

Theorem $\chi(G, x)$ is a polynomial.

Proof: Induction on $|E| + |V|$.

case $|E| = 0 \Rightarrow \chi(G, x) = x$ is a polynomial

case $|E| > 0$: $\chi(G, x) = \chi(G \setminus e, x) - \chi(G/e, x)$
 $\stackrel{!}{=} \text{polynomial} - \text{polynomial}$

□

② combinatorics

Given a finite set A_n for each $n \in \mathbb{N}$, what is $|A_n|^2$? e.g. $A_n = \{\pi \mid \pi \text{ permutation of } \{1, \dots, n\}\}$

An answer might be:

-) an "explicit formula": $|A_n| = n!$
-) a recurrence: $|A_n| = n|A_{n-1}|$, $|A_0| = 1$
-) an asymptotic formula $|A_n| \sim n^n e^{-n} \sqrt{2\pi n}$
-) an algorithm

counting principles

Let A, B, C be finite sets

$$\Rightarrow A \cap B = \emptyset \Rightarrow |A \cup B| = |A| + |B|$$

$$|A \times B| = |A| \cdot |B|$$

$$f: |A| \rightarrow |B| \text{ bijective} \Rightarrow |A| = |B|$$

Double counting: Let $A = \{a_1, \dots, a_n\}$, $B = \{b_1, \dots, b_m\}$, $R \subseteq A \times B$, $|R|^2$:

Let $R_{i,j} := \{(a_i, b_j) \in R\}$, $R_{*,j} = \{(a, b_j) \in R \mid a \in A\}$

$$\Rightarrow |R| = \sum_{j=1}^m R_{*,j} = \sum_{i=1}^n R_{i,*}$$

E.g.: $\bar{c}(n) = \text{average number of divisors of integers between 1 and } n$

$$\bar{c}(6): \begin{array}{c} \begin{array}{cccccc} 1 & 2 & 3 & 4 & 5 & 6 \\ \boxed{1} & 2 & 3 & 4 & 5 & 6 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{array} \\ \text{1 divides everything,} \\ \text{2 divides even numbers,} \\ \text{3 divides multiples of 3,} \\ \text{4 divides multiples of 4,} \\ \text{5 divides multiples of 5,} \\ \text{6 divides multiples of 6} \end{array} \Rightarrow \bar{c}(6) = (1+2+2+3+2+4) \frac{1}{6} = \frac{19}{6} = \frac{1}{6} (6+3+2+1+1+1) = \frac{1}{n} \sum_{i=1}^n \frac{1}{i}$$

↓
sum over ~~rows~~ columns
↓
sum over rows

$$\text{we only want to approximate: } \frac{1}{n} \sum_{i=1}^n \lfloor \frac{n}{i} \rfloor = \frac{1}{n} \sum_{i=1}^n \frac{n}{i} + O(1) = \sum_{i=1}^n \frac{1}{i} + O(1) \sim \log(n)$$

constant error

Pigeonhole - Principle (Schubfachprinzip)

A_1, \dots, A_k , pairwise disjoint, $|A_1 \cup \dots \cup A_k| > k \cdot r$

$$\Rightarrow \exists i: |A_i| > r$$

Alternatively: $f: |A| \rightarrow |B|$, $|A| > |B| \Rightarrow \exists b \in B: |f^{-1}(b)| \geq 2$ (not injective).

E.g. $\forall g \text{ odd } \exists i : g | 2^i - 1$, $a_i = 2^i - 1 = 1, 3, 7, 15, 31, 63, 127, \dots$

Proof: ~~if a_1, \dots, a_g are all different mod g~~ Consider $a_1, \dots, a_g \text{ mod } g$

If $\exists i : a_i \equiv 0 \pmod{g}$, we are done, otherwise $\exists i < j : a_i \equiv a_j \pmod{g}$

(a_1, \dots, a_g are ~~g~~ numbers, but can only have $g-1$ remainders mod g different from 0)

$\Rightarrow a_i - a_j = g \cdot a$ for some a

$$= 2^i (1 - 2^{j-i}) \quad , \text{ then}$$

$$g \text{ odd} \Rightarrow g | 1 - 2^{j-i} = -a_{j-i}$$

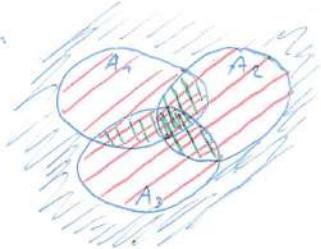
□

Principle of Inclusion/Exclusion (PIE)

Let E_1, \dots, E_n be properties, let $A_i = \{x \in A \mid x \text{ has } E_i\}$

$$\Rightarrow |A \setminus \bigcup_{i=1}^n A_i| = \left(\sum_{\substack{I \neq \emptyset \\ I \subseteq \{1, \dots, n\}}} (-1)^{|I|} \cdot |\bigcap_{i \in I} A_i| \right) + |A|$$

E.g. $n=3$:



$$\begin{aligned} |A| &= (\underline{|A_1| + |A_2| + |A_3|}) \\ &\quad + (\underline{|A_1 \cap A_2| + |A_2 \cap A_3| + |A_1 \cap A_3|}) \\ &\quad - (\underline{|A_1 \cap A_2 \cap A_3|}) \end{aligned}$$

Sketch of proof: $|\bigcup_{i=1}^m A_i| = - \sum_{\substack{\emptyset \neq I \\ I \subseteq \{1, \dots, m\}}} (-1)^{|I|} |\bigcap_{i \in I} A_i|$, consider $a \in \bigcup_{i=1}^m A_i$. Let $S = \{i \mid a \in A_i\}$

The element a contributes exactly to summands $I \subseteq S$ on the right side.

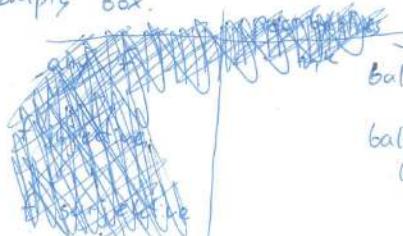
$$\begin{aligned} - \sum_{\substack{\emptyset \neq I \\ I \subseteq S}} (-1)^{|I|} &= - \sum_{k=1}^{|S|} \sum_{\substack{I \subseteq k \\ I \subseteq S}} (-1)^k = - \sum_{k=1}^{|S|} (-1)^k \left(\sum_{I \subseteq S} 1 \right) = - \sum_{k=1}^{|S|} (-1)^k \binom{|S|}{k} = - \left(\sum_{k=0}^{|S|} (-1)^k \binom{|S|}{k} - 1 \right) \\ &= -((1-1)^{|S|} - 1) = 1 \end{aligned}$$

□

Balls into boxes (k balls, n boxes), represented by $f: [k] \rightarrow [n]$.

