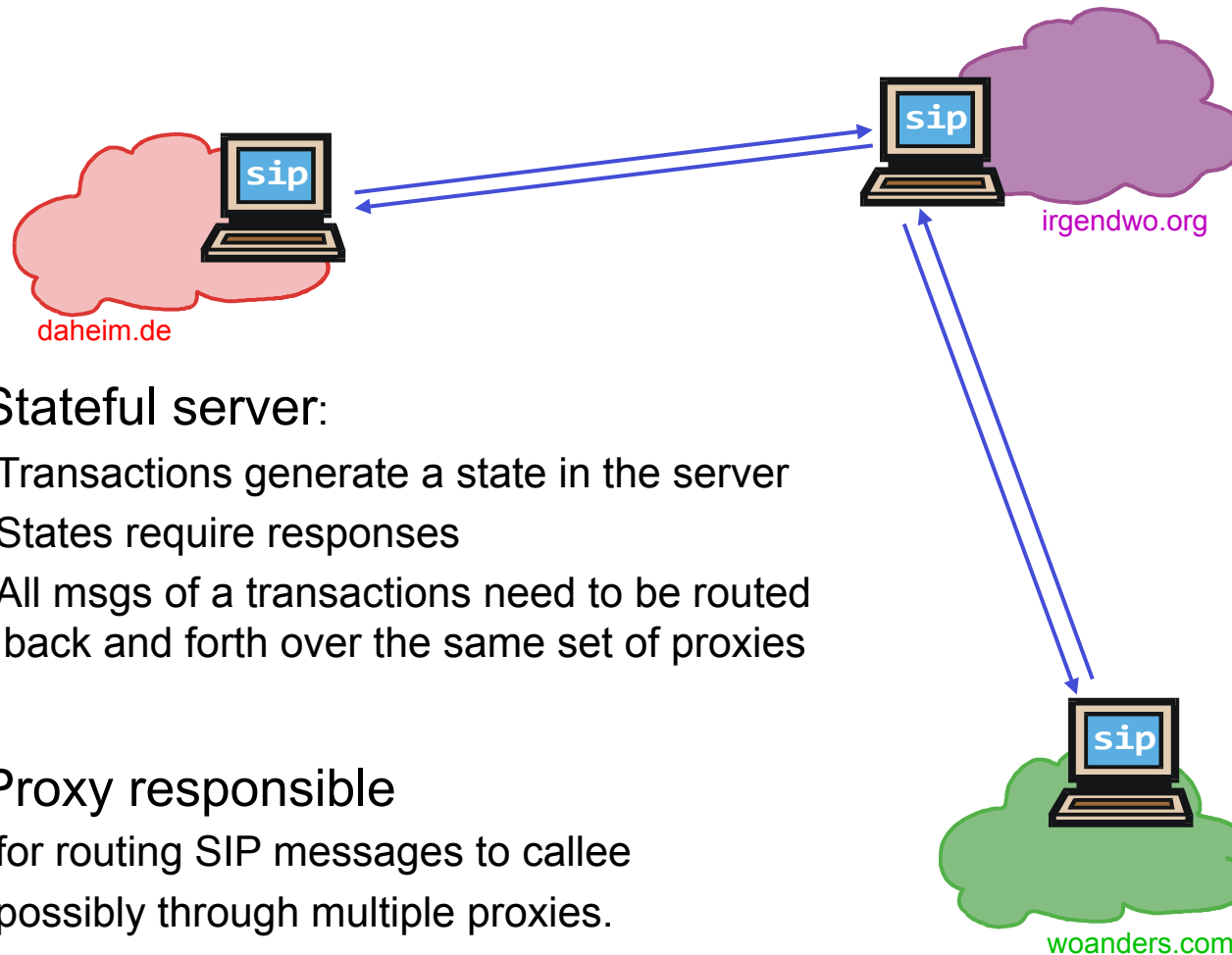


Stateless server:

- Server answers request and forgets about it



Network Server (2): Proxy Server



Stateful server:

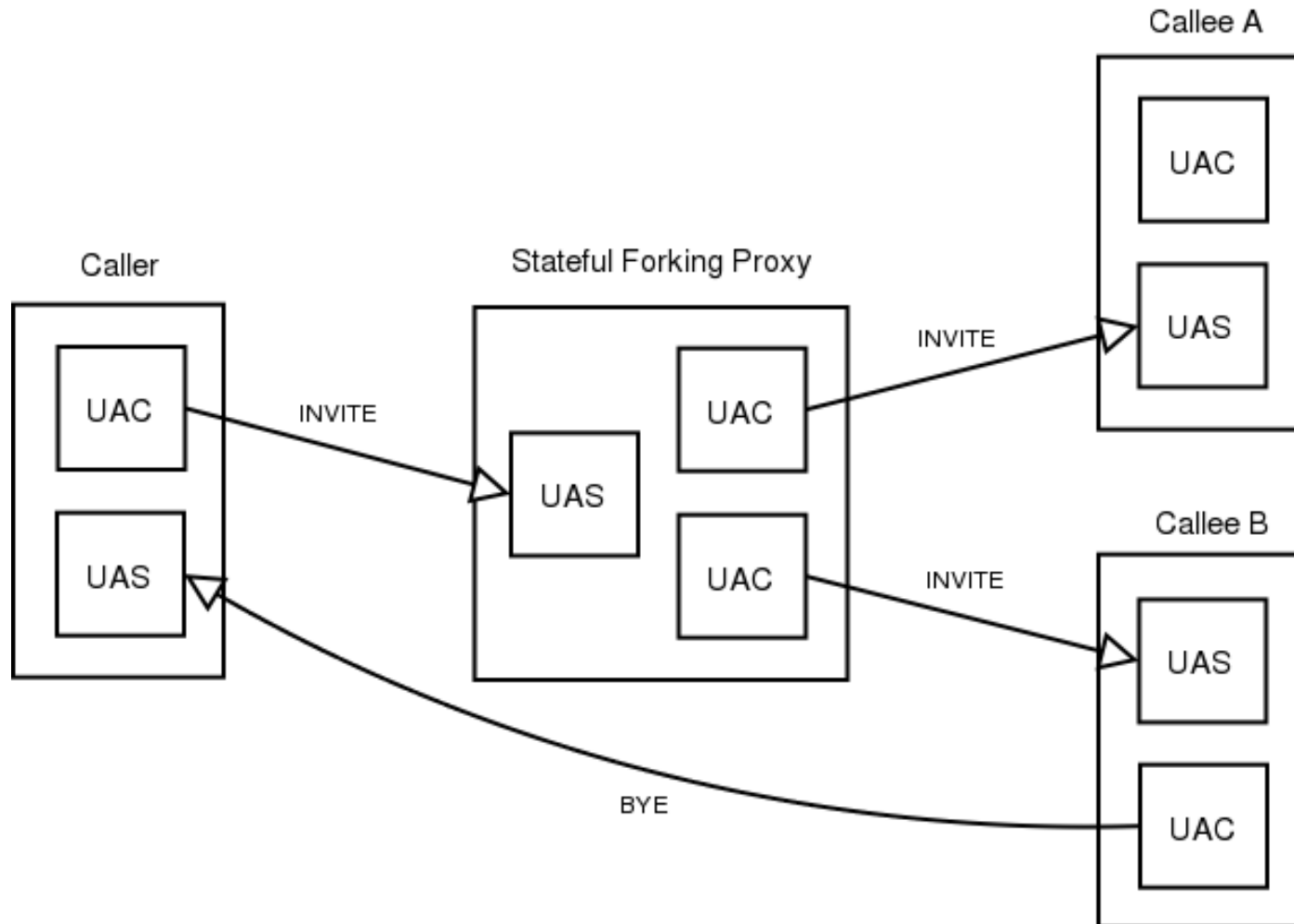
- Transactions generate a state in the server
- States require responses
- All msgs of a transactions need to be routed back and forth over the same set of proxies

Proxy responsible

- for routing SIP messages to callee
- possibly through multiple proxies.



