

VU Programm- und Systemverifikation

Assignment 1: Assertions, Testing, and Coverage

Name: _____ Matr. number: _____

Due: April 22, 2021, 1pm

1 Coverage Metrics

Consider the following program fragment and test suite:

```
unsigned gcd (unsigned x, unsigned y) {  
    unsigned m, k;  
    if (x > y) {  
        k = x;  
        m = y;  
    } else {  
        k = y;  
        m = x;  
    }  
    while (m != 0) {  
        unsigned r = k % m;  
        k = m;  
        m = r;  
    }  
    return k;  
}
```

Inputs		Outputs
x	y	return value
0	0	0
0	1	1
3	2	1

1.1 Control-Flow-Based Coverage Criteria (2.5 points)

Indicate (✓) which of the following coverage criteria are satisfied by the test-suite above.

	satisfied	
Criterion	yes	no
path coverage		✓
statement coverage	✓	
branch coverage	✓	
decision coverage	✓	

For each coverage criterion that is *not* satisfied, explain why this is the case:

In the presence of a loop, there are arbitrarily many paths; hence the path coverage criterion cannot be achieved.

1.2 Data-Flow-Based Coverage Criteria (4 points)

Indicate (✓) which of the following coverage criteria are satisfied by the test-suite above (here, the parameters of the function do not constitute definitions, the `return` statement is a c-use):

Criterion	satisfied	
	yes	no
all-defs	✓	
all-c-uses		✓
all-p-uses	✓	
all-c-uses/some-p-uses		✓
all-du-paths		✓

For each coverage criterion that is not satisfied, explain why this is the case:

- *The loop is only entered in the third test-case (in which $x > y$ holds). Consequently the use $k \% m$ of the definitions $k = y$ and $m = x$ in the else-branch of the conditional statement is never reached, and the definitions $k = y$ and $m = x$ in the then-branch never reach the return statement.*
- *Since all-c-uses is not achieved, all-c-uses/some-p-uses cannot be achieved either.*
- *Since all-c-uses is not achieved, there are du-paths that have not been covered.*

1.3 Achieving Full Coverage (1 point)

Consider the two coverage criteria below.

- If the test-suite from above does not satisfy the coverage criterion, augment it with the *minimal* number of test-cases such that this criterion is satisfied. If full coverage cannot be achieved, explain why.
- If the coverage criterion is already achieved, explain why.

all-c-uses

Inputs		Outputs
x	y	result
2	3	1
1	0	1

all-uses

Inputs		Outputs
x	y	result
2	3	1
1	0	1

1.4 General Questions (2.5 points)

Indicate whether the following statements are true or false!

Statement	True	False
If a test-suite achieves all-c-uses and all-defs coverage, it also achieves all-c-uses/some-p-uses coverage.	<input checked="" type="radio"/>	<input type="radio"/>
If a test-suite achieves path coverage, it also achieves statement coverage.	<input checked="" type="radio"/>	<input type="radio"/>
If a test-suite does not achieve all-uses coverage, it cannot achieve all-c-uses either.	<input type="radio"/>	<input checked="" type="radio"/>
For any given program, it is always possible to achieve MC/DC coverage.	<input type="radio"/>	<input checked="" type="radio"/>
If full statement coverage cannot be achieved, then the program contains unreachable code.	<input checked="" type="radio"/>	<input type="radio"/>

1.5 Formal Definition (1 point)

Let $S_{\circlearrowleft}, S_1, \dots, S_n$ be the sets of reachable states at program locations $\circlearrowleft, 1, \dots, n \in V$, respectively. Moreover, given the conditional statement at program location ℓ , let C_ℓ be corresponding decision. Assume that all branches correspond to conditional statements.

Define an abstraction function α and an abstract domain such that

$$\forall i \in v. \alpha(S_i) = \top$$

if (and only if) full branch coverage can be reached.

$$\alpha(S_i) \stackrel{\text{def}}{=} \begin{cases} \{\sigma(B_i) \mid \sigma \in S_i\} & \text{if } i \text{ is branching statement with condition } B_i \\ \top & \text{otherwise} \end{cases}$$

(We assume there are no unconditional branches.)

2 Equivalence Partitioning and Boundary Testing

If you who had contact with a person who has tested positive for COVID-19, the function `classify` below will determine whether you had a category I contact, a category II contact, or neither.

```
category classify (int time,  
                 int distance,  
                 int their_age,  
                 int your_age);
```

It uses the following data-types and parameters:

- `category` is an enum type defined as `enum priority {C1, C2, NEITHER};`
- `time` is the (cumulative) amount of time (in minutes) of exposure.
- `distance` is the minimal distance (in meters) to the infected person.
- `their_age` is the age of the person who tested positive for COVID-19.
- `your_age` is the age of the person that is to be classified.

The function `classify` is supposed to implement the following rules:

- A person who (cumulatively) spent at least 15 minutes in less than 2 meters distance from a person older than 10 years of age is classified as category I contact.
- Children younger than 10 years who were in contact (for at least 15 minutes in less than 2 meters distance) with a child also younger than 10 years that tested positive for COVID-19 are classified as category II contacts.
- A person who (cumulatively) spent less than 15 minutes with an infected person or kept a distance of at least 2 meters does not fall into category I or II.

