

Discrete Mathematics

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1 Graph Theory

1.1 Definition & Notation

- $G = (V, E)$
- Directed Edge $e = (v, w)$; $v \rightarrow w$
- Undirected Edge $e = \{v, w\} = \{w, v\} = vw$
- $d(v)$... degree
- $d^+(v)$... out-degree
- $d^-(v)$... in-degree
- $\Gamma(v)$... set of neighbours
- $\Gamma^+(v)$... successors
- $\Gamma^-(v)$... predecessors
- Walk, Trail Circuit**
- Walk:** sequence of edges (no jumps)
- Trail:** $v \rightarrow w$; $v \rightarrow w$
- Trail:** sequence of edges (no jumps, no repeating edges)
- Circuit/Closed Trail:** same start and end vertex
- Path :** no repeating edges, no repeating vertices
- Cycle/Closed Path:** same start and end vertex

Handshaking Lemma

$$\sum_{v \in V} d(v) = 2|E| \text{ or } \sum_{v \in V} d^+(v) = \sum_{v \in V} d^-(v) = |E|$$

1.2 Trees & Forests

Isomorphism

$G \cong H \Leftrightarrow \exists \varphi: V(G) \rightarrow V(H) \text{ and } w \in E(G) \Leftrightarrow \varphi(v) \varphi(w) \in E(H)$

Spanning Subgraphs & Spanning Forests

Spanning Forest \Leftrightarrow

- $V(F) = V(G)$ and $E(F) \subseteq E(G)$
- F is a forest
- F has the same connected components as G

If G connected $\Rightarrow F$ spanning tree

Minimum or Maximum Spanning Trees

- $G = (V, E)$ and weight function $w: E \rightarrow \mathbb{R}, e \mapsto w(e)$
- Then $T = (V, F) \subseteq G$ minimum spanning tree if $w(F)$ minimal
- GREEDY** (V, E, w) constructs the minimum spanning tree (MST):
- Sorts the edges of G by weight ascending
- Adds the next smallest edge to T (such that T remains a tree) until G is spanned

Matrix-Tree-Theorem (Kirchhoff)

$G = (V = \{v_1, \dots, v_n\}, E)$ simple, undirected, connected, with adjacency matrix A
 $D = \text{diag}(d(v_1), \dots, d(v_n))$; $(D - A)^{-1} := (D - A)$ without 1 row and 1 column
 $\Rightarrow \det((D - A)^{-1}) = \text{number of spanning trees of } G$

Independence Systems

Independence System (E, S) : $S \subseteq 2^E$ and S is closed under inclusion.
 $A \in S \wedge B \subseteq A \Rightarrow B \in S$

Matroids & GREEDY

An independence system $M = (E, S)$ is called a matroid if $\forall A, B \in S$ such that $|B| = |A| + 1$: $\exists e \in B \setminus A$ with $A \cup \{e\} \in S$
 $(E, S) = (V, E)$ be an undirected graph and $S = \{F \subseteq E \mid F \text{ is a forest}\}$; then (E, S) is a matroid.
 If $M = (E, S)$ is a matroid with a weight function $w: E \rightarrow \mathbb{R}$ then GREEDY computes a A maximal/minimal weight with respect to the inclusion.

1.3 Special Graphs

Planar Graphs

Definition: A graph G is planar if there is an isomorphic graph H embedded in the plane (vertices are points in the plane \mathbb{R}^2) such that no edges intersect. A graph is planar \Leftrightarrow there is no subgraph which is a subdivision of K_4 or of $K_{3,3}$.
Euler's Polyhedron Formula: if G is connected and planar graph, then $\alpha_0 - \alpha_1 + \alpha_2 = 2$. α_0 ... Vertices, α_1 ... Edges, α_2 ... Faces
Dual Graphs: Let $G = (V, E)$ be a planar graph and let F be the set of its faces. Then $G^* = (V^*, E^*)$ is defined such that $V^* = F$ and for every edge $e \in E$, set $e^* = (f_1, f_2)$ if f_1 and f_2 are the faces left and right of e .
 G^* is called the dual of G .