Forking Proxy



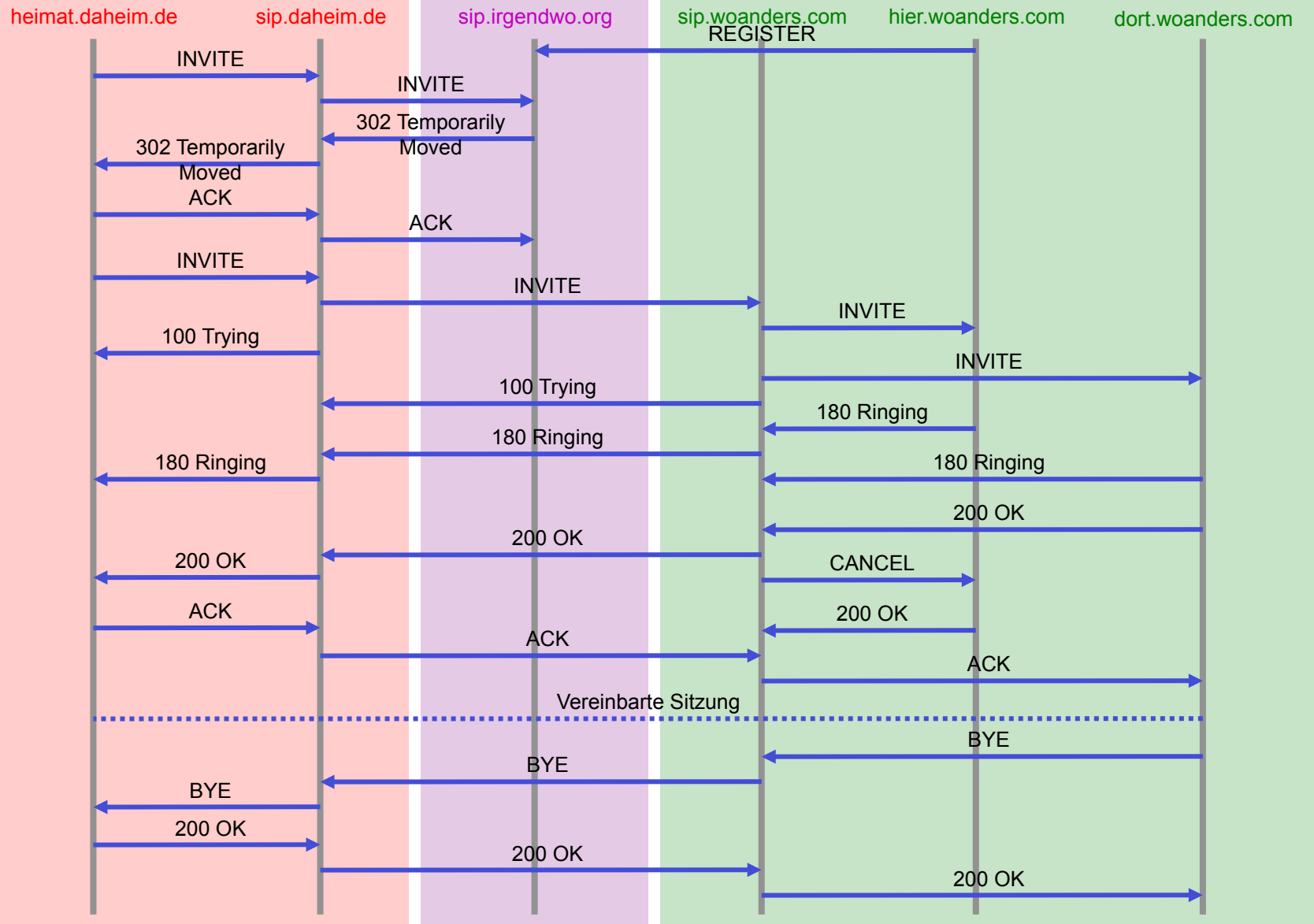


- **Stateless proxy**
 - Simple message forwarder
 - Do not take care of transactions
 - Simpler and faster than stateful proxies
- **Stateful proxy**
 - Most of today's proxies are stateful
 - Creates a state upon reception of a request and keeps it until the transaction finishes
 - Advanced functionalities
 - Absorb retransmissions
 - Advanced message routing
 - Forking (msg forwarding to different destinations)
 - Recursive traversal (try different locations to find a user)
 - Accounting
 - NAT traversal aid



How do we find appropriate SIP proxy?

- **Local SIP server usually knows users within its domain**
- **How does it find appropriate SIP proxy?**
- **Remember: Domain Name System!**
 - Translates mnemonic addresses into IP numbers
 - Hierarchical structure
 - Resource records specify request (Name, Value, Type, TTL)
 - Type may be, e.g., A, NS, CNAME, MX, ...
 - MX returns mail exchange server
 - Similar type for SIP server (SRV?)





Example of SIP message

```
INVITE sip:bob@domain.com SIP/2.0
Via: SIP/2.0/UDP 167.180.112.24
From: sip:alice@hereway.com
To: sip:bob@domain.com
Call-ID: a2e3a@pigeon.hereway.com
Content-Type: application/sdp
Content-Length: 885

c=IN IP4 167.180.112.24
m=audio 38060 RTP/AVP 0
```

- Here we don't know Bob's IP address. Intermediate SIP servers will be necessary.
- Alice specifies in Via: header that SIP client sends and receives SIP messages over UDP
- Alice sends and receives SIP messages using the SIP default port number 5060.

- HTTP-like message syntax
- Format: Unicode-Text in UTF-8-Codierung (8-bit unicode transformation format, RFC3629)
- sdp = session description protocol
- Call-ID is unique for every call.