Notation: $[n] := \{1, \dots, n\}$

If f is injective, this means balls do not share boxes, if f is surjective, there's no empty box.



Balls and boxes labelled

balls unlabelled,
boxes labelled

	any f	f injective	f surjective
Balls and boxes labelled	n^k	$n(n-1)\dots(n-k+1)$	not nice
balls unlabelled, boxes labelled	$\binom{n+k-1}{k}$	$\binom{n}{k}$	

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{n(n-1)\cdots(n-k+1)}{k!}$$

with this definition
 $n \in \mathbb{C}$ is possible

because $k! \binom{n}{k} = \underbrace{n(n-1)\cdots(n-k+1)}_{\text{number of possibilities}} \leftarrow \begin{matrix} \text{to label } k \text{ balls} \\ \text{with } 1, \dots, n \end{matrix}$
 $\binom{n}{k} = \underbrace{\text{plain balls in } n \text{ labelled boxes}}_{\text{in labelled boxes}}$

Remark: There's no division in combinatorics

Any f , balls ~~are~~ unlabelled, boxes labelled:

$$\begin{array}{ccccccc} 1 & 2 & 3 & 4 & 5 \\ 001 & 01001 & 10001 & 100001 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ n=5 & & & & k=9 \end{array}$$

as every bit carries information: 00101 $\rightsquigarrow n=3=k$

$$\# \text{ bits} = n+k-1$$

$$\# \text{ zeros} = k$$

$$\Rightarrow \binom{n+k-1}{k}$$

[if surjective, balls unlabelled, boxes labelled is skipped.]

some binomial identities:

$$*) (x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k} \quad \text{means } \text{unlabelled } n \text{ balls and 2 boxes ("x" and "y")}$$

E.g.: $2xy^3$: 2 possibilities for 1 ball in box x, 3 balls in box y

$\binom{n}{k} x^k y^{n-k} = \binom{n}{k}$ possibilities to place k balls into the x-box, $n-k$ balls into the y-box.

$$*) \sum_{m=0}^n \binom{m}{k} = \binom{n+1}{k+1}$$

$$*) \sum_{k=0}^n \binom{m+k}{k} = \binom{m+n+1}{n}$$

$$*) \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}, \text{ also works for } n \in \mathbb{C} \quad (\text{Lemma})$$

proof: easy for $n \in \mathbb{N}$ (exercise)

the left-hand side is a polynomial in n (leading coefficient: $\frac{1}{k!}$), $p(n)$, ($\deg(p(n)) = k$)

right-hand side is a polynomial in n , $q(n)$, ($\deg(q(n)) = \max(k-1, k) = k$)

$\Rightarrow p(x) = q(x) \quad \forall x \in \mathbb{C}$ because $p(n) = q(n) \quad \forall n \in \mathbb{N}$

Theorem (Vandermonde) $\binom{x+y}{n} = \sum_{k=0}^n \binom{x}{k} \binom{y}{n-k}$

Proof: $x, y \in \mathbb{N}$, let $|X|=x$, $|Y|=y$, $X \cap Y = \emptyset$

$\binom{x+y}{n}$: # subsets of $X \cup Y$ of size n

$\binom{X}{k} \binom{Y}{n-k}$: # subsets of $X \cup Y$ with k elements in X , the rest in Y .

□

Sterling numbers:

Every permutation of $[n]$ is a product of cycles.

Eg: $(10\ 11)(1\ 3\ 8\ 5\ 2)(6\ 4\ 9)(7)$
(elementary)
transposition
cycle
fixed point

product: $(1\ 3\ 2)(2\ 3\ 4) = (1\ 3\ 4)(2) \rightarrow \begin{matrix} 2 \rightarrow 3 \\ 3 \rightarrow 2 \end{matrix}$
 $\downarrow \quad \downarrow \quad \downarrow$
 $1 \rightarrow 1 \quad 3 \rightarrow 4 \quad 4 \rightarrow 2$
 $1 \rightarrow 3 \quad 4 \rightarrow 9 \quad 2 \rightarrow 1$

Definition: $s_{n,k}$ is the number of permutations in S_n (permutations of $[n]$) with k cycles, the stirling number of the first kind.

E.g.: $s_{n,n} = (n-1)!$, $s_{3,1} = |\{(23), (132)\}|$

$$s_{n,n-1} = \binom{n}{2}$$

$$s_{n,n} = 1$$

$$s_{0,0} := 0, s_{n,0} = s_{0,n} = 0 \text{ for } n > 0$$

Theorem $s_{n,k} = s_{n-1,k-1} + (n-1)s_{n-1,k}$

Proof $(1 \dots)(\dots) \dots (\dots)$

case 1 fixed point: $s_{n-1,k-1}$ possibilities

otherwise: $s_{n-1,k}$ has k cycles, but element one is missing $\Rightarrow n-1$ possibilities to put it without changing the number of cycles.

□

Definition A set partition of A is a set of disjoint, non-empty sets with union A .

$S_{n,k} :=$ # set partitions of $[n]$ with k parts ~~disjoint~~ (or "blocks")

$S_{n,k}$ is called the Sterling number of the second kind. ($s_{0,0} := 1$)

$$S_{0,0} := 1, S_{n,0} = S_{0,n} = 0, n > 0$$

Theorem $S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$

Proof: case $\{n\}$ is a singleton, then $S_{n-1,k-1}$ if removed

otherwise: put n into one of the k blocks: $kS_{n-1,k}$

□

Theorem: $(x)_n := x^n := x(x-1)\cdots(x-n+1) = \sum_{k=0}^n (-1)^{n-k} S_{n,k} x^k$,

$$x^n = \sum_{k=0}^n S_{n,k} x^k$$

Remark $V_n = \{a_0 + \dots + a_n x^n \mid a \in \mathbb{C}\}$ is a vector space, $\{1, x, \dots, x^n\}$ natural basis.

$\{1, x, x^2, \dots, x^n\}$ is also a basis of V_n . The change of basis matrices are $(S_{n,k})_{n,k}$ and $((-1)^{n-k} s_{n,k})_{n,k}$ (exercises).