Bipartite Graphs and Matching

Definition: Let $G = (V, E)$ be a simple undirected graph. G is called bipartite if and only if:

$$V = V_1 \cup V_2, V_1 \cap V_2 = \emptyset$$

$$vw \in E \Rightarrow v \in V_1, w \in V_2$$

Hall's Marriage Theorem: Given a bipartite graph $G = (V, E)$, such that $V = W \cup M$ where W and M are finite and nonempty. Define the friendship relation: $F \subseteq W \times M$ with $wm \in F \Leftrightarrow wFm$. A feasible marriage is a matching of W $F_1 \subseteq F$ ($\forall w \in W: \exists y \in M$ such that wF_1y). Now the theorem states the following: there is a feasible marriage, if and only if: $\forall W_0 \subseteq W: |\{y \in M \mid \exists x \in W_0: xFy\}| \geq |W_0|$ or in different words: $|\bigcup_{w \in W_0} \Gamma(w)| \geq |W_0|$

1.4 Graph Colorings

Vertex Coloring: $c: V \rightarrow C$ where C is a set of colors. feasible if $vw \in E \Rightarrow c(v) \neq c(w)$

Edge Coloring: $\tilde{c}: E \rightarrow C$

everything that can be done with a vertex coloring can also be done with an edge coloring

- Chromatic Number:** $\chi(G)$ is the minimal number of colors such that there is a feasible coloring.
- $\chi(K_n) = n$
- $\chi(K_{n,m}) = 2$
- $\chi(T) = 2$ if T is a tree and $|V| > 1$
- $\chi(G) = 1 \Leftrightarrow E(G) = \emptyset$
- $\chi(G) = 2 \Leftrightarrow E(G) \neq \emptyset$ and G is bipartite
- $\chi(G) = 2 \Leftrightarrow E(G) \neq \emptyset$ and all cycles have even length
- $\chi(G) \leq 1 + \max_{v \in V} d(v)$
- G is planar $\Rightarrow \chi(G) \leq 5$

Ramsey Theory

- Definition:** The Ramsey Number $R(r, s)$ is the minimum n such that every red-blue coloring of K_n contains either a red K_r or a blue K_s .
- Coloring is edge coloring
- $R(r, s) \leq R(r-1, s) + R(r, s-1)$
- $R(n_1, n_2, \dots, n_r) = \min\{n \mid \text{all } r\text{-edge colorings of } K_n \text{ (colors } c_1, \dots, c_r) \text{ have a } c_j\text{-colored } K_{n_j} \text{ for some } j\}$

2 Advanced Combinatorics

Double Counting

Given two sets A and B and a relation $R \subseteq A \times B$, such that $aRb \Leftrightarrow (a, b) \in R$. Let $R_a = \{b \in B \mid aRb\}$ (all b related to a) and $S_b = \{a \in A \mid aRb\}$. Then: $|R| = \sum_{a \in A} |R_a| = \sum_{b \in B} |S_b|$

Pigeon Hole Principle

Let A_1, \dots, A_n be finite pairwise disjoint sets, $|A_1| \cup \dots \cup A_n| > k \cdot r$, for $r \in \mathbb{N}$. This implies: $\exists i: |A_i| > r$. If $r = 1$, then it follows that: $f: A \mapsto B, |A| > |B| \Rightarrow \exists b \in B: |f^{-1}(b)| \geq 2$

Principle of Inclusion & Exclusion

$|A_1 \cup \dots \cup A_n| = \sum_{i=1}^n |A_i| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k| - \dots + (-1)^{n+1} |A_1 \cap \dots \cap A_n|$
 For $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$
 If the sets are not necessarily pairwise disjoint, then:

$$A_1, \dots, A_n \subseteq I \Rightarrow \bigcap_{i=1}^n A_i = |A| + \sum_{\emptyset \neq I \subseteq \{1, \dots, n\}} (-1)^{|I|} \left| \bigcap_{i \in I} A_i \right| = \sum_{I \subseteq \{1, \dots, n\}} (-1)^{|I|} \left| \bigcap_{i \in I} A_i \right|$$

