

## 7.1.

- Compute the number of additions and multiplications in Algorithms 4.2 and 4.3 (forward and backward substitution).
- Show that the product  $\mathbf{L}_1\mathbf{L}_2$  of two lower triangular matrices  $\mathbf{L}_1, \mathbf{L}_2$  is again a lower triangular matrix. Also show that the inverse of a (invertible) lower triangular matrix is lower triangular.

**Solution.**

**Algorithm 4.2 (solve  $\mathbf{Lx} = \mathbf{b}$  using *forward substitution*)**

*Input:*  $\mathbf{L} \in \mathbb{R}^{n \times n}$  lower triangular, invertible,  $\mathbf{b} \in \mathbb{R}^n$

*Output:* solution  $\mathbf{x} \in \mathbb{R}^n$  of  $\mathbf{Lx} = \mathbf{b}$

**for**  $j = 1:n$  **do**

$$x_j := \left( b_j - \sum_{k=1}^{j-1} l_{jk}x_k \right) / l_{jj}$$

[[ convention: empty sum = 0 ]]

**end for**

a)

Number of additions:  $n$  loop iterations of  $\max(j-2, 0)$  additions, where  $j$  is the current loop index, for a total of  $\sum_{j=2}^n j-2 = \frac{1}{2}(n^2 - 3n + 2)$  additions.

Number of multiplications:  $n$  loop iterations of  $j-1$  multiplications, where  $j$  is the current loop index, for a total of  $\sum_{j=1}^n j-1 = \frac{1}{2}(n-1)n$  multiplications.

**Algorithm 4.3 (solution of  $\mathbf{Ux} = \mathbf{b}$  using *back substitution*)**

*Input:*  $\mathbf{U} \in \mathbb{R}^{n \times n}$  upper triangular, invertible,  $\mathbf{b} \in \mathbb{R}^n$

*Output:* solution  $\mathbf{x} \in \mathbb{R}^n$  of  $\mathbf{Ux} = \mathbf{b}$

**for**  $j = n:-1:1$  **do**

$$x_j := \left( b_j - \sum_{k=j+1}^n u_{jk}x_k \right) / u_{jj}$$

**end for**

Number of additions:  $n$  loop iterations of  $\max(n-(j+1), 0)$  additions, where  $j$  is the current loop index, for a total of  $\sum_{j=1}^{n-1} n-(j+1) = \frac{1}{2}(n-3)n$  additions.

Number of multiplications:  $n$  loop iterations of  $n-j$  multiplications, where  $j$  is the current loop index, for a total of  $\sum_{j=1}^n n-j = \frac{1}{2}(n-1)n$  multiplications.

- We show product of lower triangular matrices is triangular:

$$\text{for } j > i : (\mathbf{L}_1\mathbf{L}_2)_{ij} = \sum_{k=1}^n \underbrace{\mathbf{L}_{1ik}}_{0 \text{ for } k > i} \underbrace{\mathbf{L}_{2kk}}_{0 \text{ for } j > k} = 0$$

Since  $k \geq j \implies k > i$ .

We show inverse of triangular is triangular:

Let  $k \in \{1, \dots, n\}$  arbitrary, we consider multiplication of  $\mathbf{L}$  with the  $k$ -th column of the inverse, which has to result in the  $k$ -th unit vector:

$$\begin{pmatrix} \mathbf{L}_{11} & 0 & \cdots & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & & & \vdots \\ \mathbf{L}_{(k-1)1} & & & 0 & & 0 \\ \mathbf{L}_{k1} & & & \mathbf{L}_{kk} & \ddots & 0 \\ \vdots & & & & \ddots & \vdots \\ \mathbf{L}_{(n-1)1} & & & & & \ddots & 0 \\ \mathbf{L}_{n1} & \cdots & \cdots & \mathbf{L}_{kk} & \cdots & \mathbf{L}_{nn} \end{pmatrix} \begin{pmatrix} \mathbf{L}_{1k}^{-1} \\ \vdots \\ \mathbf{L}_{n(k-1)}^{-1} \\ \mathbf{L}_{nk}^{-1} \\ \vdots \\ \mathbf{L}_{(n-1)k}^{-1} \\ \mathbf{L}_{nk}^{-1} \end{pmatrix} = \begin{pmatrix} 0 = \mathbf{L}_{11} \cdot \mathbf{L}_{1k}^{-1} \\ \vdots \\ 0 = \sum_{j=1}^{k-1} \mathbf{L}_{1j} \cdot \mathbf{L}_{jk}^{-1} \\ 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} \begin{matrix} \implies \mathbf{L}_{1k}^{-1} = 0 \\ \implies \mathbf{L}_{2k}^{-1}, \dots, \mathbf{L}_{(k-2)k}^{-1} = 0 \\ \implies \mathbf{L}_{(k-1)k}^{-1} = 0 \end{matrix}$$

Therefore in the  $k$ -th column,  $\mathbf{L}_{ik}^{-1} = 0$  for  $k > i$ . Since  $k$  was chosen arbitrarily, this means that  $\mathbf{L}^{-1}$  has to be lower triangular.