Proof of the theorem: Induction on n : $x^0 = 1 = S_{0,0} x^0$ ✓

$$\begin{aligned} x^n &= x^{n-1}(x-n+1) = (x-n+1) \sum_{k=0}^{n-1} (-1)^{n-1-k} S_{n-1,k} x^k = \sum_k (-1)^{n-1-k} S_{n-1,k} x^{k+1} + (n-1) \sum_k (-1)^{n-1-k} S_{n-1,k} x^k \\ &= \sum_k (-1)^{n-k} S_{n-1,k-1} x^k + (n-1) \sum_k (-1)^{n-1-k} S_{n-1,k} x^k \end{aligned}$$

□

Generating functions

Power series: a sequence $(a_n)_{n \in \mathbb{N}_0}$ (one for example). Consider $\sum_{n \geq 0} a_n z^n$ is the series with coefficients a_n . $\sum a_n z^n$ may or may not converge for given complex number $z \in \mathbb{C}$.

A formal power series, written $\sum a_n z^n$ is the same information as the sequence $(a_n)_{n \in \mathbb{N}}$.

~~Formal power series~~

$$\sum_{n=0}^{\infty} a_n z^n \stackrel{\text{power series}}{=} \lim_{N \rightarrow \infty} \sum_{n=0}^N a_n z^n, \quad \text{i.e. as a limit of a sequence of complex numbers.}$$

Theorem: $\lim_{N \rightarrow \infty} \sum_{n=0}^N a_n z^n$ exists if $|z| < \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|a_n|}} := R$, \limsup being the limit superior.

If $|z| = \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|a_n|}}$ an ad hoc analysis is necessary, if $|z| > \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|a_n|}}$ it diverges.

Remark: $\{z \mid \sum a_n z^n \text{ converges}\}$ is the domain of convergence, essentially a circle centered at the origin.

E.g.: $\sum_{n=0}^{\infty} z^n = \frac{1}{1-z}$ geometric series, $R=1$.

$\sum_{n=0}^{\infty} \frac{z^n}{n!} = e^z$ exponential series, $R=\infty$ (converges everywhere)

$$\sum_{n=0}^{\infty} \binom{\alpha}{n} z^n = (1+z)^\alpha, \quad \alpha \in \mathbb{C}$$

Theorem (Identity theorem for power series)

$f(z) = \sum_{n=0}^{\infty} a_n z^n$ converges for $|z| < R$ and $R > 0$.

$$\text{Then } a_n = \frac{f^{(n)}(0)}{n!}$$

Corollary: $f(z) = \sum a_n z^n = \sum b_n z^n \quad \text{if } |z| < R \Rightarrow a_n = b_n \forall n$

Operations on formal power series

$A(z) := \sum_{n \geq 0} a_n z^n, B(z) := \sum_{n \geq 0} b_n z^n$ (however, z is not a complex number now!)

Write $(a_n) \leftrightarrow A(z), (b_n) \leftrightarrow B(z), (0, 1, 0, 0, \dots) \leftrightarrow z, (1, 0, 0, \dots) \leftrightarrow 1, \dots$

- $(\alpha a_n + M b_n)_{n \in \mathbb{N}_0} \leftrightarrow: \alpha A(z) + M B(z)$
- $\left(\sum_{k=0}^n a_k b_{n-k} \right)_{n \in \mathbb{N}_0} \leftrightarrow: A(z) B(z) \quad (\text{Convolution})$
- $(a_n \delta^n)_{n \in \mathbb{N}_0} \leftrightarrow: A(\sqrt[n]{z})$
- $(a_{n-1})_{n \in \mathbb{N}_{\geq 1}} \leftrightarrow: z A(z) = a_0 z + a_1 z^2 + \dots$
- $(n a_n)_{n \in \mathbb{N}_0} \leftrightarrow: z A'(z)$

Remark: formal power series are not functions, generating functions are not functions.

E.g. $\frac{1}{1-z} = \sum_{n \geq 0} (-1)^n z^n$ is an equality of FPS.

$$\frac{2}{(1-z)^2} = z \cdot \left(\frac{1}{1-z} \right)' = \sum_{n \geq 0} n z^n$$

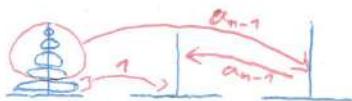
$$\frac{1}{(1-z)^k} = \sum_{n \geq 0} \binom{n+k-1}{k-1} z^n$$

Remark: if $A(z) = B(z)$ as FPS and $A(z)$ and $B(z)$ converge as FPS for $|z| < R$, then $A(z) = B(z)$ as power series.

$\sum_{n \geq 0} n! z^n$ is a FPS which converges only at 0 as a power series.

Why are FPS useful?

E.g.: (Towers of Hanoi)



Goal: Move disks to another peg, but a disk can only be put at a larger disk.

There is a recurrence for $a_n = \# \text{moves required}$. To move n disks: $a_n := 1 + 2a_{n-1}, a_0 = 0$

We want an explicit expression for a_n :

$$a_n = 2a_{n-1} + 1 \quad | z^n$$

$$a_n z^n = 2a_{n-1} z^n + z^n \quad | \quad \cancel{\sum_{n \geq 1}}$$

$$\cancel{\sum_{n \geq 1} (2a_{n-1} + 1) z^n} = \cancel{2 \sum_{n \geq 1} a_{n-1} z^n} + \sum_{n \geq 1} z^n$$

$$\sum_{n \geq 1} a_n z^n = 2 \underbrace{\sum_{n \geq 1} a_{n-1} z^n}_{a_0 z + a_1 z^2 + \dots} + \underbrace{\sum_{n \geq 1} z^n}_{\frac{1}{1-z} - 1}$$

$$=: A(z) - a_0$$

$$\Rightarrow A(z) - a_0 = 2z A(z) + \frac{1}{1-z} - 1 \quad \text{as FPS!}$$

$$(1-2z) A(z) = \frac{1}{1-z} - 1 + a_0 \stackrel{a_0=0}{=} \frac{z}{1-z}$$

$$A(z) = \underbrace{\frac{z}{(1-z)(1-2z)}}_{\text{as FPS! (no function involved)}}$$

$$= \frac{-1}{1-z} + \frac{1}{1-2z} = - \sum_{n \geq 0} z^n + \sum_{n \geq 0} 2^n z^n = \sum_{n \geq 0} \underbrace{(2^{n-1})}_{a_n} z^n$$

E.g.: $F_0 = 0, F_1 = 1, F_{n+2} = F_{n+1} + F_n, n \geq 0$ (Fibonacci)

$$F(z) := \sum_{n \geq 0} F_n z^n$$

$$\sum_{n \geq 0} F_{n+2} z^{n+2} = \sum_{n \geq 0} F_{n+1} z^{n+2} + \sum_{n \geq 0} F_n z^{n+2}$$

$$F(z) - F_0 - F_1 z = z(F(z) - F_0) + z^2 F(z)$$

$$(1 - z - z^2) F(z) = 2F_0 + F_0 + F_1 z$$

$$F(z) = \frac{z}{1-z-z^2} = \frac{\frac{z}{\sqrt{5}} \cdot \frac{1-\sqrt{5}}{2}}{z - \frac{-1+\sqrt{5}}{2}} + \frac{-\frac{z}{\sqrt{5}} \cdot \frac{1+\sqrt{5}}{2}}{z - \frac{-1-\sqrt{5}}{2}}$$

