

# Discrete Mathematics

Kellerkinder

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## 1 Graph Theory

### 1.1 Definition & Notation

- $G = (V, E)$
- Directed Edge  $e = (v, w): v \rightarrow w$
- Undirected Edge  $e = \{v, w\} = \{w, v\} = vw$
- $d(v)$ : degree
- $d^+(v)$ : out-degree
- $d^-(v)$ : in-degree
- $\Gamma(v)$ : set of neighbours
- $\Gamma^+(v)$ : successors
- $\Gamma^-(v)$ : predecessors

Walk, Trail Circuit

- Walk: sequence of edges (no jumps)
- 3 walk  $v: v \sim w$
- Trail: sequence of edges (no jumps, no repeating edges)
- Circuit/Closed Trail: same start and end vertex
- Path: no repeating edges, no repeating vertices (colors  $c_1, \dots, c_r$ ) have a  $c_j$ -colored  $K_{n_j}$  for some  $j$
- Cycle/Closed Path: same start and end vertex

### Handshaking Lemma

$$\sum_{v \in V} d(v) = 2|E| \quad \sum_{v \in V} d^+(v) = \sum_{v \in V} d^-(v) = |E|$$

### 1.2 Trees & Forests

$G \cong H \iff \exists \varphi: V(G) \mapsto V(H) \text{ and } vw \in E(G) \iff \varphi(v)\varphi(w) \in E(H)$

Spanning Subgraphs & Spanning Forests

Spanning Forest  $\iff$

1.  $V(F) = V(G)$  and  $E(F) \subseteq E(G)$

2.  $F$  is a forest

3.  $F$  has the same connected components as  $G$

If  $G$  connected  $\Rightarrow F$  spanning tree

Minimum or Maximum Spanning Trees

•  $G = (V, E)$  and weight function  $w: E \rightarrow \mathbb{R}, e \mapsto w(e)$

Then  $T = (V, F) \subseteq G$  minimum spanning tree if  $w(F)$  minimal

• GREEDY( $V, E, w$ ) constructs the minimum spanning tree (MST):

- Sorts the edges of  $G$  by weight ascending
- Adds the next smallest edge to  $T$  (such that  $T$  remains a tree) until  $G$  is spanned

Matrix-Tree-Theorem (Kirchhoff)

$G = (V = \{v_1, \dots, v_n\}, E)$  simple, undirected, connected, with adjacency matrix  $A = \text{diag}(d(v_1), \dots, d(v_n))$ ;  $(D - A)^{-1} = (D - A)$  without 1 row and 1 column  $\Rightarrow \det((D - A)^{-1})$  = number of spanning trees of  $G$

Independence System  $(E, S) \subseteq 2^E$  and  $S$  is closed under inclusion.

$A \subseteq S \wedge B \subseteq A \Rightarrow B \in S$

Matroids & GREEDY

An independence system  $M = (E, S)$  is called a matroid if  $\forall A, B \in S$  such that  $|B| = |A| + 1 \Rightarrow \exists e \in B, A \cup \{e\} \in S$

Let  $G = (V, E)$  an undirected graph and  $S = \{F \subseteq E \mid F \text{ is a forest}\}$ ; then  $(E, S)$  is a matroid.

If  $M = (E, S)$  is a matroid with a weight function  $w: E \rightarrow \mathbb{R}$  then GREEDY

computes a  $A$  maximal/minimal weight with respect to the inclusion.

### 1.3 Special Graphs

Planar Graphs

Definition: A graph  $G$  is planar if there is an isomorphic graph  $H$  embedded in the plane (vertices are points in the plane  $\mathbb{R}^2$ ) such that no edges intersect. A graph is planar  $\iff$  there is no subgraph which is a subdivision of  $K_5$  or of  $K_3,3$ . Euler's Polyhedron Formula: if  $G$  is connected and planar graph, then  $\alpha_0 - \alpha_1 + \alpha_2 = 2 - \alpha_0$ . Vertices,  $\alpha_0$ : Edges,  $\alpha_2$ : Faces

Dual Graph: Let  $G = (V, E)$  be a planar graph and let  $F$  be the set of its faces. Then  $G^* = (V^*, E^*)$  is defined such that  $V^* = F$  and for every edge  $e \in E$ , set  $e^* = (f_1, f_2)$  if  $f_1$  and  $f_2$  are the faces left and right of  $e$ .

$G^*$  is called the dual of  $G$ .