Note: *The requirements are unclear about positively tested contact persons aged exactly 10. Therefore, we will have separate equivalence classes with `age=10`.*

2.1 Equivalence Partitioning (2 points)

From the specification above, derive equivalence classes for the function `classify`. Use the table below to partition them into *valid equivalence classes* (valid inputs) and *invalid equivalence classes* (invalid inputs). Label each of the equivalence classes clearly with a number (in the according column). For each correct equivalence class you can score $\frac{1}{2}$ a point (up to 2 points).

(Do not provide test-cases here – that’s task 2.2)

2.1.1 Valid Equivalence Classes

Condition	ID
<code>(time ≥ 15) ∧ (distance < 2) ∧ (their_age > 10)</code>	1
<code>(time ≥ 15) ∧ (distance < 2) ∧ (your_age < 10) ∧ (their_age < 10)</code>	2
<code>(time ≥ 15) ∧ (distance < 2) ∧ (your_age < 10) ∧ (their_age = 10)</code>	3
<code>(time < 15) ∨ (distance ≥ 2)</code>	4

Note that we didn’t choose to create separate equivalence classes for `(time ≥ 15)`, `(distance < 2)`, ...; the reason is, that the outcome depends on a combination of the parameters.

2.1.2 Invalid Equivalence Classes

Condition	ID
<code>(time < 0)</code>	5
<code>(distance < 0)</code>	6
<code>(your_age < 0)</code>	7
<code>(their_age < 0)</code>	8

2.2 Boundary Value Testing (2 points)

Use *Boundary Value Testing* to derive a test-suite for the function `triage`. Specify the inputs points for `classify`. Indicate clearly which equivalence classes each test-case covers by referring to the numbers from task (a). You can receive up to 2 points ($\frac{1}{2}$ a point per test-case), where redundant test-cases and test-cases that do not represent boundary values do not count.

Input	Output	Classes Covered
(15, 1, 18, 11)	C1	1
(15, 0, 0, 11)	C1	1
(15, 1, 9, 9)	C2	2
(15, 0, 9, 0)	C2	2
(15, 1, 9, 10)	C2 (?)	3
(14, 0, 18, 42)	NEITHER	4
(15, 2, 18, 42)	NEITHER	4
(-1, 0, 18, 42)	error	5
(15, -1, 18, 42)	error	6
(15, 2, -1, 42)	error	7
(15, 2, 18, -1)	error	7
...		

3 Invariants (5 points)

Consider the following program, where t , x , y , and z are integer values in \mathbb{Z} (that means no over- or underflow can happen):

```
z = 1;
if (x > y) {
    int t = x; x = y; y = t;
}
while (y > x) {
    x = x + 1;
    y = y - 1;
    z = z + (y - x);
}
```

Consider the formulas below; tick the correct box () to indicate whether they are loop invariants for the program above.

- If the formula is an inductive invariant for the loop, provide an informal argument that the invariant is inductive.
- If the formula P is an invariant that is *not* inductive, give values of x , y , and z before and after the loop body demonstrating that the Hoare triple

$$\{P \wedge B\} \quad x = x + 1; y = y - 1; z = z + (y - x) \quad \{P\}$$

(where B is $(y > x)$) does not hold. This means that you need to find a starting state in which $P \wedge B$ evaluates to true and which results in a violation of P after executing the loop body.

- Otherwise, provide values of x , y , and z that correspond to a reachable state showing that the formula is *not* an invariant.

$(x \leq y)$	<input type="checkbox"/> Inductive Invariant	<input type="checkbox"/> Non-inductive Inv.	<input checked="" type="checkbox"/> Neither
Justification: Before: $x = 0, y = 1, z$ arbitrary After: $x = 1, y = 0$			
$(x - y) \leq 1$	<input checked="" type="checkbox"/> Inductive Invariant	<input type="checkbox"/> Non-inductive Inv.	<input type="checkbox"/> Neither
Justification: Base: $x \leq y$ implies $(x - y) \leq 1$ Induction: Pre-condition of $x - y \leq 1$ is $x - y < 0$, implied by $y > x$			
$(y - x) + 2 \neq 0$	<input checked="" type="checkbox"/> Inductive Invariant	<input type="checkbox"/> Non-inductive Inv.	<input type="checkbox"/> Neither
Justification: Note: This invariant is also implied by $(x - y) \leq 1$. Base: $x \leq y$ implies that $y - x$ is positive, hence adding 2 can't result in 0 Induction: if $y > x$, then $y - x$ is at least one, and after the loop at least -1. Hence adding 2 can't make the term 0.			
$z \geq 0$	<input type="checkbox"/> Inductive Invariant	<input checked="" type="checkbox"/> Non-inductive Inv.	<input type="checkbox"/> Neither
Justification: Before: $z = 0, y = 1, x = 0$ After: $x = 1, y = 0, z = 0 + (0 - 1)$ We know that $-1 \leq (y - x)$ when z is assigned, and $(y - x) = -1$ only if $y = x + 1$ upon loop entry. So after z is decreased once by 1, the loop terminates. Note that z is initially 1.			
$z \geq 1$	<input type="checkbox"/> Inductive Invariant	<input type="checkbox"/> Non-inductive Inv.	<input checked="" type="checkbox"/> Neither
Justification: Start with $z = 1, x = 0, y = 1$, the $z = 0$ after the loop.			