

**Formal Methods in Computer Science**  
**Block 2: Satisfiability problems**  
**From Uninterpreted Functions to Equality Logic**  
**WS 2013**

---

**Ackermann's Reduction**

(on page 67 in D.Kroening, O. Strichman. Decision Procedures, Springer, 2008)

**Algorithm: ACKERMANN'S-REDUCTION**

**Input:** An EUF-formula  $\varphi^{EUF}$  with  $k$  uninterpreted functions  $F^{(1)}, \dots, F^{(k)}$  of arity  $> 0$ . For each  $F^{(l)}$  ( $1 \leq l \leq k$ ),  $m_l$  instances occur in  $\varphi^{EUF}$ .

**Output:** An equality logic formula  $\varphi^E$  such that  $\varphi^E$  is valid if and only if  $\varphi^{EUF}$  is valid.

1. Assign indices to the uninterpreted function instances from subexpressions outwards (separately for each function symbol). Denote by  $F_i^{(l)}$  the instance  $F^{(l)}$  that is given the index  $i$ . Denote by  $arg(F_i^{(l)})$  the argument(s) of  $F_i^{(l)}$ .
2. Let  $flat^E(\varphi^{EUF}) := \mathcal{T}(\varphi^{EUF})$ , where  $\mathcal{T}$  is a function that takes an EUF-formula (or term) as input and transforms it to an equality formula (or term respectively) by replacing each uninterpreted function instance  $F_i^{(l)}$  with a new term variable  $f_i^{(l)}$ . In the case of nested functions, only the variable corresponding to the most external instance remains. Constants and variables remain unchanged under  $\mathcal{T}$ .
3. Let  $FC_{F^{(l)}}^E(\varphi^{EUF})$  denote the following conjunction of functional consistency constraints for a function symbol  $F^{(l)}$ :

$$FC_{F^{(l)}}^E(\varphi^{EUF}) := \bigwedge_{i=1}^{m_l-1} \bigwedge_{j=i+1}^{m_l} (\mathcal{T}(arg(F_i^{(l)})) \doteq \mathcal{T}(arg(F_j^{(l)}))) \rightarrow f_i^{(l)} \doteq f_j^{(l)}.$$

If the arity of  $F^{(l)}$  is  $> 1$ , then read  $(s_1, \dots, s_n) \doteq (t_1, \dots, t_n)$  as  $\bigwedge_{p=1}^n (s_p \doteq t_p)$ .

4. Let  $FC^E(\varphi^{EUF})$  be  $\bigwedge_{l=1}^k FC_{F^{(l)}}^E(\varphi^{EUF})$ . Moreover, let

$$\varphi^E := FC^E(\varphi^{EUF}) \rightarrow flat^E(\varphi^{EUF}).$$

Return  $\varphi^E$ .

## Bryant's Reduction

(on page 70 in D.Kroening, O. Strichman. Decision Procedures, Springer, 2008)

### Algorithm: BRYANT'S-REDUCTION

**Input:** An EUF-formula  $\varphi^{EUF}$  with  $m$  instances of an uninterpreted function  $F$ .

**Output:** An equality logic formula  $\varphi^E$  such that  $\varphi^E$  is valid if and only if  $\varphi^{EUF}$  is valid.

1. Assign indices to the uninterpreted function instances from subexpressions outwards. Denote by  $F_i$  the instance  $F$  that is given the index  $i$ , and by  $arg(F_i)$  its argument(s).
2. Let  $flat^E(\varphi^{EUF}) := \mathcal{T}^*(\varphi^{EUF})$ , where  $\mathcal{T}^*$  is a function that takes an EUF-formula (or term) as input and transforms it to an equality formula (or term respectively) by replacing each uninterpreted function instance  $F_i$  with a new term variable  $F_i^*$  (in the case of nested functions, only the variable corresponding to the most external instance remains). Constants and variables remain unchanged under  $\mathcal{T}^*$ .
3. For  $i \in \{1, \dots, m\}$ , let  $f_i$  be a new variable. Let

$$F_i^* \doteq \left( \begin{array}{ccc} case & \mathcal{T}^*(arg(F_1)) & \doteq & \mathcal{T}^*(arg(F_i)) & : & f_1 \\ & \vdots & & \vdots & & \vdots \\ & \mathcal{T}^*(arg(F_{i-1})) & \doteq & \mathcal{T}^*(arg(F_i)) & : & f_{i-1} \\ & TRUE & & & : & f_i \end{array} \right)$$

be a case statement (as defined in the lecture slides) and let

$$C(F_i^*) := \bigvee_{j=1}^i (F_i^* \doteq f_j \wedge \mathcal{T}^*(arg(F_j)) \doteq \mathcal{T}^*(arg(F_i)) \wedge \bigwedge_{k=1}^{j-1} \mathcal{T}^*(arg(F_k)) \neq \mathcal{T}^*(arg(F_i)))$$

be the associated formula.

Finally, let

$$FC^E(\varphi^{EUF}) := \bigwedge_{i=1}^m C(F_i^*).$$

4. Let

$$\varphi^E := FC^E(\varphi^{EUF}) \rightarrow flat^E(\varphi^{EUF}).$$

Return  $\varphi^E$ .

Remarks:

- The generalization to more function symbols and to function symbols with arity  $> 1$  works as in Ackermann's reduction.
- In the aforementioned book, there are two small errors, which we corrected:
  1. In the case statement we replaced the wrong statements of the form  $arg(F_j^*)$  by the correct ones of the form  $arg(F_j)$ .
  2. In the definition of  $FC^E(\varphi^{EUF})$  we replaced the wrong statement of the form  $F_i^*$  by the correct one of the form  $C(F_i^*)$ .