Extensiv Partialbruchzerlegung:

$$\frac{z}{(z-2)(z-5)^2} = \frac{A}{z-2} + \frac{B}{z-5} + \frac{C}{(z-5)^2} \quad | \quad (z=2)$$

$$\frac{z}{(z-5)^2} = A + (z-2) \cdot (\text{something which exists for } z=2)$$

$$z=2: \frac{z}{(z-5)^2} = A + 0 \cdot \text{something} \Rightarrow A = \frac{2}{9} \quad (\text{multiplicity doesn't always work})$$

$$\frac{z}{z-2} = (z-5) \left((z-5) \frac{A}{z-2} + B \right) + C$$

$$z=5: \frac{5}{3} = C \quad \sim (z-5) \text{ doesn't work but we can plug in } C$$

$$\text{Plugging in } C: \frac{z}{(z-2)(z-5)^2} - \frac{5}{3} \frac{1}{(z-5)^2} = \frac{A}{z-2} + \frac{B}{z-5} \quad | \quad (z-5) \text{ now works:}$$

$$= \frac{2 - \frac{5}{3}(z-2)}{(z-2)(z-5)} = -\frac{2}{3} \frac{1}{(z-2)(z-5)}$$

$$z=5: [...]$$

(start topic again.)

In general $a_{n+k} + q_1 a_{n+k-1} + \dots + q_k a_n = 0$ for $n \geq 0$, a_0, \dots, a_{k-1} are given as initial conditions. We are interested in $A(z) = \sum_{n \geq 0} a_n z^n$ | $\cdot z^{n+k}$, $\sum_{n \geq 0}$

$$\Rightarrow \sum_{n \geq 0} a_{n+k} z^{n+k} + q_1 \sum_{n \geq 0} a_{n+k-1} z^{n+k} + \dots + q_k \sum_{n \geq 0} a_n z^{n+k} = 0$$

$$A(z) - a_0 - a_1 z - \dots - a_{k-1} z^{k-1} + q_1 z (A(z) - \sum_{i=0}^{k-2} a_i z^i) + \dots + q_k z^k (A(z) \cancel{- a_0 - \dots - a_{k-1} z^{k-1}}) = 0$$

$$A(z) \cancel{(1 + q_1 z + \dots + q_k z^k)} = p(z), \quad p \text{ a polynomial at most } k-1 \\ (\text{essentially the initial condition})$$

$1-z-z^2$ for Fibonacci

$$F_0 + z(F_1 - F_0)$$

$\therefore q(z)$ contains as information the recurrence, has degree k

$$\Rightarrow A(z) = \frac{p(z)}{q(z)}$$

We have proven that $A(z)$ is a rational function!

partial fraction decomposition: $q(z) = \prod_{i=1}^r (z - z_i)^{\lambda_i}, \quad \sum_{i=1}^r \lambda_i = k$

Fibonacci: $q(z) = (z - \frac{1+\sqrt{5}}{2})(z - \frac{1-\sqrt{5}}{2})$

Ansatz: $\frac{p(z)}{q(z)} = \sum_{i=1}^r \sum_{j=1}^{\lambda_i} \frac{A_{ij}}{(z - z_i)^j}$ ④ goal: expand into generating function
example: $\frac{1}{3z-2} \frac{1}{z-2} = \frac{a}{1-9z} = a \cdot \sum (9z)^n$

\oplus $= \sum_{i=1}^r \sum_{j=1}^{\lambda_i} \frac{A_{ij}}{(z - z_i)^j}$

$\underbrace{p_n(n) \text{ polynomial of degree } \lambda_i - 1}_{\lambda_i}$

$= \sum_{n \geq 0} \left(A_{n1} + \binom{n+1}{1} A_{n2} + \dots + \binom{n+\lambda_i-1}{\lambda_i-1} A_{n\lambda_i} \right) \left(\frac{z}{z_i}\right)^n + \dots + \cancel{\dots}$

$+ \sum_{n \geq 0} \left(A_{11} + \binom{n+1}{1} A_{12} + \dots \right) \left(\frac{z}{z_1}\right)^n + \dots$

$= \sum_{n \geq 0} \underbrace{\left(p_1(n) \left(\frac{1}{z_1}\right)^n + \dots + p_r(n) \left(\frac{1}{z_r}\right)^n \right)}_{= a_n} z^n$

Definition: $X(z) = z^k + q_1 z^{k-1} + \dots + q_k$ is called the characteristic polynomial of the recurrence relation.

Remark: $(X(z) = q(z)|_{z^k u > z^{k-1}})$ so each root is flipped to its multiplicative inverse

$$X(z) = \prod_{i=1}^r (z - \frac{1}{z_i})^{\lambda_i}$$

E.g.: Fibonacci $F_n = \frac{1}{\sqrt{5}} \left(\underbrace{\frac{1+\sqrt{5}}{2}}_{\approx 1.618} \right)^n - \frac{1}{\sqrt{5}} \left(\underbrace{\frac{1-\sqrt{5}}{2}}_{\approx -0.618} \right)^n \in \mathbb{N}$

[for large n : $\Rightarrow F_n \approx \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n$]

unlabelled enumeration

Definition: A binary tree is a rooted tree where each node has no successors or exactly two successors. \mathcal{B} is the set of all binary trees.

$$\mathcal{B} = \{\cdot\} \cup \{ \overset{\wedge}{B_1} \overset{\wedge}{B_2} \mid B_1, B_2 \in \mathcal{B} \}$$

$$B(z) = \sum_{n \geq 0} b_n z^n \text{ where } b_n \text{ is the number of binary trees with } n \text{ internal nodes}$$

$$\Rightarrow B(z) = 1 + \underset{\substack{\text{node added} \\ \nearrow B_1 \cdot B_2}}{z} B(z)^2$$

can't appear because $\frac{1}{2}$ is not a generating function

$$\Rightarrow B(z) = \frac{1 \pm \sqrt{1-4z}}{2} = 1 + z + 2z^2 + 5z^3 + 14z^4 + \dots \quad (\text{Catalan numbers, exercise})$$

dictionary for unlabelled structures: $(A \cup \mathcal{B})(z) = A(z) + \mathcal{B}(z)$

$$(A \times \mathcal{B})(z) = A(z) \cdot \mathcal{B}(z) = \sum_{n \geq 0} \sum_{k=0}^n a_k b_{n-k} z^n$$

size of (a, b) is the size of a + size of b

$$\cancel{\text{empty sequence}} \quad 1 + \cancel{(z + z^2)} + \cancel{(z+z^2)^2} + \dots = \frac{1}{1-(z+z^2)}$$

empty sequence $\underbrace{z}_{\substack{\text{size 1} \\ \text{one object}}} + \underbrace{z^2}_{\substack{\text{size 2} \\ \text{one object}}} \quad \underbrace{(z+z^2)^2}_{\substack{\text{sequences of} \\ \text{length 2 but not} \\ \text{necessary size}}} + \dots$
 pairs of objects

