

Distributed Algorithms 182.702

Quiz 2 (Ch.3 + Ch.2, SS 2021), Form: A

Name: _____

Program code, registr. no.: _____

Date: _____

Achieved points: _____

Question 1. Tick the correct statements:

1. (a) The end-to-end delay τ of a message m is defined as the time between $comp(i)$ where m is sent and $del(i, j, m)$.
- (b) Suppose that some problem P has a worst-case lower bound of $\Omega(n \log n)$ messages for all correct algorithms solving P . An algorithm \mathcal{A} for P that has some execution where only n messages are sent cannot be correct.
2. (a) Consider two schedules σ and σ' and the configurations $C = \sigma(C^0)$ and $C' = \sigma'(C^0)$ reached when starting from C^0 in the distributed computing model of textbook. If $\sigma = del(i, j, m_1), del(k, j, m_2), comp(j)$ and $\sigma' = del(i, j, m_1), comp(j), del(k, j, m_2), comp(j)$, then $C = C'$.
- (b) If one drops $inbuf_i[*]' = \emptyset$ after a step, then an event $comp(i)$ cannot uniquely determine the corresponding step of p_i .
- (c) A norm function, bound to some assertion A , is defined upon an arbitrary partial order $(W, <)$
- (d) An assertion is a unary predicate on the transition relation.
3. (a) Any deterministic leader election algorithm requires processors with unique identifiers.
- (b) There is no asynchronous leader election algorithm in rings with message complexity at most $O(n^2)$.
- (c) A time-bounded leader election algorithm has a time complexity upper bound that is independent of the id's.
4. (a) A liveness property is a set of event traces such that any finite sequence of events has some extension that is contained in this set.
- (b) Fairness is a safety property.
- (c) An asynchronous execution of a message passing system, where every processor takes an infinite number of steps, is always admissible.

Question 2.

1. What is the asymptotic message complexity of a flooding algorithm like the one in Algorithm 2?

2. A ring where messages travel in one direction only is called _____
3. If P is a safety property and $\beta \in P$ with $\beta = \phi^1 \phi^2 \phi^3 \dots$, what can be said about the finite trace $\phi^2 \phi^3$?

4. Give the definition of a deterministic transition relation Φ_i for processor p_i :

Question 3.


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2  VAR parent = ... // initially specifies parent in spanning tree, or NULL at p0
3     start = ... // initially false, becomes true once during the execution
4     next := ⊥

5  // When start becomes true (happens only on pi ≠ p0)
6     send message M(i) to parent
7     parent := NULL
8     start := false

9  // On reception of M(k) received from pj
10     if parent = NULL then
11         next := k
12     else
13         send message M(k) to parent
14     endif
15     parent := j // In any case: adjust parent to pj
16

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Prove the invariant $A(C) = A1(C) \wedge A2(C)$ defined by

- (A1) For every processor p_i , $parent_i(C)$ is either NULL or else one of the processors p_j or p_ℓ connected by a link (p_j, p_ℓ) in the original spanning tree.
- (A2) Messages can be in transit only on a link (p_j, p_ℓ) in the original spanning tree.