- Von MMUSIC / IETF, veröffentlicht als RFC 2327 (April 1998)
 - Textformat: **<Bezeichner>=<Wert>**
 - Informationen:
 - Name und Zweck der Session
 - Zeit(en), in der (denen) die Session aktiv ist
 - Medien, die an der Session beteiligt sind
 - Informationen zum Empfang dieser Medien (Adresse, Port, Format, ...)
- zusätzlich:
- Bandbreite
 - Ansprechpartner

- **Transaktion besteht aus Request - Response (- Acknowledgement)**

- **Aufbau:**

<Start-Zeile>

<Header_1>:<Wert_1>

■ ■ ■

<Header_n>:<Wert_n>

Kopf

< ... >

Körper

- <Start-Zeile> = <Methode> <Request-URI> <SIP-Version> Request
| <SIP-Version> <Status-Code> <Begründung> Response



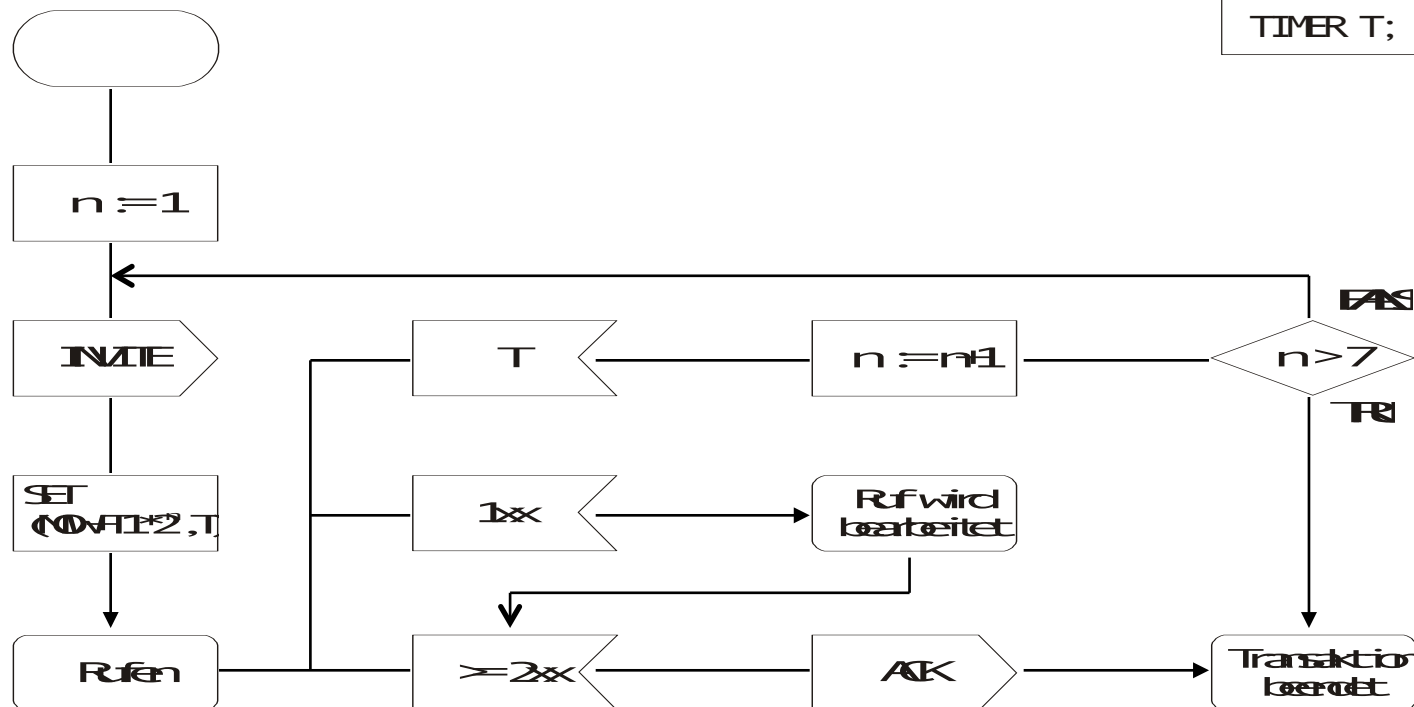
- **SIP erweiterbar um neue Header durch hierarchische Namensgebung oder Registrierung bei der Internet Assigned Numbers Authority (IANA)**
- **Wichtige Standard-Header:**
 - From: der anrufende Partner
 - To: der gerufene Partner
 - Call-ID: eindeutiger Bezeichner der Sitzung
 - CSeq: Bezeichner für eine Transaktion (Nr + Methode)
 - Via: Route der Nachricht (bisher)
 - Contact: Alternativ-Adresse(n)
 - Require: Erforderliche Optionen
 - Unsupported: Nicht unterstützte Optionen
 - Content-Type/-Length/-Encoding: Nachrichten-Körper



SDL Diagram for Client Finite State Machine

Aus Sicht des Client, Server-Sicht analog!

DCL
n Integer;
TIMER T;





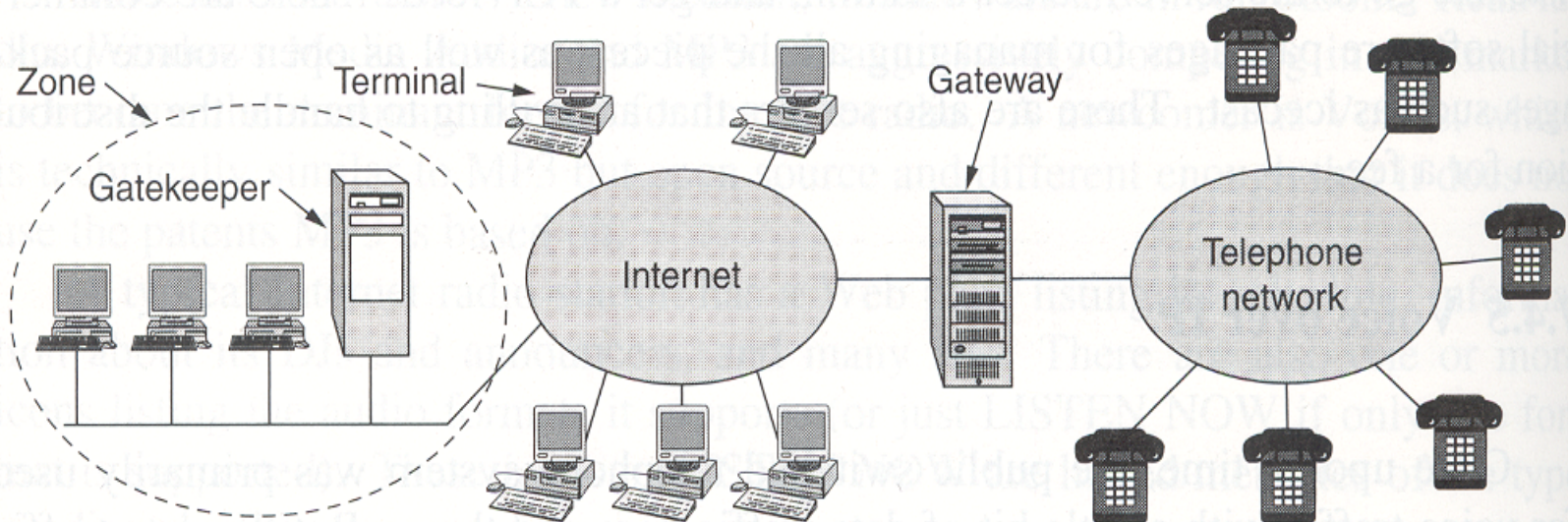
- **Nachrichtenfluss bei den übrigen Methoden:**
 - Request/Response bei BYE, CANCEL, OPTIONS, REGISTER analog
 - Retransmission des Request ebenfalls mit Back-Off-Algorithmus
 - Keine Bestätigung mit ACK
 - ⇒ keine period. Wiederholung des Response
 - Wiederholung des Response nur nach Wiederholung des Request
- **Bemerkung:**
Wiederholungen bei TCP i.d.R. unnötig, da zuverlässige Verbindung