E.g. sequences of 1's and 2's (compositions with parts in $\{1, 2\}$)

$$\emptyset, \overset{1}{1}, \overset{2}{1+1}, \overset{1+1}{2}, \overset{1+1+1}{2+1}, \overset{1+1+1+1}{2+1+1}, \overset{1+2+1}{1+2}, \overset{1+1+2}{1+1}, \overset{2+2}{2+2}$$

(size 0), (size 1), (size 2), (size 3), (size 4)

extract coefficient of
a generating function

[$A(z) = \sum_{n \geq 0} (\# \text{ elements of size } n) z^n$] Notation: $[z^n] A(z) = a_n$

$$(\text{sequences of objects in } A) = 1 + A(z) + A^2(z) + \dots = \frac{1}{1-A(z)}$$

Example: red, blue and yellow balls, 2 or 3 red, at least one blue, at most one yellow, n balls, what's the generating function?

$$\underbrace{(r^2 + r^3)}_{\text{2 or 3 red}} \cdot \underbrace{(b + b^2 + b^3 + \dots)}_{\text{at least one blue}} \cdot \underbrace{(1+y)}_{\text{at most one yellow}}$$

$$= (r^2 + r^3) \cdot (b + b^2 + b^3 + \dots) \cdot (1+y)$$

$$= \left(\frac{b}{1-b} \right)$$

In total: $A(z) = ((r^2 + r^3) \frac{b z}{1-bz}) (1+y z)$

for n balls: $[2^n] A(z) = \text{generating function in } r, b, y$

$$A(z) \Big|_{r=b=y=1} = (2^2 + 2^3) \left(\frac{z}{1-z} \right) (1+z) = \frac{z^3 + 2z^4 + 5z^5}{1-z} \quad \text{so } [2^n] A(z) = 4 \text{ for } n \geq 5$$

Example: combinations without repetitions: balls a_1, a_2, \dots, a_n

generating function for selecting balls (selecting a ball or not)

$$(1+a_1)(1+a_2) \cdots (1+a_n), \text{ set } a_i := z \text{ (not caring which ball)}$$

$$(1+z)^n = \sum_{n \geq 0} \binom{n}{n} z^n$$

$$\text{with repetitions: } (\sum_{n \geq 0} a_1^n)(\sum_{n \geq 0} a_2^n) \cdots (\sum_{n \geq 0} a_n^n), \quad a_i := z \quad \left(\frac{1}{1-z} \right)^n = \sum_{n \geq 0} \binom{n+N-1}{n} z^n$$

Labelled enumeration

$$\text{exponent generating function: } \sum_{n \geq 0} a_n \frac{z^n}{n!} = \hat{A}(z)$$

$$\text{unlabelled enumeration: } \sum_{n \geq 0} a_n z^n = A(z), \quad a_n = \# \text{objects of size } n$$

$$\text{Example: permutations: } A(z) = \sum_{n \geq 0} n! z^n, \quad \text{cyclic permutations: } \hat{A}(z) = \sum_{n \geq 0} n! \frac{z^n}{n!} = \frac{1}{1-z}$$

dictionary:

$$\widehat{(A \cup B)}(z) = \hat{A}(z) + \hat{B}(z)$$

$$\widehat{(A \times B)}(z) = \hat{A}(z) \cdot \hat{B}(z)$$

$$\widehat{(\text{sets of objects})}(z) = e^{\hat{A}(z)}$$

$$\widehat{(\text{cycles of objects in } A)}(z) = \log \frac{1}{1-\hat{A}(z)}$$

Remark: $\widehat{A} \cdot \widehat{B}(z) = \widehat{A}(\widehat{B}(z))$

Definition: Let A and B be sets closed under relabelling. Let $A[1, 2, \dots, n]$ be the set of objects with labels $1, 2, \dots, n$.

Then $A \times B[1, \dots, n]$ is the set of pairs (a, b) with $a \in A, b \in B$ such that the total set of labels is $1, \dots, n$.

Formally, $A \times B[1, \dots, n] = \bigcup_{\substack{U \subseteq \{1, \dots, n\} \\ V = \{1, \dots, n\} \setminus U}} A[U] \times B[V]$

$$[z^n] \widehat{A} \times \widehat{B}(z) = \sum_{k=0}^n \binom{n}{k} a_k b_{n-k}$$

$$\begin{aligned} \widehat{A} \times \widehat{B}(z) &= \sum_{n \geq 0} \sum_{k=0}^n \binom{n}{k} a_k b_{n-k} \frac{z^n}{n!} = \sum_{n \geq 0} \sum_{k=0}^n \frac{\frac{z^k}{k!}}{\frac{z^{n-k}}{(n-k)!}} \cdot \frac{1}{k!} a_k b_{n-k} z^k z^{n-k} \\ &= \sum_{n \geq 0} \sum_{k=0}^n \frac{a_k z^k}{k!} \frac{b_{n-k} z^{n-k}}{(n-k)!} = \sum_{k \geq 0} \sum_{l \geq 0} \frac{a_k z^k}{k!} \frac{b_l z^l}{l!} = \widehat{A}(z) \widehat{B}(z) \end{aligned}$$

Definition: Let A be a set closed under relabelling. Then $\text{set}(A)$ is the set of objects $\{a_1, \dots, a_i, b\}, a_i \in A$, such that the total set of labels is $\{1, \dots, n\}$.

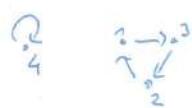
E.g.: sets of cycles with labels $\{1, 2, 3, 4\}$.