2.1 Enumerative Combinatorics

Counting Sets

- 6 Basic Problems:**
- # of permutations: $n!$
 - # of k -subsets of A : $\binom{n}{k}$
 - # of ordered k -subsets: $k! \binom{n}{k}$
 - # of k -multisets (elements can be used more than once): $\binom{n+k-1}{k-1}$
 - # of arrangements of a multiset (element b_i appears k_i times): $\frac{n!}{k_1! k_2! \dots k_r!}$
 - # of ordered k -multisets: n^k
- Stars and bars:** Number of assignments of n items into k boxes $\binom{n+k-1}{k-1}$
- Stirling Numbers**
- First kind:** $s_{n,k}$ = # of permutations of $\{1, \dots, n\}$ with k cycles
- $s_{n,1} = (n-1)!$, $s_{n,n-1} = \binom{n}{2}$, $s_{n,n} = 1$
- $s_{n,0} = 0$, $s_{0,0} = 0$, $s_{0,0} = 1$ (!)
- $\sum_{k=1}^n s_{n,k} = n!$
- $\forall n, k > 0: s_{n,k} = s_{n-1,k-1} + (n-1)s_{n-1,k}$
- Second kind:** $S_{n,k}$ = # of k partitions of $\{1, \dots, n\}$
- $\Leftrightarrow \forall i: A_i \neq \emptyset, i \neq j \Rightarrow A_i \cap A_j = \emptyset, \bigcup_{i=1}^k A_i = A$
- $S_{n,1} = 1$, $S_{n,n} = \binom{n}{k}$, $S_{n,n} = 1$
- $S_{n,0} = 0$, $S_{0,0} = 0$, $S_{0,0} = 1$ (!)
- $S_{n,2} = 2^{n-1} - 1$
- $\forall n, k > 0: S_{n,k} = S_{n-1,k-1} + kS_{n-1,k}$

2.2 Generating Functions

Definition & Properties

Let $(a_n)_{n \geq 0}$, $a \in \mathbb{C}$ be a sequence; Then $A(z) = \sum_{n \geq 0} a_n z^n$ is the **generating function** of a_n .

- $A(z) + B(z) = \sum_{n \geq 0} (a_n + b_n) z^n$
- $A(z)B(z) = \sum_{n \geq 0} c_n z^n$ (linearity)
- $A(z)/B(z) = C(z)$, when $b_0 \neq 0$

Operations on Generating Functions

- $(\alpha a_n + \beta b_n)_{n \geq 0} \Leftrightarrow \alpha A(z) + \beta B(z)$, $\alpha, \beta \in \mathbb{C}$
- $(a_{n-1})_{n \geq 1} \Leftrightarrow zA(z)$ (index shift)
- $(na_n)_{n \geq 0} \Leftrightarrow zA'(z)$ (derivative)
- $\frac{1}{1-x} z^n = (1-x)^{-1} = \sum_{k \geq 0} \binom{n+k-1}{k} z^k$
- $(a^n)_{n \geq 0} \Leftrightarrow \frac{1}{1-az}$ $= \sum_{n \geq 0} a^n z^n$
- $\frac{1-x^{n+1}}{1-x} = \sum_{k=0}^n x^k$ (transform back)

Unlabelled Combinatorial Structures

Combinatorial structure: $(A, |\cdot|)$, $A \neq \emptyset$; $|\cdot|$ is weight function; $\forall x \in A: |x|$