□

## 7.2

- Explain Crout's algorithm from Chapter 4.3.1 in the lecture notes
- Modify the algorithm to compute a Cholesky factorization

$$\mathbf{C}^T \mathbf{C} = \mathbf{A}$$

and realize your algorithm in Matlab/Python.

### **Solution.**

- We want to solve a system of  $n^2$  equations for  $n^2$  unknowns, and in order to do so traverse these equations in a specific order ("Crout ordering"). This involves alternating between row and column traversal: first we go through the first row, then the first column, then the second row, .. This is done to allow our equations to either consist of lots of zeros and/or elements already computed in previous steps.
- See Python.

□

## 7.4.

Consider the matrix

$$\mathbf{A} = \begin{pmatrix} 10^{-6} & 1 \\ 1 & 0 \end{pmatrix}$$

- Compute by hand the  $LU$ -factorization of  $\mathbf{A}$ . Calculate  $\mathbf{A}^{-1}$ ,  $\mathbf{L}^{-1}$ ,  $\mathbf{U}^{-1}$  and compute the three condition numbers  $\kappa_{\infty}(\mathbf{A}) = \|\mathbf{A}\|_{\infty} \|\mathbf{A}^{-1}\|_{\infty}$ ,  $\kappa_{\infty}(\mathbf{L}) = \|\mathbf{L}\|_{\infty} \|\mathbf{L}^{-1}\|_{\infty}$ ,  $\kappa_{\infty}(\mathbf{U}) = \|\mathbf{U}\|_{\infty} \|\mathbf{U}^{-1}\|_{\infty}$ . Here,  $\|\cdot\|_{\infty}$  is the row-sum norm

- b) Repeat the calculation of a) for the matrix  $\tilde{\mathbf{A}}$  that is obtained from  $\mathbf{A}$  by interchanging the two rows. What do you observe?

**Solution.**

a)

$$\mathbf{A} = \begin{pmatrix} 10^{-6} & 1 \\ 1 & 0 \end{pmatrix} = \mathbf{L}\mathbf{U} \stackrel{CROUT}{=} \begin{pmatrix} 1 & 0 \\ \mathbf{L}_{21} & 1 \end{pmatrix} \begin{pmatrix} 10^{-6} & 1 \\ 0 & \mathbf{U}_{22} \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 1/10^{-6} & 1 \end{pmatrix} \begin{pmatrix} 10^{-6} & 1 \\ 0 & -10^6 \end{pmatrix}$$

Solve

$$\mathbf{A}\mathbf{A}^{-1} = I, \begin{pmatrix} 10^{-6} & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \mathbf{A}_{11}^{-1} \\ \mathbf{A}_{21}^{-1} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 10^{-6} & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \mathbf{A}_{12}^{-1} \\ \mathbf{A}_{22}^{-1} \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \mathbf{A}^{-1} = \begin{pmatrix} 0 & 1 \\ 1 & -10^{-6} \end{pmatrix}$$

$$\mathbf{L}\mathbf{L}^{-1} = I, \begin{pmatrix} 1 & 0 \\ 10^6 & 1 \end{pmatrix} \begin{pmatrix} \mathbf{L}_{11}^{-1} \\ \mathbf{L}_{21}^{-1} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 10^6 & 1 \end{pmatrix} \begin{pmatrix} \mathbf{L}_{12}^{-1} \\ \mathbf{L}_{22}^{-1} \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \mathbf{L}^{-1} = \begin{pmatrix} 1 & 0 \\ -10^6 & 1 \end{pmatrix}$$

$$\mathbf{U}\mathbf{U}^{-1} = I, \begin{pmatrix} 10^{-6} & 1 \\ 0 & -10^6 \end{pmatrix} \begin{pmatrix} \mathbf{U}_{11}^{-1} \\ \mathbf{U}_{21}^{-1} \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 10^{-6} & 1 \\ 0 & -10^6 \end{pmatrix} \begin{pmatrix} \mathbf{U}_{12}^{-1} \\ \mathbf{U}_{22}^{-1} \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix},$$