Cycles of $\{1\}, \dots$ 

Cycles of $\{1, 2\}, \dots$ 

Cycles of $\{1, 2, 3\}, \dots$ 







E.g.: $A[1, 2, 3] = \{1, 2, 3\}$ is not closed under relabelling. It produces all objects with labels $\{1, 2, 3\}$

E.g.: $A[1, 2] := \{12, 21\}$, $B[1, 2] := \{12, 21\}$

$$A \times B[1, 2, 3, 4] = \{(13, 42), (12, 34), (14, 23), (31, 42), (21, 34), (41, 23), (13, 24), (12, 43), (14, 32), (31, 24), (21, 43), (41, 32)\}$$

cardinality: 24

$$A[\emptyset, \{2\}] = \{\emptyset, \{2\}\}, A[\emptyset] = \{\emptyset\}$$

$$\widehat{\text{sets}}(z) = e^z, \quad e^z = 1 + z + \frac{z^2}{2!} + \frac{z^3}{3!} + \dots$$

$$\widehat{\text{cycles}}(z) = \log \frac{1}{1-z}$$

$$\widehat{\text{sets(cycles)}}(z) = e^{\log \frac{1}{1-z}} = \frac{1}{1-z} = \widehat{\text{permutations}}(z)$$

$B :=$ sets of non-empty sets

$\widehat{B}(z) = e^{(e^z-1)}$ is the (exponential) generating function for set partitions.

Partially ordered sets

Definition: A partial order is a set P together with a relation \leq , such that

• $a < b \Rightarrow b \neq a$ (antisymmetry)

• $a \neq a$

• $a < b, b < c \Rightarrow a < c$ (transitivity) (poset, partially ordered set)

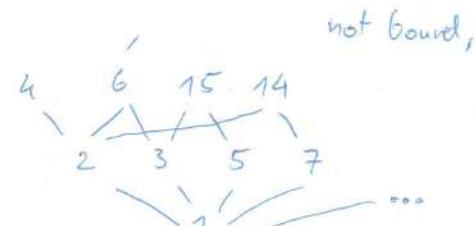
Notation: $a \leq b$ means $a < b$ or $a = b$.

The Hasse-diagram is the digraph with vertices P and (a, b) is an arc if $a < b$ and $\nexists c: a < c < b$.

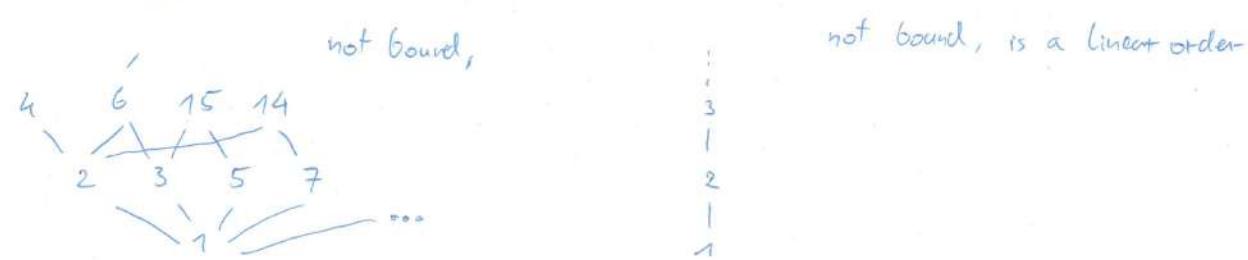
Remark: equivalently define (P, \leq) with \leq reflexive, transitive.

E.g.:

$$(N_{\geq 1}, \mid)$$



$$(N_{\geq 1}, \leq)$$



Definition A minimal element is an $a \in P$ such that $\forall b \in P : a \leq b$,
a maximal element is an $a \in P$, such that $\forall b \in P : b \leq a$.

Definition: An interval is a subset $[x, y] := \{z \mid x \leq z \leq y\}$ of P

P is bounded if $\exists M \subseteq P : \forall x \in P \exists y \in M : x \leq y$,

$\exists M \subseteq P : \forall x \in P \exists y \in M : y \leq x$

(P, \leq) is locally finite if all intervals have finite cardinality.

$$|[x, y]| < \infty \quad \forall x, y \in P.$$

E.g.:



is not locally finite

$$P = \{x \in Q \mid 0 \leq x \leq 1\}, \quad a \leq b \text{ if } a=0 \text{ or } b=1$$

Definition Let (P, \leq) be locally finite, bounded (or at least bounded below)

$$f : P \rightarrow \mathbb{R}, \quad \text{def} \quad S_f(x) := \sum_{z \leq x} f(z)$$

E.g.:

$$S_f(x) = f(a) + f(c) + f(b) + f(x)$$

Goal: find f for given $S_f(x)$.

Definition: Let (P, \leq) be a poset, locally finite with a minimal element "0".

$\mu : P \times P \rightarrow \mathbb{R}$ is called Möbius function of P if and only if

$$\forall x, y : \sum_{z \in [x, y]} \mu(z, y) = \delta_{xy} = \begin{cases} 0 & x \neq y \\ 1 & x = y \end{cases}$$

Remark This relation determines μ uniquely.

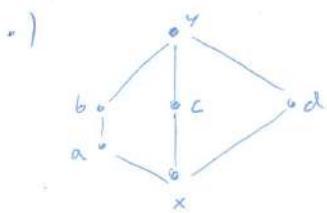
If $x \neq y$, then $\mu(x, y) = 0$

Examples: *) $[x, x] = \{x\} \Rightarrow \mu(x, x) = 1$

*) $[x, y] = \{x, y\} \Rightarrow \mu(x, y) + \mu(y, y) = 0, \mu(y, y) = 1 \Rightarrow \mu(x, y) = -1$

*) $(\mathbb{N}, \leq) : \mu(n, n) = 1, \mu(n, n+1) = -1, \mu(n, m) = 0 \quad \forall m > n+2 \text{ or } m < n$

because $\mu(n, n+2) = -\underbrace{\mu(n, n+1)}_{-1} - \underbrace{\mu(n, n)}_1$



$$\begin{aligned}\mu(x, y) &= - \left(\underbrace{\mu(a, y)}_{-1} + \underbrace{\mu(b, y)}_{-1} + \underbrace{\mu(c, y)}_{-1} + \underbrace{\mu(d, y)}_{-1} + \underbrace{\mu(y, y)}_1 \right) \\ \mu(a, y) &= - \left(\underbrace{\mu(b, y)}_{-1} + \underbrace{\mu(c, y)}_1 \right) = 0 \\ \Rightarrow \mu(x, y) &= 2\end{aligned}$$

Definition: Let $(P_1, \leq_1) \times (P_2, \leq_2) := (P_1 \times P_2, \leq_3)$ (product poset) has
 $(x_1, x_2) \leq_3 (y_1, y_2) \Leftrightarrow (x_1 \leq_1 y_1) \text{ and } (x_2 \leq_2 y_2)$

Theorem If P_1 and P_2 both have a unique minimal element, then $P_1 \times P_2$ has a unique minimal element and $\mu_{P_1 \times P_2}(x, y) = \mu_{P_1}(x_1, y_1) \mu_{P_2}(x_2, y_2)$

Proof is left as exercise (plug in definition).