- $|x| \in \mathbb{N}$
- Counting sequence** of A : $(a_n)_{n \geq 0}$, $a_n = |\{x \in A: |x| = n\}| < \infty$; "How many x of weight n ?"
- The counting sequence has a generating function $A(z)$, also called generating function of A .
- $C = A \cup B, A \cap B = \emptyset$; $(\hat{A}, |\cdot|_A), (\hat{B}, |\cdot|_B)$; $|x|_C = \begin{cases} |x|_A & \text{if } x \in A \\ |x|_B & \text{if } x \in B \end{cases}$
 $c_n = a_n + b_n \Leftrightarrow C(z) = A(z) + B(z)$
 - $C = A \times B$, $(x, y) \in C$, $x \in A, y \in B$; $|(x, y)| = |x|_A + |y|_B$
 $c_n = \sum_{k=0}^n a_k b_{n-k} \Leftrightarrow C(z) = A(z) \cdot B(z)$
 - $C = \text{seq}(A) = \{\varepsilon\} \cup \{(x_1, x_2, \dots, x_k)_{k \geq 0}\} = \{\varepsilon\} \cup A \cup A^2 \cup A^3 \cup \dots$; $|x|_C = \sum_{i=1}^k |x_i|_A$
 $C(z) = \frac{1}{1-A(z)}$, only if we define $a_0 = 0$
 - $C = \text{seq}_{\leq p}(A) = \{\varepsilon\} \cup \{(x_1, x_2, \dots, x_k)_{k \leq p}\}$, $C(z) = \frac{1-A(z)^{p+1}}{1-A(z)}$ (finite sequence)
 - $C = \text{set}(A) = 2^A$ (power set)
 $C(z) = \exp\left(\sum_{k \geq 1} \frac{(-1)^{k-1}}{k} A(z)^k\right)$

Labelled Combinatorial Structures

- Labelled combinatorial structure:** Elements are numbered \Rightarrow switching of two equal elements does make a difference! $\hat{A}(z) = \sum_{n \geq 0} \sum_{m!} a_n \frac{z^n}{m!}$
- $\hat{C} = \hat{A} \cup \hat{B}, \hat{A} \cap \hat{B} = \emptyset$; $(\hat{A}, |\cdot|_A), (\hat{B}, |\cdot|_B)$; $|x|_C = \begin{cases} |x|_A & \text{if } x \in A \\ |x|_B & \text{if } x \in B \end{cases}$
 $c_n = a_n + b_n \Leftrightarrow \hat{C}(z) = \hat{A}(z) + \hat{B}(z)$
 - $\hat{C} = \hat{A} \times \hat{B}$, $(x, y) \in C$, $x \in \hat{A}, y \in \hat{B}$; x, y labelled in order preserving way; $|(x, y)| = |x|_A + |y|_B$; $\hat{C}(z) = \hat{A}(z) \cdot \hat{B}(z)$
 - $\hat{C} = \text{seq}(\hat{A}) = \{\varepsilon\} \cup \{(x_1, x_2, \dots, x_k)_{k \geq 0}\} = \{\varepsilon\} \cup \hat{A} \cup \hat{A}^2 \cup \hat{A}^3 \cup \dots$
 $\hat{C}(z) = \frac{1}{1-\hat{A}(z)}$
 - $\hat{C} = \text{set}(\hat{A}) = 2^{\hat{A}}$ (labelled power set)
 $\hat{C}(z) = \sum_{k \geq 0} \frac{\hat{A}(z)^k}{k!} = \exp(\hat{A}(z))$
 - $\hat{C} = \text{cyc}(\hat{A})$; $\hat{C}(z) = \sum_{k \geq 1} \frac{\hat{A}(z)^k}{k} = \log\left(\frac{1}{1-\hat{A}(z)}\right)$
 - $\hat{C} = \text{set}(\text{cyc}(\hat{A})) = \frac{1}{1-\hat{A}(z)}$

3 Number Theory

3.1 Divisibility & Factorization

GCD

- Let $a, b \in \mathbb{Z}$; then $d = \gcd(a, b) \Leftrightarrow$
- $d|a$ and $d|b$ (makes it divisor)
 - $t|a$ and $t|b \Rightarrow t|d$ (makes it greatest)

$a, b \in \mathbb{Z} \setminus \{0\}$, $d = \gcd(a, b) \Leftrightarrow \exists c \in \mathbb{Z}$ such that $d = ac + bf$ (*Bézout's Identity*)

Commutative Ring (with 1)

- $(R, +)$ is an **Abelian group** (0-element, $\forall a: -a$)
- (R, \cdot) is a **semigroup** (additionally: \exists neutral element, distributive law holds)
- R is an **integral domain** $\Leftrightarrow \forall a, b \in R \setminus \{0\}: a \cdot b = 0$

