

<b>Status</b>	Beendet
<b>Begonnen</b>	Montag, 26. Mai 2025, 15:27
<b>Abgeschlossen</b>	Montag, 26. Mai 2025, 15:41
<b>Dauer</b>	14 Minuten 7 Sekunden
<b>Bewertung</b>	10,00 von 10,00 (100%)

**Frage 1**

Richtig

Erreichte Punkte 1,00 von 1,00

Let

$$t : \forall x \exists y (P(x, g(y)) \vee Q(g(x), f(c, b)))$$

be the starting formula of a tableau branch. Which of the following terms are suitable substitutes for the variable  $x$  in the subsequent steps of a tableau proof, when applying the corresponding quantifier rule?

**Note:** Multiple answers may be correct.

- $b$  ✓
- $f(b, c)$  ✓
- $f(b, b)$  ✓
- $e$  ✓

**Frage 2**

Richtig

Erreichte Punkte 1,00 von 1,00

In this question we consider tableau proofs with only one branch and we will represent them as a sequence. That is, the sequence  $(t : P(a), f : \neg P(a))$  corresponds to the tableau proof:

$$t : P(a)$$

$$f : \neg P(a)$$

Considering this, which one of the following tableau proofs are closed? **Multiple** answers may be correct.

- a.  $(t : \forall x (Q(x) \rightarrow R(x)), t : \forall x R(x), t : Q(a), f : R(a), t : R(a))$  ✓
- b.  $(t : \forall x (Q(x) \rightarrow R(x)), f : \exists x \neg R(x), t : Q(a), t : \neg R(a), f : \neg \neg R(b))$
- c.  $(t : \forall x (Q(x) \vee R(x)), t : Q(b), t : \forall x \neg \neg R(x), f : \neg R(a), t : \neg R(a))$  ✓
- d.  $(t : \forall x (Q(x) \vee R(x)), t : Q(b), f : \exists x \neg \neg R(x), t : Q(b), t : \neg R(a), f : \neg R(a))$  ✓

**Frage 3**

Richtig

Erreichte Punkte 1,00 von 1,00

Which is the formula in First-Order Logic that best captures the natural language sentence "Liam teaches civics and geography."?

- a.  $Teaches(liam, civics) \wedge Teaches(liam, geography)$  ✓
- b.  $Liam(teaches, civics \wedge geography)$
- c.  $Liam(teaches, civics) \wedge Liam(teaches, geography)$
- d.  $Teaches(teaches, civics \wedge geography)$

**Frage 4**

Richtig

Erreichte Punkte 1,00 von 1,00

Let  $A$  be the  $\mathcal{T}_E$ -formula  $a = b \vee a \neq b$ . Which of the following statements is true?

- a.  $A$  is satisfiable but not valid in  $\mathcal{T}_E$ .
- b.  $A$  is valid in  $\mathcal{T}_E$ . ✓
- c.  $A$  is not satisfiable in  $\mathcal{T}_E$ .

**Frage 5**

Richtig

Erreichte Punkte 1,00 von 1,00

Which one of the formulas below is a prenex normal form of  $\neg\exists y((\forall yR(y) \wedge P(b)) \rightarrow (\exists xP(x, y) \vee \neg R(y)))$ ?

- a.  $\forall x_0 \forall x_1 \forall x_2 ((R(x_1) \vee P(b)) \wedge (\neg P(x_2, x_0) \wedge R(x_0)))$
- b.  $\forall x_0 \forall x_1 \forall x_2 ((R(x_1) \wedge P(b)) \wedge (\neg P(x_2, x_0) \wedge R(x_0)))$
- c.  $\forall x_0 \forall x_1 \forall x_2 ((R(x_1) \wedge P(b)) \wedge (\neg P(x_2, x_0) \wedge R(x_0)))$  ✓
- d.  $\forall x_0 \exists x_1 \exists x_2 ((\neg R(x_1) \vee \neg P(b)) \vee (P(x_2, x_0) \vee \neg R(x_0)))$

**Frage 6**

Richtig

Erreichte Punkte 1,00 von 1,00

What is the first-order formula which best captures the following natural language sentence: "If there is a dog that barks loudly, then the neighbor will be annoyed."?

- a.  $\forall x (Dog(x) \wedge BarksLoudly(x) \rightarrow Annoyed(neighbor))$
- b.  $\exists x (Dog(x) \wedge BarksLoudly(x) \rightarrow Annoyed(neighbor))$  ✓
- c.  $Annoyed(neighbor) \rightarrow \exists x (Dog(x) \wedge BarksLoudly(x))$
- d.  $\exists x (Dog(x) \wedge BarksLoudly(x) \rightarrow Annoyed(neighbor))$

**Frage 7**

Richtig

Erreichte Punkte 1,00 von 1,00

Let  $\mathcal{A}$  be the  $\mathcal{L}$ -formula  $(f(b) \neq b \wedge f(a) = a \wedge b = f(a))$ . Which are the congruence classes in the congruence closure of  $\{=\}$  over  $\mathcal{A}$ ?

- a.  $\{[a], [f(b), f(a), b]\}$
- b.  $\{[a, b], [f(b), f(a)]\}$
- c.  $\{[a], [b], [f(b), f(a)]\}$
- d.  $\{[a, f(b), f(a), b]\}$  ✓

**Frage 8**

Richtig

Erreichte Punkte 1,00 von 1,00

Consider the formula  $(x \leq c - 1)$  to be interpreted in the standard way over integers, with  $x$  denoting a variable and  $c$  being a constant. Let  $\mathcal{I}$  be an interpretation such that  $\mathcal{I}(x) = \mathcal{I}(c) = 0$ .

- a.  $\mathcal{I}$  is a model of  $(x \leq c - 1)$
- b.  $(x \leq c - 1)$  is satisfiable in  $\mathcal{I}$  ✓
- c.  $(x \leq c - 1)$  is unsatisfiable
- d.  $(x \leq c - 1)$  is valid

**Frage 9**

Richtig

Erreichte Punkte 1,00 von 1,00

Which one of the following is a suitable new tableau rule that could be applied to formulas of the form  $\neg \exists x \neg F(x)$  in tableau proofs?

- a.  $\frac{\neg \exists x \neg F(x)}{F(c/x)} \sim \text{for } c \text{ a new constant}$  ✓
- b.  $\frac{\neg \exists x \neg F(x)}{F(c/x)} \sim \text{for } c \text{ a new constant}$
- c.  $\frac{\neg \exists x \neg F(x)}{F(t/x)} \sim \text{for } t \text{ a variable-free term}$
- d.  $\frac{\neg \exists x \neg F(x)}{F(t/x)} \sim \text{for } t \text{ a variable-free term}$

**Frage 10**

Richtig

Erreichte  
Punkte 1,00 von  
1,00

Let  $F$  be the formula  $\forall x \exists y (x \neq y \rightarrow y = x)$ , where  $(x, y)$  are variables. Let  $(\alpha)$  be the sort of  $(x, y)$ .

- a.  $(F)$  is valid in the class of interpretations that interpret  $(\alpha)$  to be the domain  $(D = \{1\})$  ✓
- b.  $(F)$  is unsatisfiable in the class of interpretations that interpret  $(\alpha)$  to be the domain  $(D = \{1\})$
- c.  $(F)$  is satisfiable but not valid in the class of interpretations that interpret  $(\alpha)$  to be the domain  $(D = \{1\})$