Example: $A = \{a_1, \dots, a_n\}$, consider $(2^A, \subseteq) \cong (\{0,1\}^n, \leq)$ componentwise not lexicographic
 binary strings of length n
 e.g. $x = \{a_1, a_5\}, y = \{a_1, a_2, a_3, a_5\}, n=5$
~~then~~ $x = 01001, y = 11101$

$$\begin{aligned}\text{then } \mu(x, y) &= \mu(0, 1)\mu(1, 1)\mu(0, 1)\mu(0, 0)\mu(1, 1) \\ &= (-1) \cdot 1 \cdot (-1) \cdot 1 \cdot 1 = 1\end{aligned}$$

(for this example)

In general: $\mu(x, y) \in \{-1, 1\}$ (if x, y are comparable, $x \leq y$)

$$\Rightarrow \mu(x, y) = (-1)^{\# \text{different positions}} = (-1)^{|x|-|y|} = (-1)^{|y|-|x|} = (-1)^{|y-x|}$$

Theorem (Möbius Inversion): Let (P, \leq) be locally finite with "0" (unique minimal element).

$$f: P \rightarrow \mathbb{R}, \quad S_f(x) = \sum_{z \in [0, x]} f(z) \quad \Rightarrow \quad f(x) = \sum_{z \in [0, x]} S_f(z) \mu(z, x)$$

$$\begin{aligned}\text{Proof: } \sum_{0 \leq z \leq x} S_f(z) \mu(z, x) &= \sum_{0 \leq z \leq x} \sum_{0 \leq y \leq z} f(y) \mu(z, x) \stackrel{\text{double counting}}{=} \sum_{0 \leq y \leq x} \sum_{y \leq z \leq x} f(y) \mu(z, x) \\ &= \sum_{y \in [0, x]} f(y) \sum_{z \in [y, x]} \mu(z, x) = \sum_{y \in [0, x]} f(y) \cdot \delta_{yx} = f(x)\end{aligned}$$

□

Example: (\mathbb{N}, \leq) , $\mu(m,n) = \begin{cases} 1 & m=n \\ -1 & m+1=n \\ 0 & \text{otherwise} \end{cases}$

$$f: \mathbb{N} \rightarrow \mathbb{R}, \quad S_f(n) = \sum_{k=0}^n f(k) \Rightarrow f(n) = \sum_{k=0}^n S_f(k) \mu(k, n) = \begin{cases} S_f(n) - S_f(n-1), \\ S_f(n) \quad \text{if } n=0 \end{cases}$$

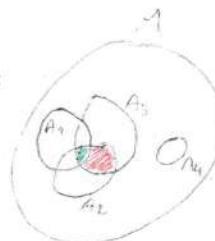
Example (PIE): $A_1, \dots, A_m \subseteq M$, consider $(2^{\{1, \dots, m\}}, \leq)$ poset of indices

$$\mathcal{I} \subseteq \{1, \dots, m\}, \quad f(\mathcal{I}) := |\bigcap_{i \in \mathcal{I}} A_i \cap \bigcap_{j \in \{1, \dots, m\} \setminus \mathcal{I}} \overline{A_j}| \quad \cancel{|\text{elements precisely in all } A_i, i \in \mathcal{I}|}$$

$$S_f(\mathcal{I}) = \sum_{J \supseteq \mathcal{I}} f(J) = |\bigcap_{i \in \mathcal{I}} A_i| \quad \cancel{|\text{elements in } A_i \text{ for } i \in \mathcal{I}|} \quad \text{e.g.:}$$

$\bigcap_{i \in \mathcal{I}} A_i$

$= |\text{elements in } A_i \text{ for } i \in \mathcal{I}|$



$$f(2, 3) = |A_2 \cap A_3 \cap \overline{A}_4 \cap \overline{A}_1|$$

$$S_f(2, 3) = \underline{f(2, 3)} + \underline{f(1, 2, 3)} + \\ + \cancel{f(1, 2, 3, 4)} + \cancel{f(1, 2, 3, 4)}$$

empty

Möbius inversion: $f(\mathcal{I}) = \sum_{J \supseteq \mathcal{I}} S_f(J) \mu(J, \mathcal{I}) = \sum_{J \supseteq \mathcal{I}} (-1)^{|\mathcal{I} \cup J|} |\bigcap_{j \in J} A_j|$

In particular: $f(\emptyset) = |\bigcap_{j \in \{1, \dots, m\}} \overline{A_j}| = \sum_{J \subseteq \{1, \dots, m\}} (-1)^{|J|} |\bigcap_{j \in J} A_j|$

Example: "classical" number theoretic Möbius function.

$$(\mathbb{N}_+, \mid), \quad m = p_1^{e_1} \cdots p_r^{e_r}, \quad p_i \in \mathbb{P}, \quad e_i \in \mathbb{N}_0$$

$$n = p_1^{f_1} \cdots p_s^{f_s}, \quad m \mid n \Leftrightarrow e_i \mid f_i \quad \forall i$$

def

$$\cong (\mathbb{N}_0, \leq) \times (\mathbb{N}_0, \leq) \times \dots$$

finite entries $\neq 0$ in the infinite product

$$\mu(n) := \prod_{i=1}^r \mu_{(\mathbb{N}_+, \mid)}(1, n) = \mu_{(\mathbb{N}_+, \mid)}(0, e_1) \mu_{(\mathbb{N}_+, \mid)}(0, e_2) \cdots \mu_{(\mathbb{N}_+, \mid)}(0, e_r) \underbrace{\mu_{(\mathbb{N}_+, \mid)}(0, 0)}_{=1} \cdots$$

$$\mu(0, k) = \begin{cases} 1 & k=0 \\ -1 & k=1 \\ 0 & k>1 \end{cases} \Rightarrow \mu(n) = \begin{cases} 1 & n=1 \\ (-1)^r & n=p_1 \cdots p_r, \quad p_i \neq p_j \text{ for } i \neq j \\ 0 & \text{otherwise} \end{cases}$$

conclusion:

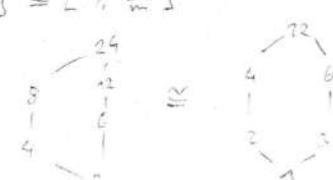
$$f: \mathbb{N}_+ \rightarrow \mathbb{R}, \quad S_f(n) = \sum_{d \mid n} f(d)$$

$$\Rightarrow f_{\text{conv}} = \sum_{d \mid n} S_f(d) \mu\left(\frac{n}{d}\right)$$

Remark $\mu_{(\mathbb{N}_+, \mid)}(m, n) = \mu\left(\frac{m}{n}\right)$

e.g.: $[m, n] \cong [1, \frac{m}{n}]$

$[2, 24]$



Lattices (Verband in DC)

Definition: (P, \leq) poset, $x, a, b \in P$, $a \leq x \leq b$ ($\Rightarrow a$ lower bound/ b upper bound for x). Let $x \vee y$ ("join") be the smallest common upper bound of x and y and $x \wedge y$ ("meet") the largest common lower bound of x and y , if it exists.