- **Multimedia communication standard by ITU**

- Development start in 1996 to make IP telephony equipment from different vendors interoperable
- Revised in 1998
- Basis for first widespread Internet telephony systems
- Not a single protocol but an architectural overview comprising many different protocols and other stuff
 - Speech codecs
 - Call setup
 - Signalling
 - Data transport
 - Interoperability with telephone network
- Network elements
 - Terminals
 - Gatekeeper
 - Zone
 - Gateway







- **Codecs**
 - Requirement: G.711 (64 kbps PCM voice, uncompressed)
 - Many others permitted, e.g., G.723.1 (predictive coding to compress speech to 24 or 20 bytes / 30 ms = 6.4 and 5.3 kbps, compression factor 10 and 12!)
- **H.245**
 - Capability information and parameter negotiation
 - Codecs, bit rates, ...
- **RTP and RTCP required for data transport**
- **ITU Q.931**
 - Standard telephony signalling
 - Establishing and releasing connections
 - Providing dial tones, making ringing sounds
 - Rest of standard telephony features



- **H.225**

- Used for communication with gatekeeper
- Manages PC-to-gatekeeper channel „RAS“ (Registration/ Admission/Status), allows terminals
 - To join and leave the zone
 - To request and return bandwidth
 - To provide status updates and other stuff

Speech	Control			
G.7xx	RTCP	H.225 (RAS)	Q.931 (Call signaling)	H.245 (Call control)
RTP				
UDP			TCP	
IP				
Data link protocol				
Physical layer protocol				

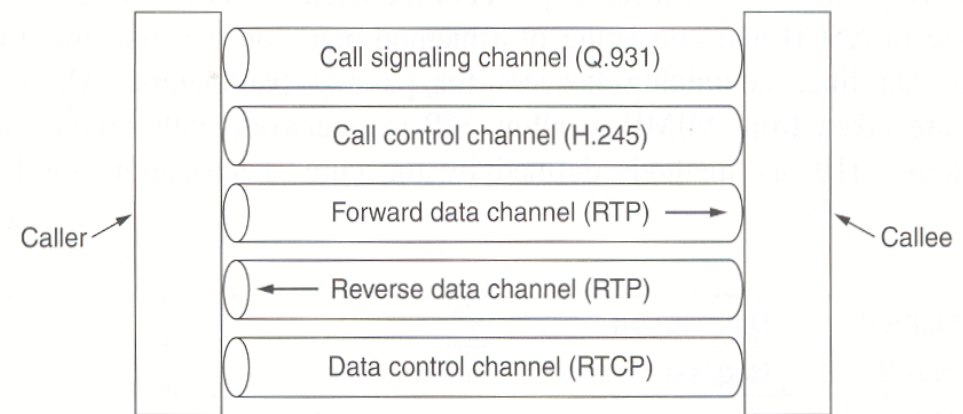


- **PC discovers gatekeeper by broadcasting a UDP gatekeeper discovery packet to port 1718**
- **Gatekeeper responds and PC learns gatekeeper's IP address**
- **PC registers at gatekeeper with a UDP msg**
- **After acceptance, PC requests bandwidth with a RAS msg in UDP**
 - QoS achieved by local admission control
- **If successful, connection setup may start over TCP towards gatekeeper**
 - PC sends Q.931 SETUP (including telephone number / IP address)
 - Gatekeeper responds with Q.931 CALL PROCEEDING and forwards SETUP towards gateway
 - Gateway is half computer and half telephone switch
 - Forwards SETUP in appropriate way and forwards Q.931 ALERT to PC (ringing has begun)
 - If remote peer picks up the telephone, a CONNECT msg is sent back to the PC



H.323 Signaling Example

- **From now on, the gatekeeper is no longer in the loop and communication is directly PC – Gateway**
 - H.245: capability and parameter (e.g. codec) negotiation
 - Different codecs may be used for both directions
 - Two unidirectional data channels are set up
 - Data flow begins over RTP
 - RTCP helps to control congestion and to synchronize audio and video
- **Q.931 channel is finally used to tear down the connection**
- **When call is terminated, PC contacts gatekeeper to release the reserved bandwidth**