$$\mathbf{U}^{-1} = \begin{pmatrix} 10^6 & 1 \\ 0 & -10^{-6} \end{pmatrix}$$

$$\kappa_{\infty}(\mathbf{A}) = \|\mathbf{A}\|_{\infty} \|\mathbf{A}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |a_{ij}| \max_i \sum_{j=1}^n |a_{ij}^{-1}| = (10^{-6} + 1)(1 + 10^{-6})$$

$$\kappa_{\infty}(\mathbf{L}) = \|\mathbf{L}\|_{\infty} \|\mathbf{L}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |l_{ij}| \max_i \sum_{j=1}^n |l_{ij}^{-1}| = (10^6 + 1)(10^6 + 1)$$

$$\kappa_{\infty}(\mathbf{U}) = \|\mathbf{U}\|_{\infty} \|\mathbf{U}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |u_{ij}| \max_i \sum_{j=1}^n |u_{ij}^{-1}| = (0 + 10^6)(10^6 + 1)$$

b)

$$\tilde{\mathbf{A}} = \begin{pmatrix} 1 & 0 \\ 10^{-6} & 1 \end{pmatrix} = \tilde{\mathbf{L}}\tilde{\mathbf{U}} = \begin{pmatrix} 1 & 0 \\ 10^{-6} & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

Same as in a) we get

$$\tilde{\mathbf{A}}^{-1} = \tilde{\mathbf{A}}$$

$$\tilde{\mathbf{L}}^{-1} = \begin{pmatrix} 1 & 0 \\ -10^{-6} & 1 \end{pmatrix}$$

$$\tilde{\mathbf{U}}^{-1} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\kappa_{\infty}(\tilde{\mathbf{A}}) = \|\tilde{\mathbf{A}}\|_{\infty} \|\tilde{\mathbf{A}}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |\tilde{a}_{ij}| \max_i \sum_{j=1}^n |\tilde{a}_{ij}^{-1}| = (10^{-6} + 1)(10^{-6} + 1) \approx 1$$

$$\kappa_{\infty}(\tilde{\mathbf{L}}) = \|\tilde{\mathbf{L}}\|_{\infty} \|\tilde{\mathbf{L}}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |\tilde{l}_{ij}| \max_i \sum_{j=1}^n |\tilde{l}_{ij}^{-1}| = (10^{-6} + 1)(10^{-6} + 1) \approx 1$$

$$\kappa_{\infty}(\tilde{\mathbf{U}}) = \|\tilde{\mathbf{U}}\|_{\infty} \|\tilde{\mathbf{U}}^{-1}\|_{\infty} = \max_i \sum_{j=1}^n |\tilde{u}_{ij}| \max_i \sum_{j=1}^n |\tilde{u}_{ij}^{-1}| = (0 + 1)(0 + 1) = 1$$

□

## 7.5.

Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be a tridiagonal matrix of the form

$$\mathbf{A} = \begin{pmatrix} d_1 & e_1 & & & \\ c_2 & d_2 & e_2 & & \\ & \ddots & \ddots & \ddots & \\ & & \ddots & \ddots & e_{n-1} \\ & & & c_n & d_n \end{pmatrix}$$

Assume that  $\mathbf{A}$  has an  $LU$ -factorization.

- a) Show: the factors  $\mathbf{L}$  and  $\mathbf{U}$  have the form

$$\mathbf{L} = \begin{pmatrix} 1 & & & & \\ l_2 & 1 & & & \\ & l_3 & 1 & & \\ & & \ddots & \ddots & \\ & & & l_n & 1 \end{pmatrix}, \mathbf{U} = \begin{pmatrix} u_1 & f_1 & & & \\ & u_2 & f_2 & & \\ & & \ddots & \ddots & \\ & & & \ddots & f_{n-1} \\ & & & & u_n \end{pmatrix}$$

*Hint:* do the induction step of Thm. 4.17 of the notes

- b) Formulate an algorithm that computes the  $l_i$  and the  $u_i$  for  $i = 2, \dots, n$  and realize your algorithm in Matlab/Python. Input are the vectors  $\mathbf{d}, \mathbf{e}, \mathbf{c}$  (i.e., the diagonals of  $\mathbf{A}$ ), output are the vectors  $\mathbf{l}, \mathbf{u}$ , and  $\mathbf{f}$  (i.e., the diagonals of  $\mathbf{L}$  and  $\mathbf{U}$ ).

**Solution.**

- a) On paper  
b) see Python

□