Notation: $\bigvee_{x \in V} x := x_1 \vee x_2 \vee \dots \vee x_n$

basic properties: $x \vee y = y \vee x$, $x \vee x = x$, $(x \vee y) \vee z = x \vee (y \vee z)$, $a \vee (a \wedge b) = a = a \wedge (a \vee b)$

Definition: L is a lattice if $\forall x, y \in L : \exists x \vee y$ and $x \wedge y$.

L is a join-semilattice if $\forall x, y \in L : \exists x \vee y$, analogue meet-semilattice

L is a complete lattice if $\forall X \subseteq L : \exists \bigvee_{x \in X} x$ and $\bigwedge_{x \in X} x$

Example: $(2^M, \subseteq)$ is a lattice: $A, B \subseteq M \Rightarrow A \vee B := A \cup B$, $A \wedge B := A \cap B$

Lemma: L lattice, $x, y, s, t \in L$

$$x \leq s, y \leq s \Rightarrow x \vee y \leq s$$

$$x \geq t, y \geq t \Rightarrow x \wedge y \geq t$$

$$x \leq y \Leftrightarrow x \vee y = y \Leftrightarrow x \wedge y = x$$

1 is the neutral el. for \wedge ,
0 is neutral for \vee

Lemma L a finite meet-semilattice with a 1-element (larger than all other elements)

Then L is a lattice.

Proof: $x, y \in L$, $B = \{u \in L \mid x \leq u, y \leq u\}$, $1 \in B \Rightarrow B \neq \emptyset$

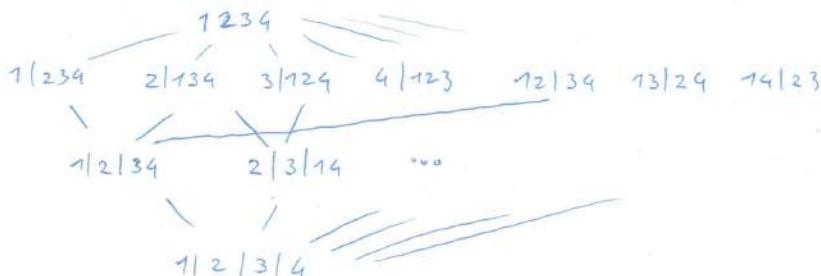
$$|B| < \infty \Rightarrow B = \{u_1, \dots, u_m\}, u := \bigwedge_{i=1}^m u_i \in B$$

$\Rightarrow u = x \vee y$ (minimal element in B) (not shown: unique) \square

Example: $\Pi_n = \{\pi \text{ a set partition of } \{1, \dots, n\}\}$

$(\Pi_n, \text{refinement})$ is a lattice. 1 is the set partition with 1 block: $\{\{1, \dots, n\}\}$, 0 is the set of all singletons $\{\{1\}, \dots, \{n\}\}$.

$\alpha, \beta \in \Pi_n : \alpha \wedge \beta$: set partition, where i, j are in the same block if and only if they are in the same block in α and β



Theorem L lattice with, $0, 1$, $b \neq 1, b \in L$

$$\Rightarrow \mu(0, 1) = - \sum_{\substack{x: x \wedge b = 0 \\ x \neq 0}} \mu(x, 1)$$

Proof:

~~$\sum_{x: x \wedge b = 0} \mu(x, 1)$~~ $\oplus \quad (\star) \quad 0 = \sum_{x: x \wedge b = 0} \mu(x, 1)$

$$N(y) := \sum_{x: x \wedge b = y} \mu(x, 1) \quad \text{for } y \leq b$$

Given x and $b \rightarrow y = x \wedge b \leq b$
only one y

sum over an interval
of the Möbius function

$$S(b) = \sum_{y: y \leq b} N(y) = \sum_{y \leq b} \sum_{x: x \wedge b = y} \mu(x, 1) \stackrel{!}{=} \sum_{x \in L} \mu(x, 1) = \sum_{x \in L \setminus \{0\}} \mu(x, 1) = 0$$

$$\stackrel{\text{Möbius inversion}}{\Rightarrow} N(b) = \sum_{y \leq b} S(y) \mu(y, b) = 0 \quad \Rightarrow \quad N(0) = 0$$

□

$$\text{Corollary: } \mu_{\prod_b}(0, 1) = (-1)^{n-1} (n-1)!$$

Proof: exercise (induction, choose $b = \{1, \dots, n-1\}, \{n\}$ (split of largest element))

③ Number theory

Definition $a, b \in \mathbb{Z}$, then $a|b$: $\Leftrightarrow \exists c \in \mathbb{Z} : a \cdot c = b$ (a divides b)

(more generally for anything \mathbb{A} , e.g. $\mathbb{Z}[\mathbf{x}]$, $\mathbb{Z}_m = \mathbb{Z}/m\mathbb{Z}$, ...)

Definition: $a, b \in \mathbb{Z}$, $d = \gcd(a, b) \Leftrightarrow$

-) $d|a, d|b$ (greatest common divisor)
-) $+1|a, +1|b \Rightarrow +1|d$

$b > 0 \Rightarrow \exists q, r \in \mathbb{Z} : \bullet) a = bq + r \quad (\text{division with } \cancel{\text{rest}} \text{ remainder})$
 $\bullet) 0 \leq r < b$

Euclidean algorithm (theorem)

$$a = bq_0 + r_0$$

$$\Rightarrow b > r_0 > r_1 > \dots > r_k > 0$$

$$b = r_0 q_1 + r_1$$

$$\Rightarrow r_1 = \gcd(a, b)$$

$$r_0 = r_1 q_2 + r_2$$

⋮

$$r_{k-2} = r_{k-1} q_k + r_k \quad ; \quad r_{k-1} = r_k q_{k+1} \quad (\text{no remainder } \cancel{\text{any more}})$$