Bipartite Graphs and Matching

Definition: Let  $G = (V, E)$  be a simple undirected graph.  $G$  is called bipartite if and only if

- $V = V_1 \cup V_2, V_1 \cap V_2 = \emptyset$
- $vw \in E \Rightarrow v \in V_1, w \in V_2$

Hall's Marriage Theorem: Given a bipartite graph  $G = (V, E)$ , such that  $V = W \cup M$  where  $W$  and  $M$  are finite and nonempty. Define the friendship relation:  $F \subseteq W \times M$  with  $vw \in F \iff \exists y \in M$  such that  $xFy$ . Now the theorem states the following: there is a feasible marriage, if and only if:  $\forall W_0 \subseteq W: |\{y \in M \mid \exists x \in W_0: xFy\}| \geq |W_0|$  or in different words:  $|\bigcup_{w \in W_0} \Gamma(w)| \geq |W_0|$

### 1.4 Graph Colorings

• Vertex Coloring:  $c: V \rightarrow C$  where a  $C$  is a set of colors. feasible if  $vw \in E \Rightarrow c(v) \neq c(w)$

• Edge Coloring:  $\bar{c}: E \rightarrow C$

• everything that can be done with a vertex coloring can also be done with an edge coloring

- Chromatic Number:  $\chi(G)$  is the minimal number of colors such that there is a feasible coloring.
  - $\chi(K_n) = n$
  - $\chi(K_{n,m}) = 2$
  - $\chi(T) = 2$  if  $T$  is a tree and  $|V| > 1$
  - $\chi(G) = 1 \iff E(G) = \emptyset$
  - $\chi(G) = 2 \iff E(G) \neq \emptyset$  and  $G$  is bipartite
  - $\chi(G) = 2 \iff E(G) \neq \emptyset$  and all cycles have even length
  - $\chi(G) \leq 1 + \max_{v \in V} d(v)$
  - $G$  is planar  $\iff \chi(G) \leq 5$

- Ramsey Theory
- Definition: The Ramsey Number  $R(r, s)$  is the minimum  $n$  such that every red-blue coloring of  $K_n$  contains either a red  $K_r$  or a blue  $K_s$ .
  - Coloring is edge coloring
  - $R(r, s) \leq R(r-1, s) + R(r, s-1)$
  - $R(n_1, n_2, \dots, n_p) = \min\{n \mid \text{all } r\text{-edge colorings of } K_n \text{ have a } c_j\text{-colored } K_{n_j} \text{ for some } j\}$

### 2 Advanced Combinatorics

#### Double Counting

Given two sets  $A$  and  $B$  and a relation  $R \subseteq A \times B$ , such that  $aRb \iff (a, b) \in R$ : Let  $R_i = \{b \in B \mid aRb\}$  (all  $b$  related to  $a$ ) and  $S_j = \{a \in A \mid aRb_j\}$ . Then:  $|R| = \sum_i |R_i| = \sum_j |S_j|$

Pigeon Hole Principle

Let  $A_1, \dots, A_k$  be finite pairwise disjoint sets,  $|A_1 \cup \dots \cup A_k| > k \cdot r$ , for  $r \in \mathbb{N}$ ; This implies:  $|A_i| > r$ , if  $i = 1$ , then it follows that:  $f: A \mapsto B, A \mapsto B \Rightarrow \exists b \in B: |f^{-1}(b)| \geq r$

Principle of Inclusion & Exclusion

$$|A_1 \cup \dots \cup A_k| = \sum_{i=1}^k |A_i| - \sum_{1 \leq i < j \leq k} (-1)^{i+j} |\bigcap_{l=1}^k A_l|$$

For  $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$  If the sets are not necessarily pairwise disjoint, then:

$$A_1, \dots, A_n \subseteq A \setminus \bigcap_{i=1}^n A_i = |A| + \sum_{1 \leq i < j \leq n} (-1)^{|i|+|j|} |\bigcap_{l=1}^n A_l|$$

### 2.1 Enumerative Combinatorics

#### Counting Sets

Useful theorems:

$$\begin{aligned} 1. \# \text{ of permutations: } n! \\ 2. \# \text{ of } k\text{-subsets of } \binom{[n]}{k} \\ 3. \# \text{ of ordered } k\text{-subsets: } k! \binom{n}{k} \\ 4. \# \text{ of } k\text{-multisets (elements can be used more than once): } \binom{n+k-1}{k} \\ 5. \# \text{ of arrangements of a multiset (element } b_i \text{ appears } k_i \text{ times): } \frac{n!}{k_1! k_2! \dots k_n!} \\ 6. \# \text{ of ordered } k\text{-multisets: } n^k \end{aligned}$$

Stars and bars: Number of assignments of  $n$  items into  $k$  boxes

$$\binom{n+k-1}{k-1}$$

#### Stirling Numbers

First kind:  $s_{n,k} = \# \text{ of permutations of } \{1, \dots, n\} \text{ with } k \text{ cycles}$

$$\begin{aligned} \bullet s_{n,1} = (n-1)!, \quad s_{n,n-1} = \binom{n}{2}, \quad s_{n,n} = 1 \\ \bullet s_{n,0} = 0, \quad s_{0,0} = 0, \quad s_{0,1} = 1 \quad (!!) \end{aligned}$$

$$\sum_{k=1}^n s_{n,k} = n!$$

$$\forall n, k > 0: s_{n,k} = s_{n-1,k-1} + (n-2)$$

$$\forall n, k > 0: s_{n,k} = s_{n-1,k-1} + k s_{n-1,k}$$

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