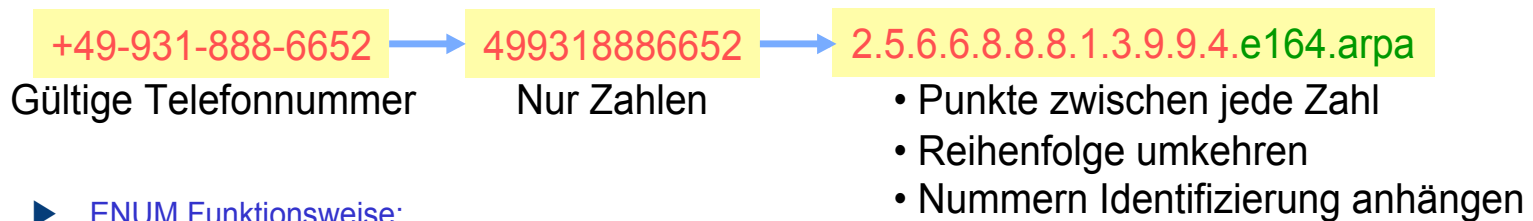


Comparison SIP vs. H.323

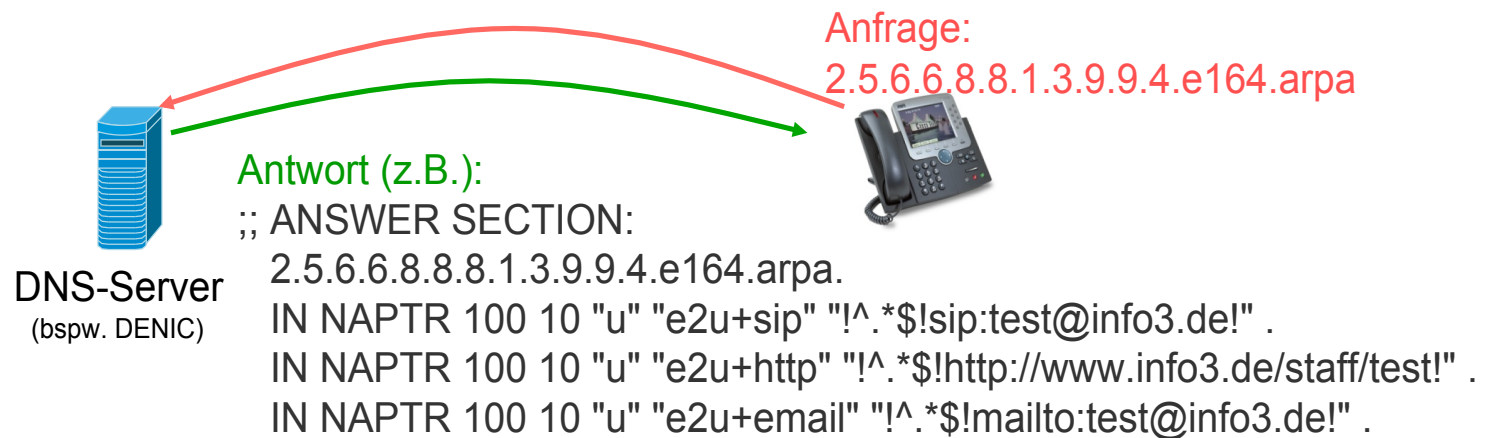
- **SIP is a single component.**
Works with RTP, but does not mandate it. Can be combined with other protocols and services.
- **SIP comes from IETF: Borrows much of its concepts from HTTP.**
- **SIP has a Web flavor.**
 - States kept in end devices
 - Proxy just helps but it is not required
- **SIP uses the KISS principle: Keep it simple stupid.**
- **H.323 is a complete, vertically integrated suite of protocols for multimedia conferencing: signaling, registration, admission control, transport and codecs.**
- **H.323 comes from the ITU (telephony).**
- **H.323 has a telephony flavor.**
 - End devices (telephones) are very primitive
 - States kept in network devices in PSTN; H.323 uses also a gatekeeper



- ▶ tElephone NUMber Mapping (ENUM): Anwendung des DNS zur Übersetzung von Telefonnummern in Internet-Adressen (RFC3761)
- ▶ Klassische Telefonnummer durch ITU E.164 festgelegt
- ▶ ENUM Umsetzung:
 - Eingabe: E.164 konforme Nummer
 - Ausgabe: gültige „absoluteURI“ (RFC2396)
 - Beispiel:



- ▶ ENUM Funktionsweise:
 - NAPTR (Naming Authority Pointer) im Resource Record des DNS
 - IN: Internet





Next Generation Network (NGN) – im engeren Sinne



- **Ziel: Konvergenz von Telefon- und Datennetzen**

- paketvermittelte Verbindungen zwischen zwei und mehr Teilnehmern
 - Zusammenarbeit zwischen der leitungsvermittelnden und der paketorientierten Domäne
 - eine End-to-End-Aushandlung der Dienstgüte (Quality of Service)
 - dienstabhängige Kostenabrechnung
 - Bereitstellung der Heimnetzumgebung in Fremdnetzen
 - Unterstützung verschiedener Medientypen
 - schnelle und flexible Erstellung von Diensten durch Service Enabler (vordefinierte Dienstbausteine)
 - Dienste sollen unabhängig vom Zugangsnetz sein
- Netzwerktechnologie für Multimedia Anwendungen

- **Ansätze in der Standardisierung**

- TC TISPAN (Technical Committee Telecoms & Internet converged Services & Protocols for Advanced Networks) von ETSI (European Telecommunications Standards Institute)
- IP Multimedia Subsystem (IMS) von 3GPP (3rd Generation Partnership Project, weltweite Kooperation von 5 Normierungsorganisationen zur Standardisierung im Mobilfunk, darunter auch ETSI)
- Y.2001 der ITU-T (International Telecommunication Union – Standardization Sector)
- Mehrheitliche Benutzung von SIP

Quelle: Wikipedia, 29.6.2009
WS 2010/11



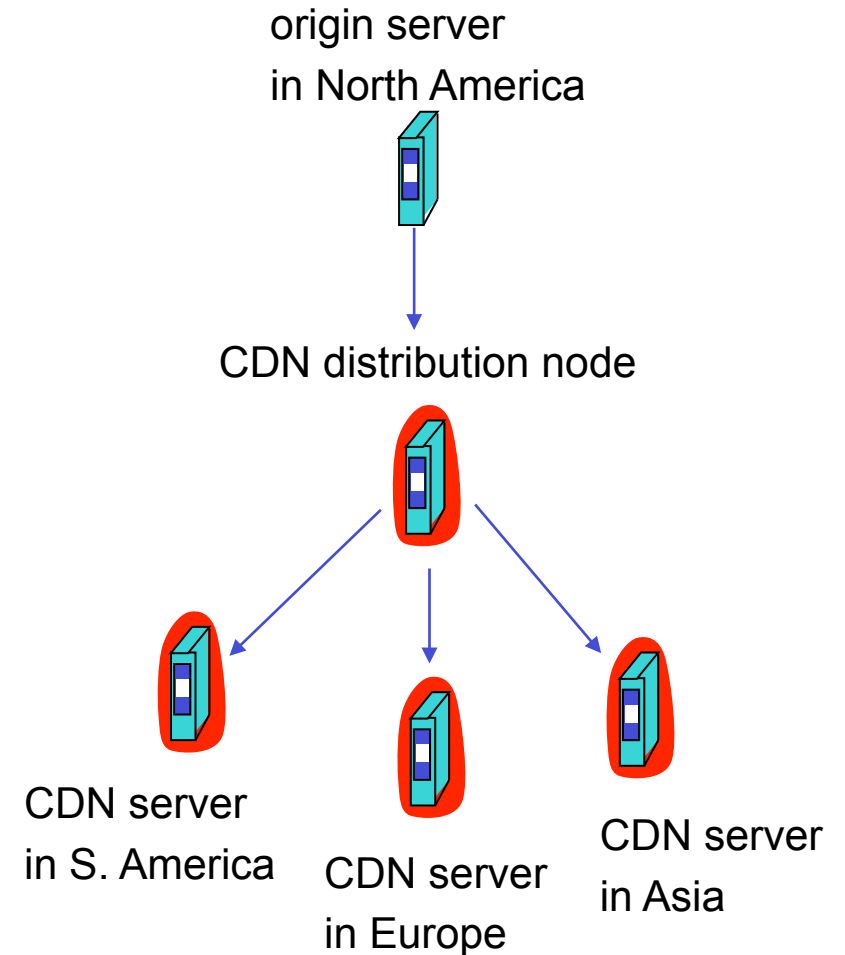
Overview:

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- ▶ 2.2 Streaming stored audio and video
- ▶ 2.3 Real-time Multimedia: Internet Phone study
- ▶ 2.4 Protocols for Real-Time Interactive Applications
 - RTP, RTCP
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- ▶ **2.6 Distributing Multimedia: content distribution networks**