Euclidean Ring

R is a **Euclidean ring** if R is an integral domain and there is a *euclidean function* $n: R \rightarrow \mathbb{N}$ such that $\forall a, b \in R$ with $b \neq 0$ there exists $q, r \in R$:

- $a = b \cdot q + r$ with $r = 0$ or $n(r) < n(b)$
- $\forall a, b \in R \setminus \{0\}: n(a) \leq n(a \cdot b)$

Units: $U = \{a \in R: \exists a^{-1}\}$

Prime Numbers

- $p \in \mathbb{P} \Rightarrow \pm 1$ and $\pm p$ are the only divisors of p
- let $n \in \mathbb{N}^+$, then $\exists p_1, p_2, \dots, p_r \in \mathbb{P}$ such that $p_1^{a_1} \cdot p_2^{a_2} \cdot \dots \cdot p_r^{a_r} = n$, for $r \geq 0$
- $\nu_p(n)$... multiplicity of p in the factorization of n (*Factorization is unique up to the order of the factors*)
- $\gcd(a, b) = \prod_{p \in \mathbb{P}} p^{\min(\nu_p(a), \nu_p(b))}$
- $\text{lcm}(a, b) = \prod_{p \in \mathbb{P}} p^{\max(\nu_p(a), \nu_p(b))}$
- $a|b \Leftrightarrow \forall p \in \mathbb{P}: \nu_p(a) \leq \nu_p(b)$
- $n = \prod_{p \in \mathbb{P}} p^{\nu_p(n)}$
- $|\mathbb{P}| = \infty$

3.2 Congruence Relations and Residue Classes

A *residue class* is defined as: $a + m \cdot \mathbb{Z} = \bar{a}$, where $a + m \cdot \mathbb{Z} = \{a + k \cdot m \mid k \in \mathbb{Z}\}$ (notice $\bar{a} \subseteq \mathbb{Z}$ and $\bar{a} = \overline{a+m}$). Definition: $\bar{a} + b = \overline{a+b}$ and $\bar{a} \cdot b = \overline{a \cdot b}$

Let $\bar{a} \in \mathbb{Z}_m$ and then the *inverse element* is defined as follows: $\bar{x} \in \mathbb{Z}_m$ such that: $\bar{x} \cdot \bar{a} = \bar{1}$ then $\bar{x} = \bar{a}^{-1}$.

$\exists \bar{a}^{-1} \in \mathbb{Z}_m \Leftrightarrow \gcd(a, m) = 1$

Prime Residue Classes

$\mathbb{Z}_m^* = \{\bar{x} \in \mathbb{Z}_m \mid \gcd(a, m) = 1\}$ contains all invertible elements, can further be defined as $\mathbb{Z}_m^* = \{x \in \mathbb{Z}_m \mid \exists x^{-1}: x \cdot x^{-1} = 1\}$ (also called **Group of Units**)

Some rules:

- $a \cdot b \equiv ac(am) \Rightarrow a \equiv c(m)$
- $a \cdot b \equiv ac(m) \Rightarrow b \equiv c(m)$ if $ax \equiv 1(m)$ has a solution $\Leftrightarrow \gcd(a, m) = 1$

- ### 3.3 Systems of Congruences
- #### Chinese Remainder Theorem
- Given a system of congruence equations $x \equiv a_i(m_i)$ where $1 \leq i \leq r$ and if $i \neq j$ then $\gcd(m_i, m_j) = 1$, then the system has a unique solution, modulo $m = \prod_{i=1}^r m_i$. The solution is given by: $x \equiv \sum_{i=1}^r a_i b_i \cdot b_j \cdot a_j(m)$. Where $b_j = \left(\frac{m}{m_j}\right)^{-1} \pmod{m_j}$.
- The chinese remainder theorem can be applied to all **euclidean rings**

