

FMSP: Solution Attempt for exam 2021-09

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Preamble

Feel free to edit it to add more examples, fix mistakes, etc. If you do so, please update both the .tex and .pdf files. And obviously no guarantee for anything.

Problem 1

(a)

$$\begin{aligned} P &= (\nu k_1: \text{Sym}K^{HH}[HH])(\nu k_2: \text{Sig}K^{HH}[LH, LL])A|B \\ A &= (\nu m: LL)(\nu n: HH)\bar{c}\langle\{|n|\}_{k_1}^s, m\rangle.0 \\ B &= c(x). \text{ case } x \text{ of } [y, m']_{\nu k(k_2)} \text{ in case } y \text{ of } \{|n'\}_{k_1}^s \text{ in } 0 \end{aligned}$$

(b)

We will let $\Gamma = \{c: LL\}$.

Restrictions and parallel .

$$\frac{\frac{\Gamma, k_1: \text{Sym}K^{HH}[HH], k_2: \text{Sig}K^{HH}[LH, LL] \vdash A \quad \Gamma, k_1: \text{Sym}K^{HH}[HH], k_2: \text{Sig}K^{HH}[LH, LL] \vdash B}{\Gamma, k_1: \text{Sym}K^{HH}[HH], k_2: \text{Sig}K^{HH}[LH, LL] \vdash A|B} \text{PAR}}{\frac{\Gamma, k_1: \text{Sym}K^{HH}[HH], k_2: \text{Sig}K^{HH}[LH, LL] \vdash A|B}{\Gamma, k_1: \text{Sym}K^{HH}[HH] \vdash (\nu k_2: \text{Sig}K^{HH}[LH, LL])A|B} \text{RES}}{\Gamma \vdash (\nu k_1: \text{Sym}K^{HH}[HH])(\nu k_2: \text{Sig}K^{HH}[LH, LL])A|B} \text{RES}}$$

For simplicity, let

$$\Gamma_1 = \Gamma, k_1: \text{Sym}K^{HH}[HH], k_2: \text{Sig}K^{HH}[LH, LL]$$

Proving A .

With

$$\Gamma_2 = \Gamma_1, m: LL, n: HH$$

$$\begin{array}{c}
\frac{(2)}{\Gamma_2 \vdash [\{|n\}_{k_1}^s, m]_{k_2} : LL} \quad \frac{(0.2)}{\Gamma_2 \vdash \diamond} \text{ STOP} \quad \frac{(4)}{\Gamma_2 \vdash c : C^{LL}[LL]} \\
\hline
\text{OUT} \\
\frac{\Gamma_1, m : LL, n : HH \vdash \bar{c}(\{|n\}_{k_1}^s, m).0}{\Gamma_1, m : LL \vdash (\nu n : HH)\bar{c}(\{|n\}_{k_1}^s, m).0} \text{ RES} \\
\frac{\Gamma_1 \vdash (\nu m : LL)(\nu n : HH)\bar{c}(\{|n\}_{k_1}^s, m).0}{\Gamma_1 \vdash (\nu m : LL)(\nu n : HH)\bar{c}(\{|n\}_{k_1}^s, m).0} \text{ RES} \\
\hline
(1)
\end{array}$$

$$\begin{array}{c}
\frac{(0.2)}{\Gamma_2 \vdash \diamond} \quad \frac{c : LL \text{ in } \Gamma_2}{\Gamma_2 \vdash c : LL} \text{ ATOM} \quad LL \leq C^{LL}[LL] \\
\hline
\Gamma_2 \vdash c : C^{LL}[LL] \text{ SUBSUMPTION} \\
\hline
(4)
\end{array}$$

Doing the signature .

$$\begin{array}{c}
\frac{(0.2)}{\Gamma_2 \vdash \diamond} \quad k_2 : \text{Sig}K^{HH}[LH, LL] \text{ in } \Gamma_2 \quad \text{ATOM} \quad \frac{(5)}{\Gamma_2 \vdash [\{|n\}_{k_1}^s, m] : [LH, LL]} \quad L = \dots \\
\hline
\Gamma_2 \vdash k_2 : \text{Sig}K^{HH}[LH, LL] \quad \frac{\Gamma_2 \vdash [\{|n\}_{k_1}^s, m]_{k_2} : LH}{\Gamma_2 \vdash [\{|n\}_{k_1}^s, m]_{k_2} : LH} \text{ DigSig} \quad LH \leq LL \\
\hline
\Gamma_2 \vdash [\{|n\}_{k_1}^s, m]_{k_2} : LL \text{ SUBSUMPTION} \\
\hline
(2)
\end{array}$$

Doing the encryption .

$$\begin{array}{c}
\frac{(0.2)}{\Gamma_2 \vdash \diamond} \quad k_1 : \text{Sym}K^{HH}[HH] \text{ in } \Gamma_2 \quad \text{ATOM} \quad \frac{(0.2)}{\Gamma_2 \vdash \diamond} \quad n : HH \text{ in } \Gamma_2 \quad \text{ATOM} \quad \frac{(0.2)}{\Gamma_2 \vdash \diamond} \quad m : LL \text{ in } \Gamma_2 \quad \text{ATOM} \\
\hline
\Gamma_2 \vdash k_1 : \text{Sym}K^{HH}[HH] \quad \Gamma_2 \vdash n : HH \quad \text{SymEnc} \quad \Gamma_2 \vdash m : LL \\
\hline
\Gamma_2 \vdash \{|n\}_{k_1}^s : LH \quad \Gamma_2 \vdash m : LL \text{ LIST} \\
\hline
\Gamma_2 \vdash [\{|n\}_{k_1}^s, m] : [LH, LL] \\
\hline
(5)
\end{array}$$

Proving B .

For simplicity we let

$$Q = \text{case } y \text{ of } \{|n'\}_{k_1}^s \text{ in } 0$$

Proving the wellformedness The $\Gamma_1 \vdash \diamond$, $\Gamma_2 \vdash \diamond$ and $\Gamma_3 \vdash \diamond$ hold, because all channels and keys in it are of type HH .

Conclusion If I didn't forget about any branch, this should be finished and typechecks. We have $img(\Gamma) = LL$, thus it preserves secrecy and integrity ☺