Content replication

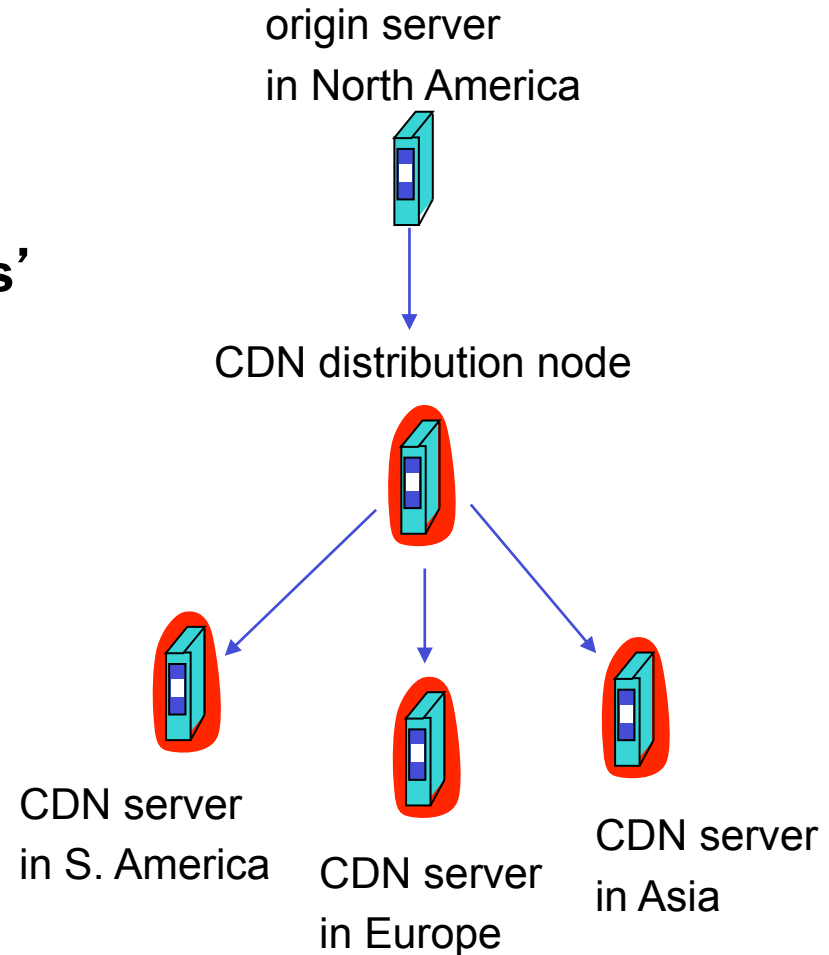
- **Challenging to stream large files (e.g., video) from single origin server in real time**
- **Solution: replicate content at hundreds of servers throughout Internet**
 - content downloaded to CDN servers ahead of time
 - placing content “close” to user avoids impairments (loss, delay) of sending content over long paths
 - CDN server typically in edge/access network





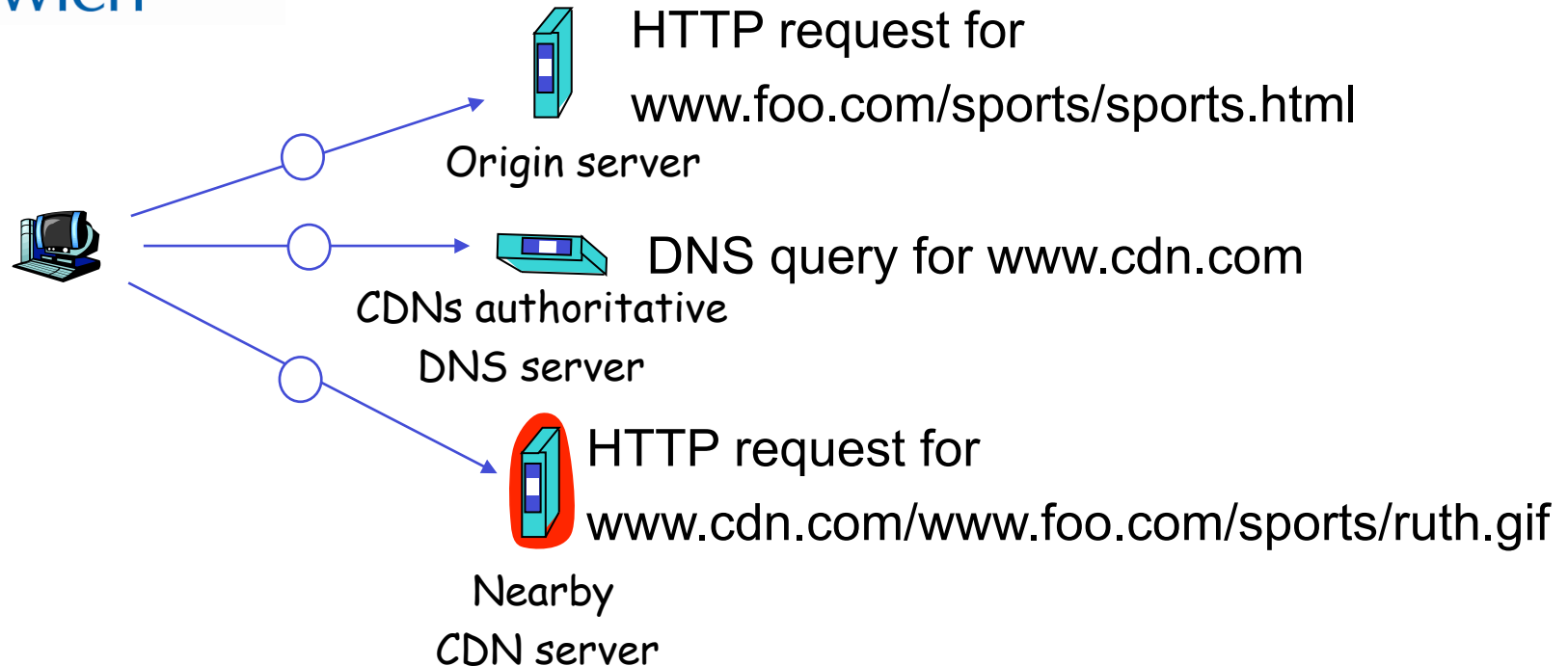
Content replication

- **CDN (e.g., Akamai) customer is the content provider (e.g., CNN)**
- **CDN replicates customers' content in CDN servers. When provider updates content, CDN updates servers**





universität wien CDN example



origin server (www.foo.com)

- distributes HTML
- replaces:
<http://www.foo.com/sports/ruth.gif>
with
<http://www.cdn.com/www.foo.com/sports/ruth.gif>

CDN company (cdn.com)

- distributes gif files
- uses its authoritative DNS server to route redirect requests



routing requests

- **CDN creates a “map”, indicating distances from leaf ISPs and CDN nodes**
- **when query arrives at authoritative DNS server:**
 - server determines ISP from which query originates
 - uses “map” to determine best CDN server
- **CDN nodes create application-layer overlay network**