3.4 Euler-Fermat Theorem

- #### Euler's Totient Function
- $\varphi(m) = |\{x \mid 0 \leq x \leq m-1, \gcd(x, m) = 1\}| = |\mathbb{Z}_m^*|$
- Two cases:
- Suppose $m \in \mathbb{P}$, then $\varphi(m) = m - 1$
 - $m = p^k$, for $p \in \mathbb{P}$ and $k \geq 1$: Now $\bar{x} \in \mathbb{Z}_m \Rightarrow \gcd(a, p^k) = \begin{cases} 1 & \text{if } p \nmid x \\ p^i & \text{with } 1 \leq i < k \end{cases}$
- Let $m = p_1^{a_1} p_2^{a_2} \dots p_r^{a_r}$, then: $\varphi(m) = m \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \dots \left(1 - \frac{1}{p_r}\right)$
- #### Fermat's Little Theorem
- If $\gcd(a, m) = 1$ then $a^{\varphi(m)} \equiv 1(m)$. In particular: $p \in \mathbb{P}, p \nmid a \Rightarrow a^{p-1} \equiv 1(p)$
- Theorem: Let $p, q \in \mathbb{P}$ such that $p \neq q$. Let $m = p \cdot q$ and $v = \text{lcm}(p-1, q-1)$; then $\forall a, k \in \mathbb{Z}: a^{kv+1} \equiv a(m)$

3.5 RSA-Algorithm

- (m, e) ... public key
- (m, d) ... private key
- Encryption:** $E(a_j) = b_j = a_j^e \pmod{m}$
- Decryption:** $D(b_j) = a_j = b_j^d \pmod{m}$
- # number of fixpoints: $\gcd(e-1, p-1) \cdot \gcd(e-1, q-1)$

Groups and Cyclic Groups

- Group:** Let G be a group and $x \in G$, then $\text{ord}_G(x) = \min\{i \in \mathbb{N}^+ \mid x^i = e\}$. If e is the neutral element, then $\text{ord}_G(e) = 1$ since $e^1 = e$. If $a \neq e$, then $\text{ord}_G(a) > 1$
- Cyclic Group:** Let $\langle x \rangle$ denote the group generated by x . Examples $\langle e \rangle = \{e\}$ is the trivial group
- $\langle x \rangle = \{e, x, x^2, x^3, \dots\}$
- \mathbb{Z}_m is cyclic
- $\Leftrightarrow \exists$ primitive root mod m
- $\Leftrightarrow m \in \{2, 4\} \cup \{p^k \mid p \in \mathbb{P} \setminus \{2\}, k \geq 1\} \cup \{2p^k \mid p \in \mathbb{P} \setminus \{2\}, k \geq 1\}$
- Primitive Root mod m :** $\bar{a} \in \mathbb{Z}_m^*$, then $(\bar{a}) = \mathbb{Z}_m^* \Rightarrow \mathbb{Z}_m^* = \{\bar{a}, \bar{a}^2, \bar{a}^3, \dots, \bar{a}^{\varphi(m)}\}$ (Hint: $\bar{a}^{\varphi(m)} = 1$)

The Order of Elements

- $\text{ord}_G(x) = |\langle x \rangle|$; Obviously $\text{ord}_G(x)$ divides $|G|$, since $\langle x \rangle \subseteq G$
- Let $\text{ord}(a) = r$; then $\text{ord}(a^s) = \frac{r}{\gcd(r, s)}$
 - $\exists a, b \in G: \text{ord}_G(a) = r, \text{ord}_G(b) = s \Rightarrow \exists c \in G: \text{ord}_G(c) = \text{lcm}(r, s)$
 - From the above follows: $\text{ord}_G(a)$ maximal $\Leftrightarrow \forall b \in G: \text{ord}_G(b) \mid \text{ord}_G(a)$
 - In general: a is generator of finite field $K \setminus \{0\}$ $\Leftrightarrow \bar{a} \in K \setminus \{0\} \mid \Lambda K \setminus \{0\} = \{\bar{a}, \bar{a}^2, \bar{a}^3, \dots, \bar{a}^{\text{ord}(K)-1}\}$
 - Let a_1, \dots, a_r