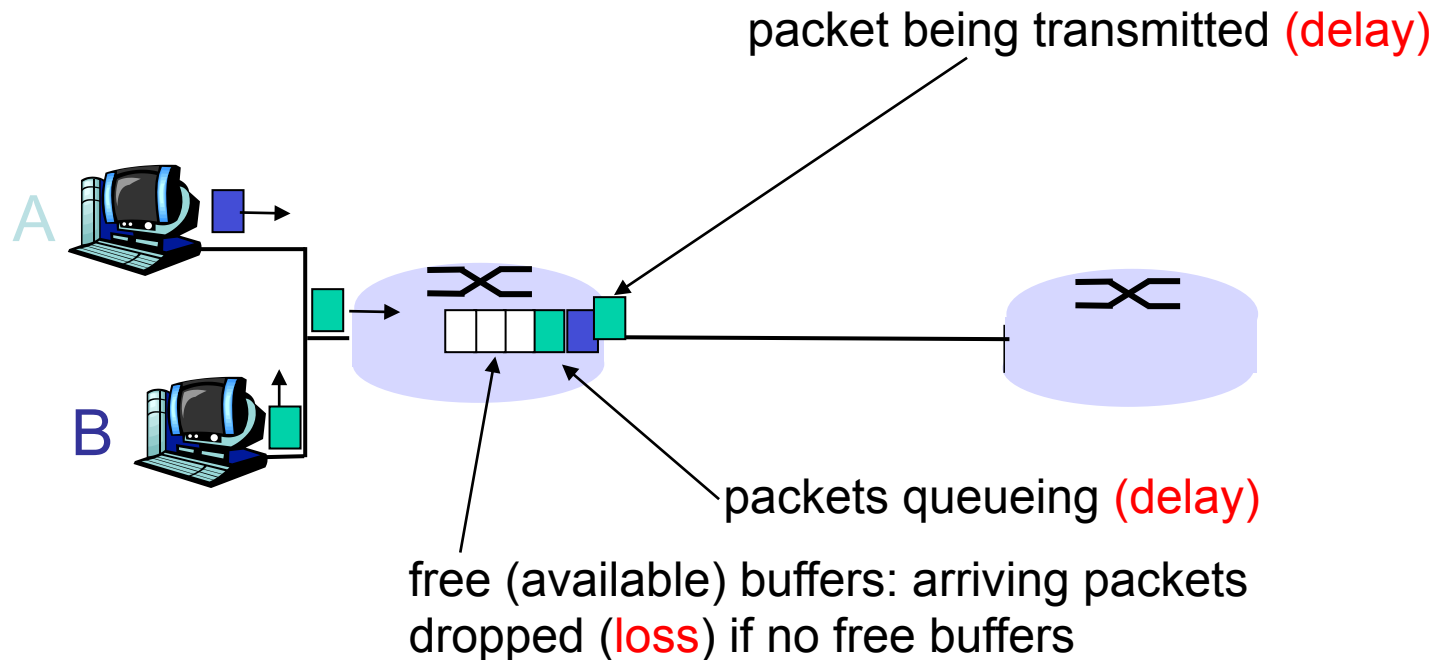
- Overview:
- **3.1 Packet loss & delay**
- 3.2 What's inside a router
-



How do loss and delay occur?

Packets *queue* in router buffers

- **packet arrival rate to link exceeds output link capacity**
- packets queue, wait for turn





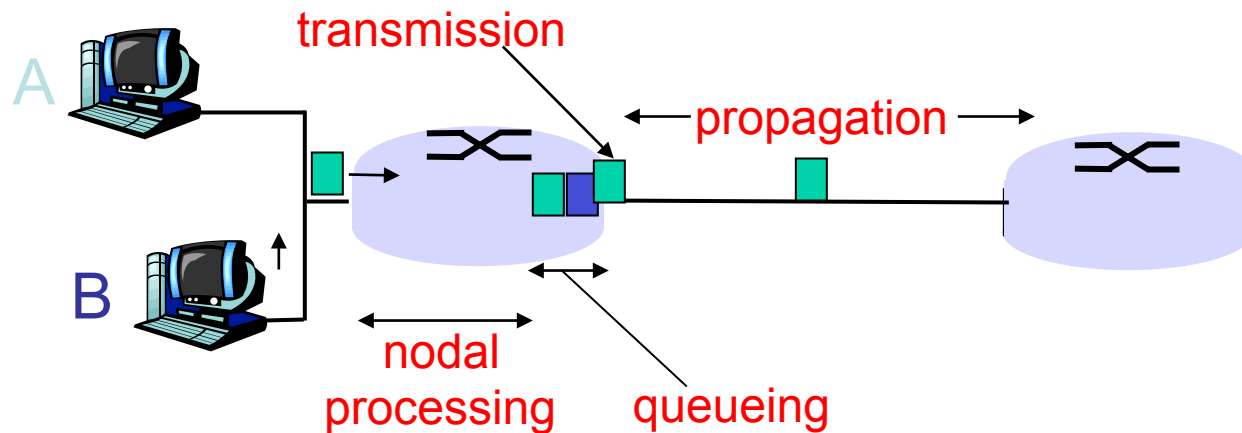
Four sources of packet delay

1. nodal processing

- check bit errors
- determine output link

2. queueing

- time waiting at output link for transmission
- depends on congestion level of router





Delay in packet-switched networks

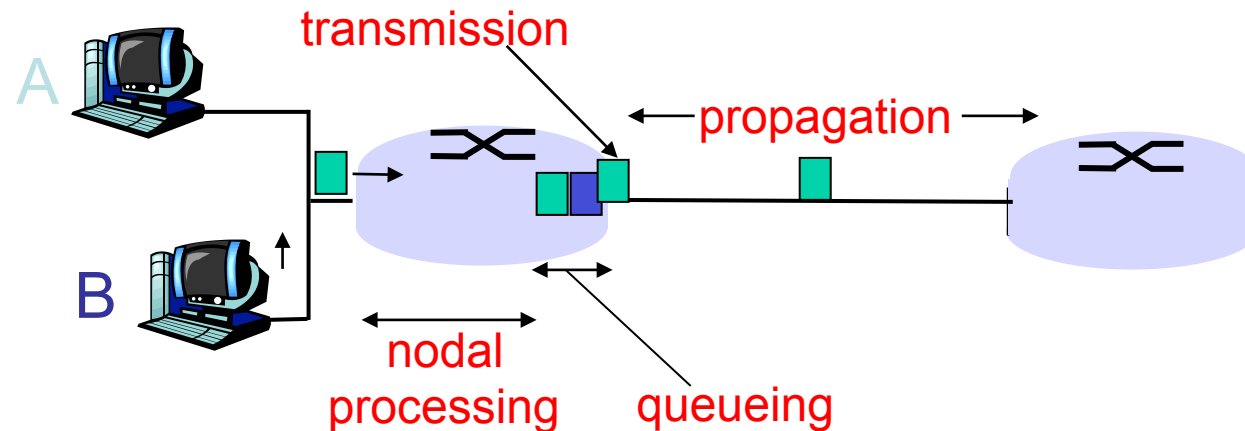
3. Transmission delay:

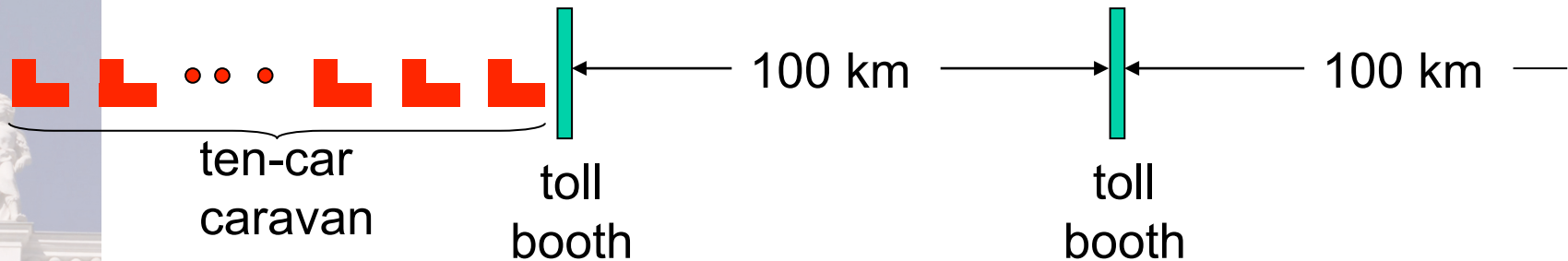
- R = link bandwidth (bps)
- L = packet length (bits)
- time to send bits into link = L/R

4. Propagation delay:

- d = length of physical link
- s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- propagation delay = d/s

Note: s and R are very different quantities!

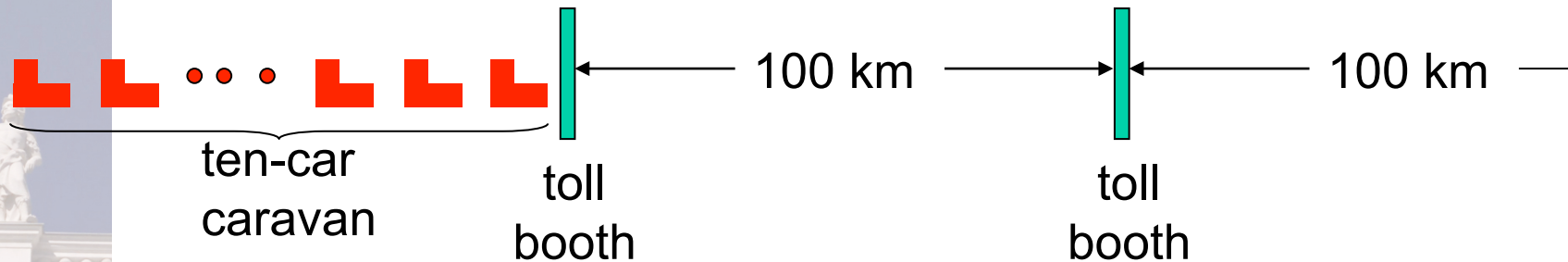




- Cars “propagate” at 100 km/hr
- Toll booth takes 12 sec to service a car (transmission time)
- car~bit; caravan ~ packet
- **Q: How long until caravan is lined up before 2nd toll booth?**
- Time to “push” entire caravan through toll booth onto highway = $12 \cdot 10 = 120$ sec
- Time for last car to propagate from 1st to 2nd toll both: $100\text{km}/(100\text{km/hr}) = 1$ hr
- **A: 62 minutes**



Caravan analogy (more)



- Cars now “propagate” at 1000 km/hr
- Toll booth now takes 1 min to service a car
- **Q: Will cars arrive to 2nd booth before all cars serviced at 1st booth?**
- **Yes!** After 7 min, 1st car at 2nd booth and 3 cars still at 1st booth.
- **1st bit of packet can arrive at 2nd router before packet is fully transmitted at 1st router!**

See Ethernet Applet at
AWL Web-Site



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

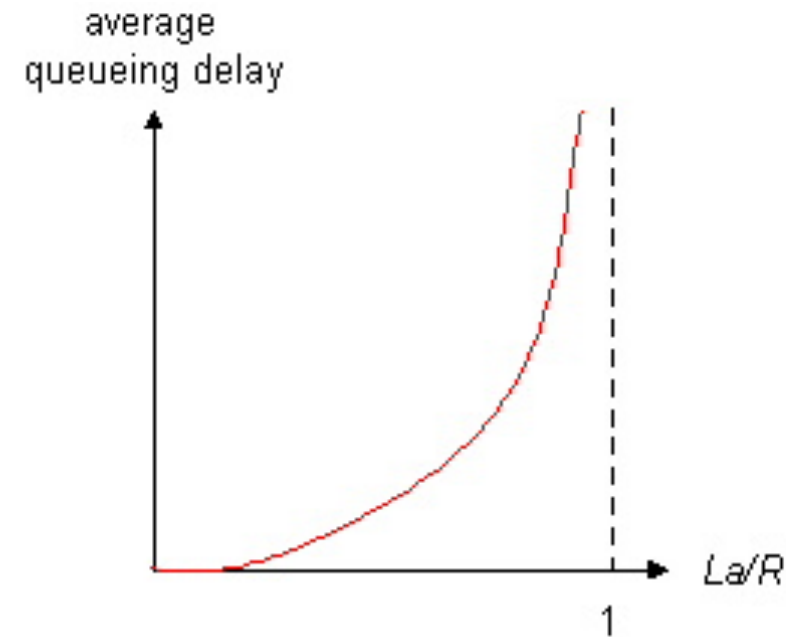
- **dproc = processing delay**
 - typically a few microsecs or less
- **dqueue = queuing delay**
 - depends on congestion
- **dtrans = transmission delay**
 - $= L/R$, significant for low-speed links
- **dprop = propagation delay**
 - a few microsecs to hundreds of msecs



Queueing delay (revisited)

- R =link bandwidth (b)
- L =packet length (bit)
- a =average packet arrival rate

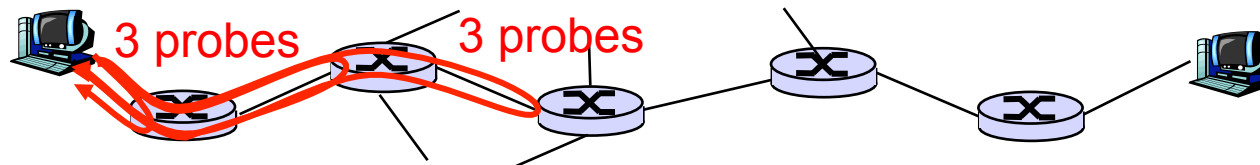
$$\text{traffic intensity} = \frac{La}{R}$$



- $La/R \sim 0$: average queueing delay small
- $La/R \rightarrow 1$: delays become large
- $La/R > 1$: more “work” arriving than can be serviced, average delay infinite!



- What do “real” Internet delay & loss look like?
- **Traceroute program:** provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.





universität “Real” Internet delays and routes wien

traceroute: gaia.cs.umass.edu to www.eurecom.fr

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 * * *
18 * * *
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

trans-oceanic link

*** means no reponse (probe lost, router not replying)**

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- Queue (aka buffer) preceding link in buffer has finite capacity
- When packet arrives to full queue, packet is dropped (aka lost)
- Lost packet may be retransmitted by previous node, by source end system, or not retransmitted at all



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Further QoS Metrics

- **Jitter**
- **Throughput**
- **Packet reordering**

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